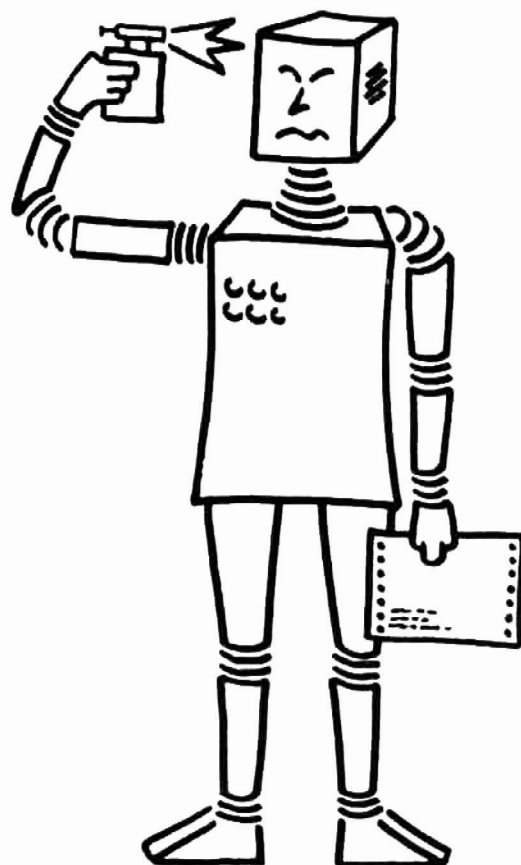


# SEKI-PROJEKT

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Programming in a Distributed  
Environment : a Collection  
of CSSA Examples

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**Abstract :**

Several CSSA scripts as solutions of mostly well known problems in the area of parallel activities are documented. All but two examples are compared with "equivalent" programs which are based on the concept of monitors. One example gives a distributed solution of the N-queens problem. Besides the purpose of demonstrating the use of CSSA concepts, the last example deserves particular attention. Here, CSSA is used for simulating and testing the logic of the "window mechanism" which is a representative for the class of communication protocols in computer networks.

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## 0. Introduction

This paper is a documentation of my first programming experiences with CSSA. The study of the examples offers an opportunity to deepen the understanding of CSSA concepts and the corresponding programming methodology. This is even more true, because a comprehensive reference manual and application guide is currently being prepared.

Perhaps with the exception of the alarm clock solution in Section 5, the decision for the ordering of the examples was based on increasing complexity. Therefore, a CSSA greenhorn is advised to profit from the sequence given. For the advanced reader a deviating reading order or picking up one specific example should cause no problems.

In addition to short comments on the well known synchronization problems discussed in Sections 2 to 5, "equivalent" solutions for the same problems using monitors are presented. For sake of coherent and unmisleading reference all these monitor solutions are programmed using syntax from the programming language CLU [Lis, Sny, Atk, Sch 77] and are taken from [Bloom 79]. Without intending a rigorous comparison between the monitor approach and the CSSA approach, some hints at specific differences are given.

Due to personal time constraints I did not succeed in making this paper self contained. I rather decided to make the following assumptions:

1. The reader should have some basic knowledge about the CSSA concepts and programming constructs.  
If this is not the case, he should first read [FRV 81] or the more concise and complete publication [FRV 82].
2. A full understanding of Sections 2.3, 3.3, 4.1.3 and 5.3 is only possible, if the reader is familiar with the monitor concept. No further comments at the syntax and semantics of the monitor programs are included in this paper.

All the CSSA examples have been compiled and executed by the BMS-CSSA-Simulation System, which was designed and implemented by the students C. Beilken, F. Mattern and M. Spenke (BMS). Without their great efforts and personal engagement programming in CSSA would have remained a paper and pencil exercise for an indefinite amount of time.

In the simulation system, translation from a CSSA source script into executable code is a two step process. First, the CSSA source is compiled into a SIMULA program, which in the second step is embedded in a SIMULA multiprocessor simulation system.

All but one example have been executed with the default processor configuration depicted at the beginning of the first session protocol (fig. 2). For the window mechanism example described in Section 7, a problem oriented configuration was designed by simply changing the default parameters of the simulation system.

In all comments, lines of the source programs are referred to by

enclosing their numbers in parantheses.

## **1. One Slot Buffer**

### **1.1. Problem Specification**

A one slot buffer is an object which may contain exactly one information element. A new information can be put into the buffer only via a write operation and later be inspected by a read operation.

The synchronization problem is specified by the following restrictions:

1. write and read operations must take place in alternating order.
2. The first operation to be executed must be a write operation.

### **1.2 CSSA Solution**

A CSSA solution for the one slot buffer problem is straightforward (fig.1). The buffer is represented by a string variable BUFFER. In two parallel facets READ and WRITE the only operations RD and WR are defined respectively. Each execution of an RD operation causes a change to facet WRITE, such that a WR operation must be executed next. The same is true vice versa. The initial statement (25) guarantees the first operation to be a WR operation.

### **1.3 General Remarks**

The history information of alternating reads and writes can be very well expressed by use of the facetting mechanism. The CSSA solution appears to be simpler than any monitor solution, since there at least one extra (boolean) variable storing this history information is needed. However, monitors do not have facets. But in my opinion, besides their practical importance the facet structure contributes much to a clear documentation of the CSSA solution.

## 2. Bounded Buffer

### 2.1 Problem Specification

A bounded buffer can be regarded as a generalization of the one slot buffer. In contrast to the one slot buffer it can store more than one element up to a fixed maximum number.

Users of the buffer can store and inspect information via two operations called **insert** and **remove**.

Like a one slot buffer agent, a bounded buffer agent has to solve the following synchronization problems:

1. New information may only be inserted into a buffer element iff
  - 1.1 either this element never has been used for insertion before, or
  - 1.2 the last information having been inserted into this element has already been removed.
2. A buffer element may only be removed iff information has been written into this element before, and after this insertion no other remove operation for this element has occurred.

So, each buffer element can be interpreted as a one slot buffer. A bounded buffer with maximum element number 1 should exhibit the same behaviour as a one slot buffer.

### 2.2 CSSA Solution

In each creation message for a bounded buffer agent the buffer size MAX (line 1) must be included (cf. fig. 3). According to the problem specification, at each time the buffer agent is in one of three possible states:

1. In facet **BUFFER\_EMPTY**, there is no buffer element which can be removed (operation **REM**).  
This is the case in the initial state (40) or later, when all inserted elements have been removed and after these removals no new elements have been inserted. Therefore, the only messages which can be processed in this facet are insert messages (operation **INS**).
2. When exactly MAX insertion messages have been processed without an intervening remove operation, the buffer agent's current facet becomes **BUFFER\_FULL**.  
In this state the enabling of an **INS** message would destroy the information in a buffer element, which has not yet been inspected by a **REM** operation. Since this would violate restriction 1.2 of the problem specification, only **REM**



operations are allowed in this state (18).

3. In facet NO\_CONSTRAINT there exists both

- 3.1 a buffer element into which information has been inserted but not removed, and
- 3.2 a buffer element which either never has been used for insertion or which already has been inspected and after inspection no further insertion has taken place.

Because of 3.1 at least one REM message can be enabled, whereas 3.2 allows the processing of at least one INS message.

In order to decide which of these facets is to be chosen next, two integer variables INS\_COUNT and REM\_COUNT are incremented modulo MAX in each INS and REM operation respectively (27,34).

At each time, INS\_COUNT is the index of the next buffer element to be used for insertion, and REM\_COUNT yields the index for the next REM operation. If REM\_COUNT becomes equal to INS\_COUNT, the next current facet must either be BUFFER\_FULL or BUFFER\_EMPTY, depending on what operation was executed last (28,35).

A general comment on the program structure will close this section. Besides the natural disjunction of facets the bounded buffer solution demonstrates how the programmer profits from the possibility to define operations at the top level of a script (3,25-37). The advantage is that he has to write down the operations only once, including them in a facet where needed (14,18,22). In this specific example, the programmer has to put up with a slight disadvantage: as he can't decide in which facet the operations will be executed at any time, each enabled message causes a new facet instantiation to be set up. I.e., a new instantiation of facet NO\_CONSTRAINT will be set up, although possibly the agent could have stayed in this very facet. Certainly, this little inefficiency could have been avoided by introducing a further variable remembering the current facet. For facets BUFFER\_EMPTY and BUFFER\_FULL, there is no such problem, because each operation execution in one of these facets must cause a change to another facet.

### 2.3 Comparison with the Monitor Solution

In script BOUNDED\_BUFFER no event conditions (buffer.nonfull, buffer.nonempty, c.f. fig. 5) are needed. Instead of starting the execution of an arbitrary operation, and then testing if the started operation may continue, always only allowed operations are selected and executed without interruption.

This strategy rules out a more general difference. As a monitor

programmer has no influence on the sequence of started operations, often one of the first actions of a monitor procedure is to check if its execution may continue. In contrast, CSSA operations are indivisible. Once started, they must be executed without interruption. Hence, a CSSA programmer is forced to extract the synchronization code from the "operations" and put it into the facet structure and possibly into assertions governing the selection of messages in the mailbox.

### **3. Protected Buffer**

#### **3.1 Problem Specification**

In the bounded buffer solution of section 2. the same agent implemented both the **buffer** and the **synchronization code** for the buffer. In other words, the synchronization code visualized in the special facet structure and pure buffer access functions have been mixed up.

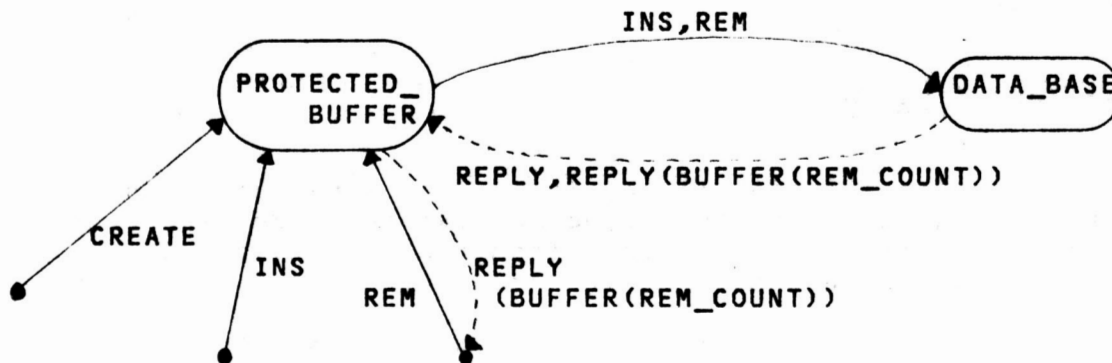
As is discussed in more detail in [Bloom 79], modifiability, modularity and understandability will be enhanced when the buffer resource object itself is no longer part of the module which implements its synchronization scheme. Defined in such separated modules, the implementation of the buffer may be changed without modifying its synchronization code. Conversely, the synchronization scheme for the buffer can be altered without changing the buffer implementation.

#### **3.2 CSSA Solution**

In fig. 6, the synchronization code is defined by script **PROTECTED\_BUFFER**, whereas the definition of the buffer with the buffer access functions is given in script **DATA\_BASE** (6-35).

An agent of type **DATA\_BASE** exactly executes the operations he is ordered to by its protector agent (an agent of type **PROTECTED\_BUFFER**). Due to the inherent assumption in the CSSA model of computation that messages may pass each other, the protector-resource relationship must be "strongly coupled", guaranteeing that protector orders are followed in the sequence given. This can be achieved by having each execution of an **INS** or **REM** operation by the protector agent only finished after receiving a reply (78,91) indicating that the corresponding order to the resource agent (76,89) has been followed.

Script **DATA\_BASE** does not seem to fit very well to the requirements **R\_COUNT** are never needed in this environment. In spite of this inconvenience, the database is included in its present structure, because the same script can be used for the readers\_writers solutions in the following section.



Users of the buffer communicate only with an agent of type `PROTECTED_BUFFER` who delegates requests to an agent of type `DATA_BASE`. It's impossible for a user to gain direct access to the real resource agent.

### 3.3 Comparison with the Monitor Solution

The two solutions differ with respect to the interpretation of where the borderline between the synchronization code and the buffer code should be defined. In the CSSA solution, a buffer is simply a data structure which can store and retrieve information provided an unique key is presented. In the monitor solution (fig. 8), a buffer still has much similarity with a bounded buffer in that it provides two functions `full` and `empty`. As a consequence, the protected buffer monitor no more has to take care of the counters `REM_COUNT` and `INS_COUNT`. In my opinion, these counters are essential elements of the synchronization code and not of the buffer definition.

Not willing to go into the particulars of these arguments, I only state that with the monitor concept and the CSSA concept corresponding solutions can be given for both interpretations. Certainly, no lack of expressiveness in either of the two concepts could be derived from such a discussion.

## 4. Readers Writers Problems

The readers\_writers problem is particularly interesting because with only slight modifications of the general problem (cf. section 4.1.1) various versions with different synchronization aspects can be formulated. In section 4.1 the readers\_writers problem with readers priority is discussed in more detail. Section 4.2 gives hints at solutions of other versions. Two results are especially noteworthy: whereas the monitor concept cannot succeed in giving a solution for the general readers\_writers problem, no satisfying solution for the readers\_writers problem with pure `first_in_first_out` scheduling can be given in CSSA.

## 4.1 Readers Writers with Readers Priority

### 4.1.1 Problem Specification

Given a database with read and write access functions, the **general readers\_writers problem** can be specified by the following constraints:

- (1) write operations are mutually excluded with other write and read operations.
- (2) read operations may execute in parallel.

Special versions of the general readers\_writers problem are defined by adding one or more further restrictions. By stating

- (3) read requests have priority over write requests

we get a specification for the readers\_writers problem with readers priority.

### 4.1.2 CSSA Solution

In fig. 9, script DATA\_BASE (3-22) implements the read and write access functions. The use of the third operation RD\_ACKN (23-26) will become clear soon. Note, that this is the same database as was used for the protected buffer solution. In the same sense as described in 3.1, the protection module RDER\_PRI0 defines the synchronization code specific to the readers priority restriction. An instantiation of one particular database can only be achieved by sending a CREATE message including the requested size of the database to an agent of type RDER\_PRI0 (44-47). Read and write requests can only be addressed to this protector agent because no acquaintance to the database can be communicated to the outside world.

The functioning of the whole system is best described by showing that restrictions (1)-(3) from 4.1 are fulfilled.

To this aim, I first want to interpret the rather intuitive notions in these assertions according to the present concrete environment.

- A read or write request occurs when a RD or WR message arrives in the mailbox of the protector agent.
- Read operations  $RD_1, RD_2, \dots, RD_j$  execute in parallel iff their requests have been accepted and executed by the protector agent, but no "corresponding" acknowledgement has been received at port R\_ACKN in the protector's write operation WR.

Hence, from the protector's point of view, these read operations are either

- just being transmitted between the protector and the database,
- in the mailbox of the database.
- just being executed in the database,
- already executed in the database.

Before receiving an acknowledgement at port R\_ACKN, the protector has no information about the state of the database. In this case, we also say that  $RD_1, RD_2, \dots, RD_i$  are in flight (between the protector and the database).

Hence, this definition of parallel reading only allows for read requests being nondeterministically executed in the data base. As the execution of operations is indivisible and sequential, the database agent can only accept and answer read requests one after another. The execution is nondeterministic, because no one can guarantee that the read operations are accepted by the database in the same order they were received in the protector's mailbox or sent away by the user agents.

The notion of "mutual exclusion" is explained more precisely as:

- **Mutual exclusion** between read and write requests and between write operations themselves is guaranteed iff no other operation ever can be in flight together with a write operation.

**Proof of assertion (1) [mutual exclusion]:**

- (1) It's easily shown that one write operation excludes other write operations because it must be acknowledged [ wait W\_ACKN do ..., (80) ] before the next write request may be accepted. Receiving this acknowledgement guarantees that the database has accepted the write operation. So, a next write operation never can't be in flight together with a preceding write operation.
- (2) In order to show that "write excludes reads" we have to prove that a write operation never can't be in flight together with read operations. According to the acknowledgement protocol described in (1), only read operations have to be taken into account.

The proof is by induction:

- (2.1) When a protector agent is created, it first must accept a CREATE message [ initial CREATION, 90 ] resulting in a change to facet WRRS (46). Because the idle operation has no effect when enabled before the first write operation [ FIRST\_OPER = true in line 85 ] the protector stays in facet WRRS until a write request will be accepted. So, the first write operation clearly is not in flight together with read operations.

(2.2) Assume the  $k$ 'th write operation was not in flight together with other read operations.

As an effect of its execution the protector's `RDER_COUNT` was reset to 0 (78). Having received an acknowledgement at port `R_ACKN` (77), the protector is assured that in the database `RCOUNT` has been reset to 0 (25).

Then the completion of the protector's `WR` operation results in a change to facet `RDERS` (85), since `FIRST_WR` and `FIRST_OPER` have been set to false (72,73).

Now, two cases have to be considered:

(2.2.1) No read operation is requested:

Then the idle operation in facet `RDERS` immediately leads back to facet `WRRS` (59).

(2.2.1.1) If no write operation is requested, the protector will loop back to facet `RDERS`. These facet changes will continue until either a write operation (2.2.1.2) or a read operation (2.2.2) will be requested.

(2.2.1.2) If a write operation is requested it will execute without any read operations being in flight.

(2.2.2) A read operation  $RD_i$  is requested.

$RD_i$  will be accepted and `RDER_COUNT` will be set to 1 (56). As long as further read operations  $RD_i$  ( $i > 1$ ) are requested before the protector has finished execution of read operation  $RD_{i-1}$ , they will be accepted and sent to the database.

(2.2.2.1)  $\forall i \in \mathbb{N}$ .

( $i > 0 \Rightarrow$  read operation  $RD_i$  has been requested before  $RD_{i-1}$  has finished execution)

In this case, an infinite number of read operations will be transmitted to the database. Therefore, no further write operation will ever be executed and "the  $(k+1)$ -st write operation is not in flight with read operations" holds trivially.

(Being polite assume here, that `CSSA` integer variables can store any natural number. So our program never will abort with integer overflow.)

(2.2.2.2) Assume  $RD_j$  finishes before  $RD_{j+1}$  has been requested. Then the idle operation will be executed causing a change to facet `WRRS`.

(2.2.2.2.1) A write operation is not requested.

Then the idle operation leads back to facet `RDERS` with `RDER_COUNT` =  $j$ .

When new read requests have arrived they will be executed either ad infinitum (case 2.2.2.1) or until  $RD_{j+k}$  finishes before  $RD_{j+k+1}$  is requested. In the latter case we have the same situation as in 2.2.2.2 with `RDER_COUNT` =  $j +$

k.

(2.2.2.2.2) Terminating the induction, we assume that a write operation (the (k+1)-st write operation) is requested.

RDER\_COUNT > 0 is the number of read operations being in flight. When the message RD\_ACKN(RDER\_COUNT) is sent to the database (76), it will only be accepted when RDER\_COUNT read operations have been executed in the database (assert COUNT = RCOUNT, 23). Therefore, the receipt of an acknowledgement at port R\_ACKN guarantees that all read operations have been executed. When the write operation is sent to the database (86), no read operation is in flight. ■

In this proof we already described a situation resulting in more than one read operation being in flight. So we have assertion (2) [parallel reading] automatically.

Assertion (3) [priority constraint] is easily shown:

Without loss of generality, assume that both a read request and a write request exists and at least one operation has been executed before. I have to prove that whatever operation read or write was executed last, the next request accepted by the protector will be the read request. Now consider the two cases:

- (1) The last operation executed was a read operation:  
Then we stay in facet RDER\_S, and obviously the read request will be satisfied next.
- (3) The last operation executed was a write operation:  
Because FIRST\_WR in facet WRRS is false, no further write request is acceptable. Therefore, the protector will change to facet RDER\_S and execute the read operation.

#### 4.1.3 Comparison with the Monitor Solution

It's rather difficult to compare two solutions when one cannot observe much similarities. The most significant difference stems from the fact that in contrast to the monitor solution (fig.11a and 11b) I don't need a queue storing user's write and read requests in the CSSA solution.

In the CSSA environment, a request was defined to occur when the corresponding message is received in the protector's mailbox. In the monitor concept, a user request can only come into existence from the programmer's point of view, when the execution of the called monitor procedure is initiated. Translated into the monitor terminology, this is the time stamp when the request enters the monitor. However, from the monitor implementation's point of view, user requests already exist before the time stamps of their



procedure entry points. On the implementation level, incoming requests are registered and stored into an internal queue, which the programmer has no access to. In some sense, this initial queue and not the programmer-defined queues must be regarded as the monitor data structure corresponding to the CSSA mailbox. Hypothetically, a monitor solution similar to the CSSA solution could be derived if the following condition were met: The initial monitor queue is not a queue, rather it is a data structure with almost arbitrary programmable access functions. In CSSA we can select nearly arbitrary messages by defining suitable facet structures and assertions.

## **4.2 Other Versions of the Readers Writers Problem**

Constraints (1) [mutual exclusion] and (2) [parallel reads] formulated in 4.1.1 are common to all versions of the readers\_writers problem. Therefore, only a third constraint will be stated to give specifications for the following specializations of the general problem.

### **4.2.1 The Fair Readers Writers Problem**

The readers\_writers problem with readers\_priority allowed writers to starve (cf. case 2.2.2.1 in the proof of assertion (1)). Substituting assertion (3) [readers\_priority] by

(3a) eventually every request will be served

a solution is required to be fair against readers and writers. That is what I call the **Fair\_Readers\_Writers Problem**. Note, that requirement (3a) makes no specific assumptions about the order of serving requests. Thus, the quality of a solution should be measured by the freedom that it leaves for selecting the next request to be served. For example, a pure first\_in\_first\_out scheduling of requests fulfills the requirements, but is too restrictive to be regarded as a good solution.

A really satisfying CSSA solution can be derived from the RDER\_PRIO script by simply placing the operations RD and WR into only one facet. The priority ensuring idle operations and the assertion FIRST\_WR for the write operation are no more needed. With operations RD and WR occurring in one facet the selection of the next request is totally left to the mailbox manager. From the programmer's point of view this message selection strategy can only assumed to be fair and nothing more. Hence, the proposed CSSA solution fits the problem specification in the best possible way. Even more, this solution is perhaps the most natural and certainly the simplest among all solutions for other versions of the readers\_writers problem.

It should be not surprising that - if at all - only a rather com-

plex and not efficient solution for the same problem can be programmed with monitors. Here, the only data structure for storing requests is a queue. Each signal statement can only release the first request of a queue. Hence, a monitor programmer can never reach the state of freedom of a CSSA programmer who can rely on the more general message selection strategy of the mailbox manager.

#### 4.2.2 The FIFO Readers Writers Problem

Another version of the readers\_writers problem is given by

(3b) each request has priority over all later requests.

Here, priority is entirely based on order of requests. Obviously, this version postulates a pure FIFO policy for serving read and write requests.

According to the definition of a "request" in 4.1.2, at each time the earliest message received has to be selected from the mailbox. However, in CSSA there is no mechanism which allows selection to be based on receive order. There is no hope to give a correct CSSA solution in the context of the definitions established in section 4.1.1. This context is important, because it is not generally impossible to program a FIFO solution in CSSA. So, I could require that each request is transmitted with a unique natural number indicating the time stamp when it had been sent away. Then the protector only accepts that message with the time stamp of the previously selected message incremented by one. Certainly, users would be obliged to coordinate in keeping the transmitted tags unique and successive. To this aim, before sending a request to the protector they could be forced to request a time stamp from a globally known "clock agent".

A monitor solution for the FIFO\_READERS\_WRITERS problem is presented in fig. 12. Although the monitor's queueing principle fits very well to the required FIFO strategy, the solution seems to be rather complex. This complexity is essentially due to the fact that at least the conventional monitor construct provides no means of identifying the process at the head of a queue or determining the conditions for which it is waiting. In fact, read and write requests are stored into a single queue (m.users) when they are not immediately satisfiable. But if a writer is dequeued when readers are in flight (readercount > 0), the writer must be queued again on a second queue until it is released by the last reader in the resource. A more profound discussion of these problems is given in [Bloom 79].

#### 4.2.3 Readers Writers with Writers Priority

Substituting assertion (3) of the readers\_writers problem with readers\_priority by

(3c) write requests have priority over read requests

yields a specification for the readers\_writers problem with writers\_priority. The construction of a CSSA solution analogue to script RDER\_PRI0 is left to the reader as a simple exercise.

## 5. The Alarmclock

### 5.1 Problem Specification

The alarmclock is a very nice representative of the class of problems where synchronization is based on the arguments of requests.

It can be considered as an abstraction of a system device allowing users to fall asleep for some specified time period and to be awakened after this time has elapsed. The alarmclock cannot increment the time by itself, it rather waits for a tick request to be received from some other system facility.

### 5.2 CSSA Solution

In fig. 13, a CSSA solution is presented which looks very elegant but unfortunately is wrong. Notwithstanding, I enclosed this bad solution because it's a good demonstration of the dangers waiting for a programmer in the distributed environment. So I admit, that after having programmed this solution I accepted it as correct until the last moment before publishing this report.

Although perhaps not visible in their surface structures, the ideas behind the wrong solution and the correct one (fig. 14) are much the same. Therefore, I can afford some documenting of the wrong solution without wasting time.

An agent wishing to sleep for  $T$  time units sends a message `WAKE_ME(T)` to an agent of type `ALARM_CLOCK` [`ALARM_CLOCK` for short]. When the `WAKE_ME` message is accepted by `ALARM_CLOCK` the sleepy agent is immediately resumed if  $\text{DELAY\_TIME} (=T) \leq 0$  (23). For a  $\text{DELAY\_TIME} > 0$ , the wake request is registered by sending a private message `WAKE_UP(DELAY_TIME + NOW)` to self (24). The parameter  $\text{DELAY\_TIME} + \text{NOW}$  yields the "absolute" time stamp when the sleeping agent has to be awakened. Later the existence of this `WAKE_UP` message in its mailbox will remind `ALARM_CLOCK` of the user's reveille. Because ringing is symbolized by sending a reply, the port address implicitly included in the `WAKE_ME` message must be transferred to the `WAKE_UP` message via the `inherit` - clause. Now we arrive at the crucial point of the discussion. Each execution of a `TICK` request increments the "real time" `NOW` by 1 (28) and sets up a new instance of facet `RESUME`. The idea is, that in this facet all `WAKE_UP` requests for the current time ( `assert TIME = _NOW, 11`) should be satisfied. When no such further `WAKE_UP` request exists, the execution of the idle operation (15,16) leads back to facet `REGISTRY`.

At first sight, this strategy seems very sound. But did you get on to the problem? So assume for example, that a `WAKE_ME` message with  $\text{DELAY\_TIME} = 1$  has been received and in response `ALARM_CLOCK` has sent a `WAKE_UP` message with parameter  $\text{NOW} + 1$  to itself. After the next execution of a `TICK` operation, this `WAKE_UP` message must be served in facet `RESUME`. Now, we are faced to the problem that

ALARM\_CLOCK stays in facet RESUME only until no (further) WAKE\_UP message with TIME = NOW exists in the mailbox. When the considered WAKE\_UP request has already been received in the mailbox, it will be executed. However, if it has not been delivered for the time being, the current facet will be left via the idle operation and the sleeping agent will sleep forever.

Although from a practical point of view, the probability may be very low that a WAKE\_UP message will not have been received when it is needed, this program must be rejected as a wrong solution. The computational model of CSSA only requires message passing times to be finite, but indefinite. That's even true in this case when messages are passed to the sending agent itself.

Obviously, the idle operation is not the right criterion for leaving facet RESUME. Instead it must be assured that all WAKE\_UP requests for the current time value NOW are served before this facet is left. In the "hopefully" correct solution depicted in fig. 14, this is achieved by keeping a relation CALL\_LIST recording the number of agents to be awakened (NO\_OF\_SLEEPERS) for each time stamp which ever has been requested.

I.e., a request WAKE\_ME(DELAY\_TIME) is served as follows: when another wakeup obligation for time DELAY\_TIME + NOW exists, the corresponding call element (of type CALL\_ELEM) is assigned to the implicitly declared variable CE1, (33) and its NO\_OF\_SLEEPERS' part is incremented by 1 (34). Otherwise, a new call element with "key" WAKE\_UP\_TIME = DELAY\_TIME + NOW and NO\_OF\_SLEEPERS = 1 is inserted into the relation (36-38).

In both cases, ALARM\_CLOCK sends a WAKE\_UP message to itself (40). These private messages are still necessary since for reasons not to be discussed here, the inherit clause is the only mechanism available for saving a reply address from one operation execution to another. In this example, however, the use of private messages could be avoided, when ringing would be implemented by "normal" wakeup messages instead of reply messages. Then, in addition to the DELAY\_TIME parameter, a sleepy agent must include an acquaintance to itself in the WAKE\_ME message, which later will be used by ALARM\_CLOCK as the addressee of the corresponding WAKE\_UP message.

After incrementing the time (45), a new instantiation of facet RESUME is only set up when a call element for the new time value exists (46). Compared with the first solution this means a slight compensation for the overhead we had to introduce in the second solution. ALARM\_CLOCK then stays in facet RESUME, waiting for the first WAKE\_UP request which eventually will arrive in the mailbox, if it was not already there. So, there is no need for an idle operation in this case. Another reason is that the criterion for leaving facet RESUME can be computed during the execution of a WAKE\_UP operation: If NO\_OF\_SLEEPERS for the currently served call element has the value 1 (19), then no further wakeup obligation exists for the current time value and facet RESUME may be left. Besides, the current call element can be deleted (20). Otherwise, NO\_OF\_SLEEPERS is decremented by 1 (22) and no change of facet takes place.

### 5.3 Comparison with the Monitor Solution

In monitors, queuing is the only means to stop serving a request which later shall be resumed. Because in this example delayed requests have to be resumed in the order of their arguments, the concept of **priority queues** has to be introduced (cf. fig. 16). Most of the alarmclock code in the CSSA solution is needed to implement this data structure which is a necessary prerequisite of all monitor implementation languages. From this point of view, the seemingly more complex CSSA solution doesn't make me so unhappy, rather it indicates the provision of greater freedom and flexibility in the CSSA concepts.

An unpleasantness of this special monitor solution is due to the fact that it is not possible to examine the first element of a queue without dequeuing it first. Thus, the first process in the priority queue has to be dequeued in each tick operation, then tested for its arguments ( `while ac.now < alarmsetting` ), and possibly requeued again. As already mentioned in 5.2, this awkwardness can be avoided in the CSSA solution by not establishing an instance of facet RESUME when no wakeup obligation exists for the current time. Clearly, this shortcoming of the monitor solution is not so severe because adding a test for the priority of the first element in a priority queue as an additional language feature is all what has to be done.

## 6. The N-Queens Problem

### 6.1 Problem Specification

The following problem has to be solved:

How can you place  $N$  queens on an  $N \times N$  chessboard in such a way that no two queens can capture each other (i.e. no two queens are in the same row, column or diagonal)?

The CSSA solution described in the next section is expected to find all such possible placings for any given  $N$ .

In general, the  $N$ -queens problem is a showboy for the class of problems offering "natural" recursive solutions. In some sense, the CSSA script to be developed also defines recursive behaviour. Instead of recursive function calls, agents of a certain type are created dynamically and provided with partial solutions which they are ordered to expand.

### 6.2 CSSA Solution

The script NQUEENS in fig. 17 is best described by setting an example for one specific  $N$ , say  $N=4$ .

After having created an agent NQ of type NQUEENS, a user specifies the command

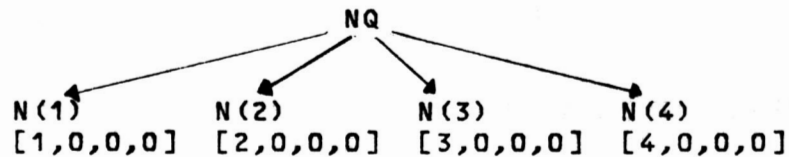
```
send FIND_SOLUTION(4) to NQ
```

Because  $N$  does not equal 1 (which would lead to the immediate answer: "For one queen solution is : 1" [75-79]), NQ sends a message FIND\_SOLUTION to each of four newly created anonymous agents of type NODE [84]. We want to identify these new agents as  $N(1)$ ,  $N(2)$ ,  $N(3)$  and  $N(4)$ . Each agent  $N(i)$ ,  $i=1, \dots, 4$ , is provided with the information that a queen has been placed by NQ in column  $i$  of the first row. This information is represented in the CSSA script by putting the value  $i$  into the first element of the array FIRST\_ROW [83]. The other array elements retain the value 0 due to default initialization [71].

In the following course of execution nothing remains to be done for NQ. One can terminate it by sending a STOPQ message [89-93], or let it survive when solutions for other sizes of the chessboard shall be generated.

The current situation is skeletonized in the following tree structure:





For example,  $N(3)$  is ordered to find a solution for the 4-queens problem with already one queen placed in column 3 of the first row.

The first action of each agent  $N(i)$  is to check whether it already was provided with a complete solution. This would be the case if a queen had been placed in some column of the last ( $N$ -th) row [ if **not**  $(Q(N) = 0), 19]$ .

The actions to be taken in this case are described later in this section when the condition will be fulfilled for our example execution. Now, for all agents  $N(i)$  this condition is not fulfilled. So, they try to find allowable places for the next queen to be placed in the current (second) row. This computation is done in three separated loops.

The first loop [34-36] determines how many queens ( $ANZ\_Q$ ) are already on the board. In our example,  $ANZ\_Q$  will have the value 1 after leaving the loop.

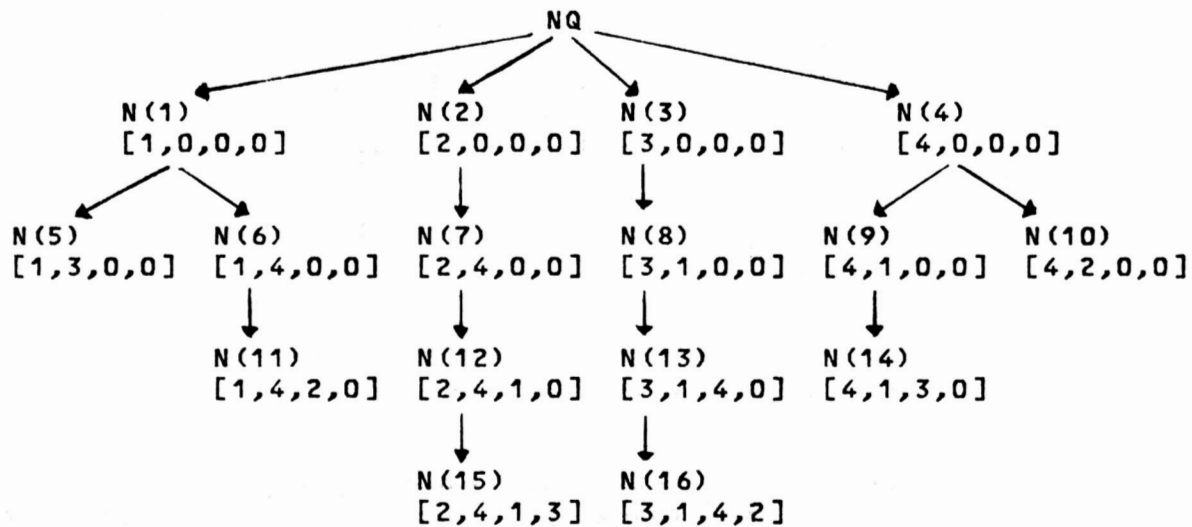
In the second loop [38-46], all columns of the current row are computed, which are **not** allowed for the next queen. To remember this information, the  $j$ -th element ( $j \in \{1, \dots, 4\}$ ) of the boolean array  $NOT\_POSSIBLE$  will become **true** iff no queen may be placed into the  $j$ -th column of the current row.

In the third loop [48-53], a new agent of type  $NODE$  will be created, one for each allowed extension of the board. If no extension is possible (i.e.  $NOT\_POSSIBLE(j)$  for  $j=1, \dots, 4$ ), the loop has no effect.

The last action of each  $NODE$  agent is to request its own termination (56). Hence, all  $NODE$  agents are alive only as long as necessary. After their work has been done they release their processors for the execution of other  $NODE$  agents which they possibly have created during their own lifetimes.

The coming into existence of new agents in our example execution is represented in the following tree structure which is an expansion of the previous one.





During the construction of the solution 16 NODE agents have been created. However, due to the self destruction of NODE agents this number should not be misunderstood as the number of agents executing simultaneously on the available processors. Agents N(5), N(11), N(14) and N(10) terminate without finding an extension of the given partial solution. For example, N(5) is ordered to place a third queen into the third row with respect to the partial solution depicted by the following board.

1	*			
2			*	
3				
4				
	1	2	3	4

Obviously, this board cannot be expanded, and no new agent will be created. Only agents N(15) and N(16) are provided with complete solutions:

1		*		
2				*
3	*			
4			*	
	1	2	3	4

		*	
*			
			*
	*		
1	2	3	4

(boards for [2,4,1,3] and [3,1,4,2])

For their executions the previously mentioned condition **not** ( $Q(N) = 0$ ) is fulfilled. Therefore, their only task is to send a message of success to the user. The envelope of this message is constructed using the string variable ASW and sent to the user ( interface ) by

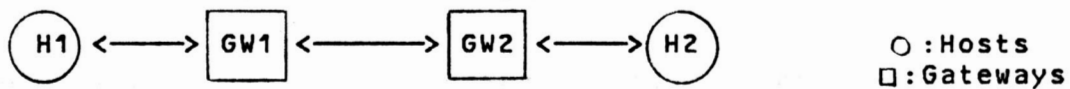
**send ANSWER(ASW) to interface reply to OK;**

The reply obligation is not really necessary. Its only purpose is to enhance the attention of the human user observing or not observing messages being displayed on the screen.

## 7. Window Mechanism

### 7.1 The Model

In this example CSSA is used to model the following computer system:



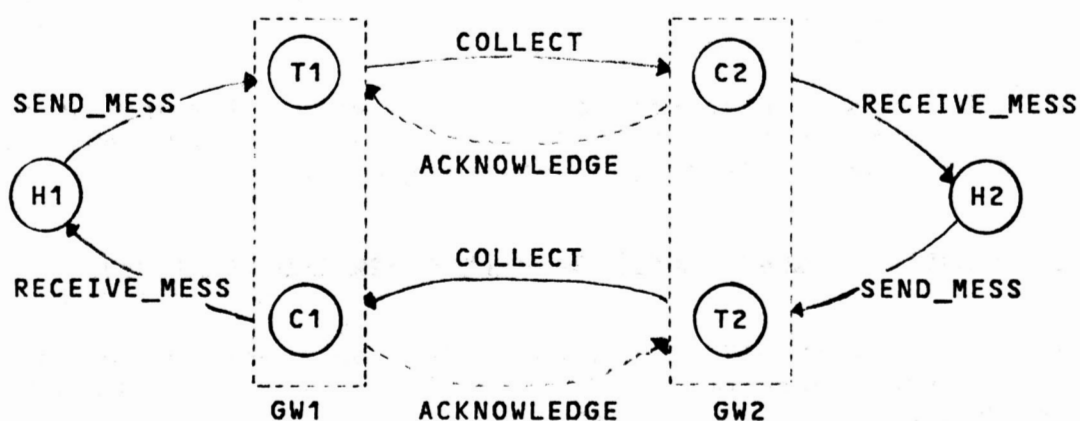
Two host computers communicate via two interface modules which are called gateways.  $H_i$ ,  $i \in \{1, 2\}$ , can send (receive) a complete message to (from)  $GW_i$  whereas only packets are transmitted and received between  $GW1$  and  $GW2$ .

In the model, messages are character strings of variable lengths and the individual characters of the strings are packets.

Each gateway provides two facilities:

- a transmitter (T) as one part of the gateway receives a message  $m$  from "its" host and sends the characters of  $m$  to the collector in the other gateway.
- a collector (C) as the other part receives packets from the transmitter in the other gateway. When the packets of one message are all received the concatenation of the packets is sent to the host as one message.

We assume that several packets may be simultaneously in flight between T and C. This extended model is described by the following picture:



In the CSSA solution all circled entities will become agents which communicate via the operations labelling the arcs. So H1 can send a message via T1 and C2 to H2 at the same time when H2 is sending a message via T2 and C1 to H1. Therefore the most effective realization of this CSSA agent structure would supply one (micro-)processor for each of the Ti's and the Ci's. Now we introduce the following realistic complications:

- packets need not reach the collector in the order they were produced by the transmitter; that means packets can pass each other;
- the transmission medium between the two gateways is unreliable so that packets can be totally lost.  
(It should be noted that this second assumption is in contrast to the CSSA model of computation which does not allow the loss of messages. In the CSSA solution we therefore simulated this unreliable transmission medium by generating two agents of type FAULTY-CHANNEL (cf. fig. 19, 353-442), which receive and normally forward all messages transmitted between a collector and a transmitter. However, depending on generated random numbers, some messages are received but not transmitted, that means they are lost).

In this situation the communication between a transmitter and a collector becomes most interesting. We are faced with the problem of providing the hosts H1 and H2 with a reliable "virtual" transmission medium which is based on an unreliable "physical" transmission medium.

In a simple solution of this problem the transmitter T would send a packet p1 to the collector C, then wait for an acknowledgement, then send a second packet p2, etc. Because the transmission medium is unreliable not alone for packets but also for acknowledgements, some timeout mechanism must be used when waiting for an acknowledgement. When time is exhausted the packet not yet acknowledged must be sent once more. The implementation of receiving duplicates is straightforward in this case.

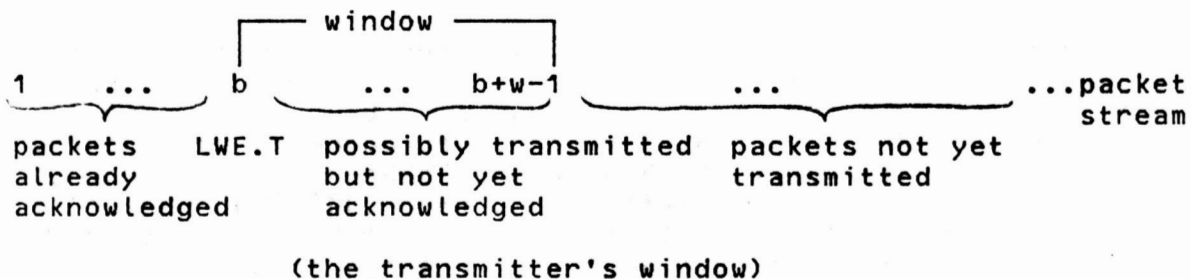
However, this solution is not really satisfying, because it does not utilize the whole capacity of the transmission medium (more than one packet may be in flight between T and C).

## 7.2 Communication between the Transmitter and the Collector

The CSSA scripts TRANSMITTER and COLLECTOR implement a mechanism which is described informally in [Cerf, Kahn 74]. This so called **window mechanism** works as follows:

The transmitter sends packets (in operation SEND\_PACKET, 71-88) to the collector. Received packets (operation COLLECT, 156-247) are acknowledged by the collector. Instead of waiting for an acknowledgement (operation ACKNOWLEDGE, 91-116) before the next packet can be transmitted, the transmitter can hurry on ahead; so several

packets can be in flight without having been acknowledged. Obviously, every packet must carry a unique identification (sequence number) which allows the reconstruction of a message from the individual packets. In order to guarantee the perception of lost packets, the transmitter eventually must stop transmitting new packets. Therefore a common upper bound for the maximum number of packets in flight is agreed upon between the sender and the receiver. This upper bound  $w$  is the size of the so called window. In the window appears a part of the packet stream which starts just behind the last packet already transmitted and acknowledged, if any. The variable  $LWE.T$  is equal to the sequence number of the leftmost packet within the transmitter's window (left window edge).



If all packets of the current window have been transmitted a next packet can only be transmitted after having received an acknowledgement  $a$  with  $LWE.T < a < LWE.T + W$ . The receipt of  $a$  indicates that the collector has accepted all packets with sequence numbers less than  $a$  and now is waiting for the packet with number  $a$ . This allows  $LWE.T$  to be shifted to the right ( $LWE.T := a$ ), thereby implicitly advancing the right window edge.

However if a packet -or the acknowledgement itself- is lost, the transmitter never would get the chance of transmitting a new packet. Therefore some timeout mechanism must be used leading to the **assumption** that a packet or an acknowledgement has been lost. Based on this hypothesis the packet at the left window edge will be transmitted for a second time.

In our CSSA solution the retransmission of a packet is processed in the `idle_operation` (60-66) of facet `TRANSMISSION2` in script `TRANSMITTER`.

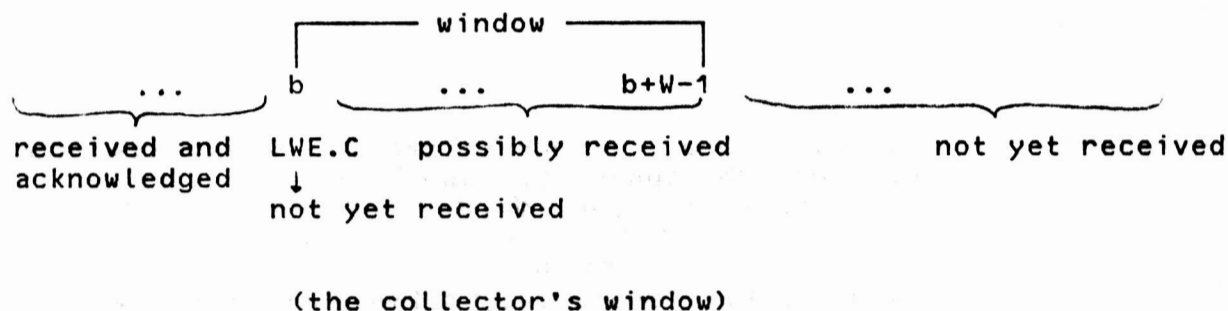
Here we use the idle operation as a simulation of a timeout mechanism, since no real time constructs are available in the current version of CSSA.

An immediately subsequent retransmission of the same packet due to the timeout mechanism is prohibited by setting the boolean variable `IDLE_ENABLED` to true and changing to facet `TRANSMISSION1`. So, a next activation of the idle operation becomes only effective after an useful acknowledgement advancing the left window edge has been received. In this implementation, every packet will be transmitted at most two times. It should be noted that this number is not the result of a principal decision. A maximal retransmission rate of 3, 4 or 100 times could have been programmed with the same ease.

In facet TRANSMISSION the transmitter works according to the following rules

1. The transmitter transmits packets with sequence numbers between  $LWE.T$  and up to  $LWE.T + W - 1$ .  
Each activation of the operation SEND\_PACKET transmits exactly one packet (81). The whole transmission mechanism is initiated by the operation SEND\_MESS (25-38) in facet START\_TRANSMISSION in which the transmitter sends a "sufficient" number of SEND\_PACKET-messages to itself (via procedure ACTIVATE\_TRANSMIT, 8-15).  
Later new SEND\_PACKET-messages are only produced in operation ACKNOWLEDGE (98,99). The last packet of the currently transmitted "host-message" is provided with the tag ENDE = true (77).
2. The receipt of an acknowledgement with number  $N$  indicates that the collector has received all packets with sequence numbers less than  $N$  and the packet with number  $N$  being the first one not yet received. Note that the acknowledgement  $N$  implicitly acknowledges all packets with sequence number less than  $N$ . If  $LWE.T < N \leq LWE.T + W$ , the left window edge  $LWE.T$  can be set to  $N$  (103). Assuming the rest of the message to be transmitted is long enough,  $N - LWE.T$  new SEND\_PACKET-messages can be produced (98); otherwise the transmitting of all remaining packets of the current message is initiated (99). If the idle operation was disabled before receipt of the acknowledgement, now it will be enabled (104-107).  
An acknowledgement with  $N \leq LWE.T$  is possible if it was overtaken by a later one with a higher number. Such an acknowledgement can simply be discarded. Only an information message will be sent to the interface in this case.  
An acknowledgement with  $N > LWE.T + W$  is not possible because at most packet number  $LWE.T + W - 1$  can have been transmitted and received before.  
An acknowledgement for a packet of an old message (MESS\_ID < CURRENT\_MESS\_ID, 92) has no effect at all.
3. On timeout (no ACKNOWLEDGE- or SEND\_PACKET - message can be enabled) the packet with sequence number  $LWE.T$  will be retransmitted if the current facet is TRANSMISSION2.

The left edge  $LWE.C$  of the collector's window (fig. 20) denotes the least packets number which has not yet been received.



Possibly received packets with sequence numbers greater than  $LWE.C$  up to the right window edge  $LWE.C + W - 1$  (window packets) are stored in the relation  $R$  under their (relative) "window position". Depending on the received packet number  $CHAR\_NO$  in operation COLLECT the following actions take place:

1.  $CHAR\_NO = LWE.C$  (179-242):

Let  $J$  be the unique number with  $0 < J < W$  such that all packets with sequence numbers less or equal to  $LWE.C + J$  have been received but the packet with number  $LWE.C + J + 1$  not yet being received. Then the currently received first part of the "host-message" can be prolonged by the substring defined by packet numbers  $LWE.C, LWE.C+1, \dots, LWE.C + J$  (194-205).

The window is adjusted by setting  $LWE.C$  to  $LWE.C + J + 1$  and the relation  $R$  is updated as an image of the new window (217-236).

As acknowledgement the number of the new left window edge is sent to the transmitter (207). This implicitly acknowledges packets with numbers less than the index of the new left window edge. When the whole "host-message" has been received it is transmitted to the destination host (184, 209-210).

2.  $LWE.C < CHAR\_NO \leq LWE.C + W - 1$ :

The currently received packet is not yet acknowledged but stored in the relation  $R$  (173-176). The window is not "moved".

3.  $CHAR\_NO < LWE.C$ :

The currently received packet is a duplicate the original of which already has been received. From the collector's point of view the receipt of a duplicate serves as a hint that the (implicit or explicit) acknowledgement for this packet might have been lost. Therefore the current left window edge  $LWE.C$  is sent as acknowledgement once again (243).

4.  $CHAR\_NO > LWE.C + W - 1$  is not possible.

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BMS - C S S A - C O M P I L E R

1981/12/01 11:11

DEFAULT-OPTIONS: NOTERM,NOTEST, CHECK, OBJECT, SOURCE,NOSTRUCT, XREF,RESWD=A,MAXP= 99,MAXE= 99,MAXD=  
 OPTIONS IN USE: NOTERM,NOTEST, CHECK, OBJECT, SOURCE, STRUCT, XREF,RESWD=A,MAXP= 99,MAXE= 99,MAXD=

	1	2	3	4	5	6	7	B L O C K N E S T I N G
1	type ONE_SLOT_BUFFER is script							+1
2								
3	var string: BUFFER;							
4								
5	facethead READ;							
6								
7	facet WRITE is							+2
8	public: WR;							
9								
10	operation WR(string: MESSAGE) is							+3
11	BUFFER := MESSAGE;							
12	replace by READ;							
13	endoperation							-3
14	endfacet							-2
15								+1
16	facet READ is							+4
17	public: RD;							
18								
19	operation RD is							+5
20	reply (BUFFER);							
21	replace by WRITE;							
22	endoperation							-5
23	endfacet							-4
24								+1
25	initial WRITE							
26								
27	endscript							-1

BMS-CSSA-COMPILER - DATE OF RELEASE: 30 SEP 1981 NO ERROR DETECTED  
 END OF COMPILING ON 1981/12/01 AT 11:16:06.00 RETURNCODE = 0  
 COMPILE-TIME (CPU) = 0.85 SEC. EXECUTION-TIME = 15.00 SEC.  
 NUMBER OF SOURCE-LINES READ = 27 NUMBER OF TOKENS = 62  
 NUMBER OF OBJECT-RECORDS GENERATED = 302

## \*\*\* CROSS-REFERENCE-TABLE \*\*\*

BUFFER	VARIABLE:STRING	3	11	20
MESSAGE	CONSTANT:STRING	10	11	
ONE_SLOT_BUF	USER_DEFINED_TYPE:SCRIPT	1		
RD	OPERATION:OPER	17	19	
READ	FACET	5	12	16
WR	OPERATION:OPER	8	10	
WRITE	FACET	7	21	25

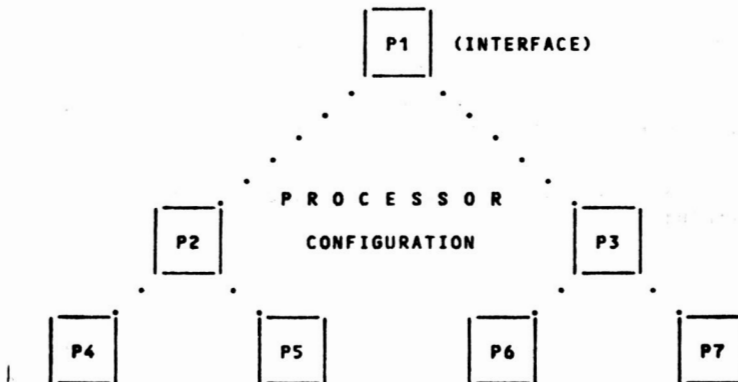
- MULTIPLE OCCURENCES ON THE SAME LINE ARE MARKED WITH '+'.  
 - TOTAL NUMBER OF IDENTIFIERS USED IN THIS PROGRAM: 7

fig. 1 : One Slot Buffer Script

=====

C S S A - S I M U L A T I O N - S Y S T E M

=====



=====

PROGRAM GENERATED ON 1981/12/01 AT 11:15:51.00  
BY BMS-CSSA-COMPILER (VERS. 30 SEP 1981)

PROTOCOL OF CSSA SESSION ON 1981/12/01 AT 11:48:14.00

=====

>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -  
==> DISPLAY

<==

IDENTIFIER	TYPE	VALUE
ONE_SLOT_BUFFER	SCRIPT	ONE_SLOT_BUFFER

```

>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -
==> VAR AGENT : OSB := NEW ONE_SLOT_BUFFER
>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -
==> OPER : RD,WR; PORT : P
>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -
==> OBSERVE; SEND RD TO OSB REPLY TO P
+++ 0.000 P1: INTERFACE(1) SENDS RD(),REPLY TO: P TO ONE_SLOT_BUFFER(1)
>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -
==> ;RUN
+++ 0.000 P2: ONE_SLOT_BUFFER(1) IS IDLE
+++ 0.100 P2: ONE_SLOT_BUFFER(1) RECEIVES RD(),REPLY TO: P
      FROM INTERFACE(1)
+++ 0.100 P2: ONE_SLOT_BUFFER(1) IS IDLE
>>> 0.100 SYSTEM TERMINATED
>>> 0.100 INTERFACE(1) : ENTER CSSA COMMAND -
==> SEND WR("A") TO OSB;RUN
+++ 0.100 P1: INTERFACE(1) SENDS WR(STRING:"A") TO ONE_SLOT_BUFFER(1)
+++ 0.200 P2: ONE_SLOT_BUFFER(1) RECEIVES WR(STRING:"A")
      FROM INTERFACE(1)
+++ 0.200 P2: ONE_SLOT_BUFFER(1) STARTING OPERATION WR(STRING:
      "A")
+++ 0.400 P2: ONE_SLOT_BUFFER(1) PERFORMS FACETTING : WRITE --> READ
+++ 0.400 P2: ONE_SLOT_BUFFER(1) STARTING OPERATION RD()
+++ 0.500 P2: ONE_SLOT_BUFFER(1) SENDS *REPLY(STRING:"A")
      TO INTERFACE(1)
+++ 0.600 P2: ONE_SLOT_BUFFER(1) PERFORMS FACETTING : READ --> WRITE
+++ 0.600 P2: ONE_SLOT_BUFFER(1) IS IDLE
ONE_SLOT_BUFFER(1)
>>> 1.100 SYSTEM TERMINATED
>>> 1.100 INTERFACE(1) : ENTER CSSA COMMAND -

```

<==

fig. 2 : One Slot Buffer Execution Protocol

```
==> MAILBOX
MAILBOX OF INTERFACE(1) :
```

```
(1) *REPLY(STRING:"A")
```

```
>>> 1.100 INTERFACE(1) : ENTER CSSA COMMAND -
==> SEND RD TO OSB REPLY TO P; SEND WR("B") TO OSB; RUN
+++ 1.100 P1: INTERFACE(1) SENDS RD(), REPLY TO: P TO ONE_SLOT_BUFFER(1)
+++ 1.100 P1: INTERFACE(1) SENDS WR(STRING:"B") TO ONE_SLOT_BUFFER(1)
+++ 1.200 P2: ONE_SLOT_BUFFER(1) RECEIVES RD(), REPLY TO: P
      FROM INTERFACE(1)
+++ 1.200 P2: ONE_SLOT_BUFFER(1) IS IDLE
+++ 1.300 P2: ONE_SLOT_BUFFER(1) RECEIVES WR(STRING:"B")
      FROM INTERFACE(1)
+++ 1.300 P2: ONE_SLOT_BUFFER(1) STARTING OPERATION WR(STRING:
      "B")
+++ 1.500 P2: ONE_SLOT_BUFFER(1) PERFORMS FACETTING : WRITE --> READ
+++ 1.500 P2: ONE_SLOT_BUFFER(1) STARTING OPERATION RD()
+++ 1.600 P2: ONE_SLOT_BUFFER(1) SENDS *REPLY(STRING:"B")
      TO INTERFACE(1)
+++ 1.700 P2: ONE_SLOT_BUFFER(1) PERFORMS FACETTING : READ --> WRITE
+++ 1.700 P2: ONE_SLOT_BUFFER(1) IS IDLE
ONE_SLOT_BUFFER(1)
>>> 2.200 SYSTEM TERMINATED
>>> 2.200 INTERFACE(1) : ENTER CSSA COMMAND -
==> MAILBOX
MAILBOX OF INTERFACE(1) :
```

```
(1) *REPLY(STRING:"A")
(2) *REPLY(STRING:"B")
```

```
>>> 2.200 INTERFACE(1) : ENTER CSSA COMMAND -
==> DUMP OSB
```

```
+++ 2.200 RUNTIME STACK OF ONE_SLOT_BUFFER(1)
```

FACET WRITE ENV: SCRIPT ONE_SLOT_BUFFER LINE: 7 ..... SCRIPT ONE_SLOT_BUFFER ENV: LINE: 27 ..... BUFFER = "B"
---

```
>>> 2.200 INTERFACE(1) : ENTER CSSA COMMAND -
==> TERMINATE
+++ 2.200 ALL EXISTING AGENTS:
```

AGENT	FACET	OPERATION	MAILBOX
P1: INTERFACE(1) *			*REPLY *REPLY
P2: ONE_SLOT_BUFF	WRITE		

#### CSSA-SESSION-STATISTICS

=====

```
SESSION STARTED AT 11:48:01.00
SESSION TERMINATED AT 11:51:08.00 ON 1981/12/01
REAL-TIME USED : 188.00 SEC.
CPU-TIME USED : 0.88 SEC.
SIMULATION TIME USED : 2.2000 SEC.
NUMBER OF AGENTS CREATED : 1
NUMBER OF MESSAGES SENT : 6
```

	1	2	3	4	5	6	7	B L O C K N E S T I N G
1	type BOUNDED_BUFFER is script(int: MAX) assert MAX > 0							+1
2								
3	public: INS, REM;							
4								
5	functionhead MOD(int: P1,P2) returns int external;							
6	var array (0..MAX-1) of string: BUFFER;							
7	var int: INS_COUNT, REM_COUNT := 0;							
8								
9	facethead BUFFER_EMPTY;							
10	facethead BUFFER_FULL;							
11	facethead NO_CONSTRAINT;							
12								
13	facet BUFFER_EMPTY is							+2
14	include: INS;							-2
15	endfacet							*1
16								+3
17	facet BUFFER_FULL is							-3
18	include: REM;							*1
19	endfacet							+4
20								-4
21	facet NO_CONSTRAINT is							*1
22	include: INS, REM;							+5
23	endfacet							-6
24								-5
25	operation INS(string: MESSAGE) is							*1
26	BUFFER(INS_COUNT) := MESSAGE;							+7
27	INS_COUNT := MOD(INS_COUNT + 1, MAX);							+8
28	if INS_COUNT = REM_COUNT then replace by BUFFER_FULL;							-8
29	else replace by NO_CONSTRAINT; endif;							-7
30	endoperation							*1
31								
32	operation REM is							
33	reply (BUFFER(REM_COUNT));							
34	REM_COUNT := MOD(REM_COUNT + 1, MAX);							
35	if INS_COUNT = REM_COUNT then replace by BUFFER_EMPTY;							
36	else replace by NO_CONSTRAINT; endif;							
37	endoperation							
38								
39								
40	initial BUFFER_EMPTY							
41								
42	endscript							-1

fig. 3 : Bounded Buffer Script

PROGRAM GENERATED ON 1981/11/26 AT 11:25:31.00  
BY BMS-CSSA-COMPILER (VERS. 30 SEP 1981)

PROTOCOL OF CSSA SESSION ON 1981/11/27 AT 11:11:40.00

```

>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -
==> VAR AGENT : BB := NEW BOUNDED_BUFFER(3)
>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -
==> OPER : INS,REM; PORT : P
>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -
==> ,DISPLAY

```

IDENTIFIER	TYPE	VALUE
BB	AGENT	BOUNDED_BUFFER(1)
BOUNDED_BUFFER	SCRIPT	BOUNDED_BUFFER
INS	LITERAL	
P	PORT	
REM	LITERAL	

```

>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -
==> OBSERVE; SEND INS("111") TO BB; SEND INS("222") TO BB;RUN
+++ 0.000 P1: INTERFACE(1) SENDS INS(STRING:"111") TO BOUNDED_BUFFER(1)
+++ 0.000 P1: INTERFACE(1) SENDS INS(STRING:"222") TO BOUNDED_BUFFER(1)
+++ 0.000 P2: BOUNDED_BUFFER(1) IS IDLE
+++ 0.100 P2: BOUNDED_BUFFER(1) RECEIVES INS(STRING:"111")
      FROM INTERFACE(1)
+++ 0.100 P2: BOUNDED_BUFFER(1) STARTING OPERATION INS(STRING:
      "111")
NO_CONSTRAINT
+++ 0.500 P2: BOUNDED_BUFFER(1) IS IDLE
+++ 0.500 P2: BOUNDED_BUFFER(1) RECEIVES INS(STRING:"222")
      FROM INTERFACE(1)
+++ 0.500 P2: BOUNDED_BUFFER(1) STARTING OPERATION INS(STRING:
      "222")
NO_CONSTRAINT
+++ 0.900 P2: BOUNDED_BUFFER(1) IS IDLE
>>> 0.900 SYSTEM TERMINATED
>>> 0.900 INTERFACE(1) : ENTER CSSA COMMAND -
==> SEND REM TO BB REPLY TO P; SEND REM TO BB REPLY TO P;SEND REM TO BB REPLY TO P
+++ 0.900 P1: INTERFACE(1) SENDS REM(),REPLY TO: P TO BOUNDED_BUFFER(1)
+++ 0.900 P1: INTERFACE(1) SENDS REM(),REPLY TO: P TO BOUNDED_BUFFER(1)
+++ 0.900 P1: INTERFACE(1) SENDS REM(),REPLY TO: P TO BOUNDED_BUFFER(1)
>>> 0.900 INTERFACE(1) : ENTER CSSA COMMAND -
==> RUN; MAILBOX; MAILBOX BB
+++ 1.000 P2: BOUNDED_BUFFER(1) RECEIVES REM(),REPLY TO: P
      FROM INTERFACE(1)
+++ 1.000 P2: BOUNDED_BUFFER(1) STARTING OPERATION REM()
+++ 1.100 P2: BOUNDED_BUFFER(1) SENDS *REPLY(STRING:"111")
      TO INTERFACE(1)
+++ 1.100 P2: BOUNDED_BUFFER(1) RECEIVES REM(),REPLY TO: P
      FROM INTERFACE(1)
NO_CONSTRAINT
+++ 1.400 P2: BOUNDED_BUFFER(1) STARTING OPERATION REM()
+++ 1.500 P2: BOUNDED_BUFFER(1) SENDS *REPLY(STRING:"222")
      TO INTERFACE(1)
+++ 1.500 P2: BOUNDED_BUFFER(1) RECEIVES REM(),REPLY TO: P
      FROM INTERFACE(1)
BUFFER_EMPTY
+++ 1.800 P2: BOUNDED_BUFFER(1) IS IDLE
+++ 1.900 P1: INTERFACE(1) RECEIVES *REPLY(STRING:"111")
      FROM BOUNDED_BUFFER(1)
+++ 1.900 P1: INTERFACE(1) RECEIVES *REPLY(STRING:"222")
      FROM BOUNDED_BUFFER(1)
>>> 2.000 SYSTEM TERMINATED

```

fig. 4 : Bounded Buffer Execution Protocol

## MAILBOX OF INTERFACE(1) :

---

```
(1) *REPLY(STRING:"111")
(2) *REPLY(STRING:"222")
```

---

## MAILBOX OF BOUNDED\_BUFFER(1) :

---

```
(1) REM(),REPLY TO: P
```

---

```
>>> 2.000 INTERFACE(1) : ENTER CSSA COMMAND -
==> SEND INS("333") TO BB; SEND INS("444") TO BB; SEND INS("555") TO BB;RUN <==
+++ 2.000 P1: INTERFACE(1) SENDS INS(STRING:"333") TO BOUNDED_BUFFER(1)
+++ 2.000 P1: INTERFACE(1) SENDS INS(STRING:"444") TO BOUNDED_BUFFER(1)
+++ 2.000 P1: INTERFACE(1) SENDS INS(STRING:"555") TO BOUNDED_BUFFER(1)
+++ 2.100 P2: BOUNDED_BUFFER(1) RECEIVES INS(STRING:"333")
      FROM INTERFACE(1)
+++ 2.100 P2: BOUNDED_BUFFER(1) STARTING OPERATION INS(STRING:
      "333")
NO_CONSTRAINT
+++ 2.500 P2: BOUNDED_BUFFER(1) STARTING OPERATION REM()
+++ 2.600 P2: BOUNDED_BUFFER(1) SENDS *REPLY(STRING:"333")
      TO INTERFACE(1)
+++ 2.600 P2: BOUNDED_BUFFER(1) RECEIVES INS(STRING:"444")
      FROM INTERFACE(1)
+++ 2.600 P2: BOUNDED_BUFFER(1) RECEIVES INS(STRING:"555")
      FROM INTERFACE(1)
BUFFER_EMPTY
+++ 2.900 P2: BOUNDED_BUFFER(1) STARTING OPERATION INS(STRING:
      "444")
+++ 3.000 P1: INTERFACE(1) RECEIVES *REPLY(STRING:"333")
      FROM BOUNDED_BUFFER(1)
NO_CONSTRAINT
+++ 3.300 P2: BOUNDED_BUFFER(1) STARTING OPERATION INS(STRING:
      "555")
>>> 3.300 REAL-TIME LIMIT EXCEEDED
>>> 3.300 INTERFACE(1) : ENTER CSSA COMMAND -
==> SEND INS("666") TO BB; SEND INS("777") TO BB;RUN <==
+++ 3.300 P1: INTERFACE(1) SENDS INS(STRING:"666") TO BOUNDED_BUFFER(1)
+++ 3.300 P1: INTERFACE(1) SENDS INS(STRING:"777") TO BOUNDED_BUFFER(1)
+++ 3.600 P2: BOUNDED_BUFFER(1) RECEIVES INS(STRING:"666")
      FROM INTERFACE(1)
+++ 3.600 P2: BOUNDED_BUFFER(1) RECEIVES INS(STRING:"777")
      FROM INTERFACE(1)
NO_CONSTRAINT
+++ 3.700 P2: BOUNDED_BUFFER(1) STARTING OPERATION INS(STRING:
      "666")
BUFFER_FULL
+++ 4.100 P2: BOUNDED_BUFFER(1) IS IDLE
>>> 4.100 SYSTEM TERMINATED
>>> 4.100 INTERFACE(1) : ENTER CSSA COMMAND -
==> MAILBOX; MAILBOX BB <==
MAILBOX OF INTERFACE(1) :
(1) *REPLY(STRING:"111")
(2) *REPLY(STRING:"222")
(3) *REPLY(STRING:"333")
```

---

## MAILBOX OF BOUNDED\_BUFFER(1) :

---

```
(1) INS(STRING:"777")
```

---

```
>>> 4.100 INTERFACE(1) : ENTER CSSA COMMAND -
==> SEND REM TO BB REPLY TO P;RUN <==
+++ 4.100 P1: INTERFACE(1) SENDS REM(),REPLY TO: P TO BOUNDED_BUFFER(1)
+++ 4.200 P2: BOUNDED_BUFFER(1) RECEIVES REM(),REPLY TO: P
      FROM INTERFACE(1)
+++ 4.200 P2: BOUNDED_BUFFER(1) STARTING OPERATION REM()
+++ 4.300 P2: BOUNDED_BUFFER(1) SENDS *REPLY(STRING:"444")
      TO INTERFACE(1)
NO_CONSTRAINT
+++ 4.600 P2: BOUNDED_BUFFER(1) STARTING OPERATION INS(STRING:
      "777")
```

---

## BUFFER\_FULL

```

+++      5.000 P2: BOUNDED_BUFFER(1) IS IDLE
+++      5.100 P1: INTERFACE(1) RECEIVES *REPLY(STRING:"444")
              FROM BOUNDED_BUFFER(1)

```

```

>>>      5.200 SYSTEM TERMINATED
>>>      5.200 INTERFACE(1) : ENTER CSSA COMMAND -
==> MAILBOX
MAILBOX OF INTERFACE(1) :

```

```

(1) *REPLY(STRING:"111")
(2) *REPLY(STRING:"222")
(3) *REPLY(STRING:"333")
(4) *REPLY(STRING:"444")

```

```

>>>      5.200 INTERFACE(1) : ENTER CSSA COMMAND -
==> MAILBOX BB
MAILBOX OF BOUNDED_BUFFER(1) :

```

```

>>>      5.200 INTERFACE(1) : ENTER CSSA COMMAND -
==> DUMP BB

```

```

+++      5.200 RUNTIME STACK OF BOUNDED_BUFFER(1)

```

<pre> FACET  BUFFER_FULL ENV:   SCRIPT  BOUNDED_BUFFER LINE:  17 ..... SCRIPT BOUNDED_BUFFER ENV: LINE:  42 ..... MAX = 3 ANONYM = (0..2) BUFFER(0) = "777" BUFFER(1) = "555" BUFFER(2) = "666" INS_COUNT = 1 REM_COUNT = 1 </pre>
--

```

>>>      5.200 INTERFACE(1) : ENTER CSSA COMMAND -
==> TERMINATE
+++      5.200 ALL EXISTING AGENTS:

```

AGENT	FACET	OPERATION	MAILBOX
P1: INTERFACE(1) *			*REPLY *REPLY *REPLY *REPLY
P2: BOUNDED_BUFFER	BUFFER_FU		

CSSA-SESSION-STATISTICS  
=====

```

SESSION STARTED AT 11:11:19.00
SESSION TERMINATED AT 11:20:26.00 ON 1981/11/27
REAL-TIME USED : 547.00 SEC.
CPU-TIME USED : 1.88 SEC.
SIMULATION TIME USED : 5.2000 SEC.
NUMBER OF AGENTS CREATED : 1
NUMBER OF MESSAGES SENT : 15

```



```

bounded_buffer = monitor is create, append, remove;

am= array[message];
rep = record[ slots:am, max:int, nonempty, nonfull: condition]

create = proc(n:int) returns (cvt);
        return (rep${slots:am$new(),
                      max:n,
                      nonempty,nonfull: condition$create()});
        end create;

append = proc(buffer:cvt, x:message) ;
        if am$size(buffer.slots) = max
            then condition$wait(buffer.nonfull);
            end;
        am$addh(buffer.slots,x);
        condition$signal(buffer.nonempty);
        end append;

remove = proc(buffer:cvt) returns (message);
        if am$size(buffer.slots) = 0
            then condition$wait(buffer.nonempty);
            end;
        x:message := am$reml(slots);
        condition$signal(buffer.nonfull);
        return (x);
        end remove;

end bounded_buffer;

```

fig. 5 : Bounded Buffer Monitor

	1	2	3	4	5	6	7	BLOCKNESTING
1	type PROTECTED_BUFFER is							
2	script							+1
3								
4	public : INS, REM;							
5								
6	type DATA_BASE is script(int: SIZE)							+2
7								
8	var array(1..SIZE) of string: FILE;							
9								
10	facet RD_OR_WR is							+3
11	public: WR, RD, RD_ACKN;							
12	var int: RCOUNT := 0;							
13								
14	operation WR(string: DATA; int: KEY) is							+4
15								
16	reply;							
17	if KEY > 0 and KEY <= SIZE then							+5
18	FILE(KEY) := DATA; endif;							-5
19	endoperation							-4
20								+3
21	operation RD(int: KEY) is							+6
22								
23	RCOUNT := RCOUNT + 1;							
24	if KEY > 0 and KEY <= SIZE then							+7
25	reply (FILE(KEY)); endif;							-7
26	endoperation							-6
27								+3
28	operation RD_ACKN(int: COUNT) assert COUNT = RCOUNT is							+8
29								
30	reply;							
31	RCOUNT := 0;							-8
32	endoperation							+3
33	endfacet							-3
34								+2
35	initial RD_OR_WR							
36	endscript;							-2
37	(* back on top level of script PROTECTED_BUFFER *)							+1
38								
39	var DATA_BASE : DB;							
40	var int : REM_COUNT, INS_COUNT;							
41	functionhead MOD(int : P1,P2) returns int external;							
42	var int : MAX;							
43								
44	facethead BUFFER_EMPTY;							
45	facethead BUFFER_FULL;							
46	facethead NO_CONSTRAINT;							
47								
48	facet CREATION is							+9
49	public : CREATE;							
50								
51	operation CREATE(->MAX) is							+10
52								+11
53	if MAX > 0 then							
54	DB := new DATA_BASE(MAX);							
55	replace by BUFFER_EMPTY;							
56	else print(" Buffer size must be greater than 0 ");							-11
	endif;							

fig. 6 : Protected Buffer Script

BMS - C S S A - C O M P I L E R

1981/12/07 16:31

	1	2	3	4	5	6	7	B L O C K N E S T I N G
57   endoperation								-10
58   endfacet								-9
59								+1
60   facet BUFFER_EMPTY is								+12
61     include : INS;								
62   endfacet								-12
63								+1
64   facet NO_CONSTRAINT is								+13
65     include : INS, REM;								
66   endfacet								-13
67								+1
68   facet BUFFER_FULL is								+14
69     include : REM;								
70   endfacet								-14
71								+1
72   operation INS(string : M) is								+15
73     port : ACCEPTED;								
74     oper : WR;								
75     -----								
76     send WR(M, INS_COUNT + 1) to DB reply to ACCEPTED;								
77     INS_COUNT := MOD(INS_COUNT + 1, MAX);								
78     wait ACCEPTED;								
79     if INS_COUNT = REM_COUNT then replace by BUFFER_FULL;								+16
80     else replace by NO_CONSTRAINT;								
81     endif;								-16
82   endoperation								-15
83								+1
84   operation REM is								+17
85     port : ACCEPTED;								
86     var string : M;								
87     oper : RD;								
88     -----								
89     send RD(REM_COUNT + 1) to DB reply to ACCEPTED;								
90     REM_COUNT := MOD(REM_COUNT + 1, MAX);								
91     wait ACCEPTED(M);								
92     reply (M);								
93     if REM_COUNT = INS_COUNT then replace by BUFFER_EMPTY;								+18
94     else replace by NO_CONSTRAINT;								
95     endif;								-18
96   endoperation								-17
97								+1
98   initial CREATION								
99								
100   endscrip								-1

fig. 6 -continued-

# All

PROGRAM GENERATED ON 1981/12/07 AT 16:36:49.00  
BY BMS-CSSA-COMPILER (VERS. 30 SEP 1981)

PROTOCOL OF CSSA SESSION ON 1981/12/07 AT 16:54:19.00  
=====

>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -  
==> DISPLAY

<==

IDENTIFIER	TYPE	VALUE
DATA_BASE PROTECTED_BUFFER	SCRIPT SCRIPT	DATA_BASE PROTECTED_BUFFER

```

>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -
==> VAR AGENT : PB := NEW PROTECTED_BUFFER
>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -
==> OPER : CREATE,INS,REM; PORT : P
>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -
==> OBSERVE; SEND CREATE(3) TO PB;RUN
+++ 0.000 P1: INTERFACE(1) SENDS CREATE(INT:3) TO PROTECTED_BUFFER(1)
+++ 0.000 P2: PROTECTED_BUFFER(1) IS IDLE
+++ 0.100 P2: PROTECTED_BUFFER(1) RECEIVES CREATE(INT:3)
      FROM INTERFACE(1)
+++ 0.100 P2: PROTECTED_BUFFER(1) STARTING OPERATION CREATE(INT:
      3)
+++ 0.300 P2: PROTECTED_BUFFER(1) CREATES DATA_BASE(1)(INT:
      3)
+++ 0.300 P3: DATA_BASE(1) IS IDLE
BUFFER_EMPTY
+++ 0.400 P2: PROTECTED_BUFFER(1) IS IDLE
>>> 0.400 SYSTEM TERMINATED
>>> 0.400 INTERFACE(1) : ENTER CSSA COMMAND -
==> SEND REM TO PB REPLY TO P; RUN
+++ 0.400 P1: INTERFACE(1) SENDS REM(),REPLY TO: P TO PROTECTED_BUFFER(1)
+++ 0.500 P2: PROTECTED_BUFFER(1) RECEIVES REM(),REPLY TO: P
      FROM INTERFACE(1)
+++ 0.500 P2: PROTECTED_BUFFER(1) IS IDLE
>>> 0.500 SYSTEM TERMINATED
>>> 0.500 INTERFACE(1) : ENTER CSSA COMMAND -
==> SEND INS("A") TO PB;RUN
+++ 0.500 P1: INTERFACE(1) SENDS INS(STRING:"A") TO PROTECTED_BUFFER(1)
+++ 0.600 P2: PROTECTED_BUFFER(1) RECEIVES INS(STRING:"A")
      FROM INTERFACE(1)
+++ 0.600 P2: PROTECTED_BUFFER(1) STARTING OPERATION INS(STRING:
      "A")
+++ 0.700 P2: PROTECTED_BUFFER(1) SENDS WR(STRING:"A",INT:
      1),REPLY TO: ACCEPTED TO DATA_BASE(1)
+++ 0.900 P2: PROTECTED_BUFFER(1) IS IDLE
ACCEPTED FROM PROTECTED_BUFFER(1)
+++ 1.600 P3: DATA_BASE(1) STARTING OPERATION WR(STRING:
      "A",INT:1)
+++ 1.700 P3: DATA_BASE(1) SENDS *REPLY() TO PROTECTED_BUFFER(1)
+++ 1.900 P3: DATA_BASE(1) IS IDLE
+++ 2.600 P2: PROTECTED_BUFFER(1) RECEIVES *REPLY() FROM DATA_BASE(1)
NO_CONSTRAINT
+++ 2.800 P2: PROTECTED_BUFFER(1) STARTING OPERATION REM()
+++ 2.900 P2: PROTECTED_BUFFER(1) SENDS RD(INT:1),REPLY TO: ACCEPTED
      TO DATA_BASE(1)
+++ 3.100 P2: PROTECTED_BUFFER(1) IS IDLE
+++ 3.600 P3: DATA_BASE(1) RECEIVES RD(INT:1),REPLY TO: ACCEPTED
      FROM PROTECTED_BUFFER(1)
+++ 3.600 P3: DATA_BASE(1) STARTING OPERATION RD(INT:1)
+++ 3.900 P3: DATA_BASE(1) SENDS *REPLY(STRING:"A") TO PROTECTED_BUFFER(1)
+++ 3.900 P3: DATA_BASE(1) IS IDLE
+++ 4.600 P2: PROTECTED_BUFFER(1) RECEIVES *REPLY(STRING:
      "A") FROM DATA_BASE(1)
+++ 4.700 P2: PROTECTED_BUFFER(1) SENDS *REPLY(STRING:"A")
      TO INTERFACE(1)
BUFFER_EMPTY

```

<==

<==

<==

<==

<==

fig. 7 : Protected Buffer Execution Protocol

```

+++      4.900 P2: PROTECTED_BUFFER(1) IS IDLE
PROTECTED_BUFFER(1)
>>>      5.600 SYSTEM TERMINATED
>>>      5.600 INTERFACE(1) : ENTER CSSA COMMAND -
==> MAILBOX; SEND INS("B") TO PB; SEND INS("C") TO PB; RUN
MAILBOX OF INTERFACE(1) :
(1) *REPLY(STRING:"A")

+++      5.600 P1: INTERFACE(1) SENDS INS(STRING:"B") TO PROTECTED_BUFFER(1)
+++      5.600 P1: INTERFACE(1) SENDS INS(STRING:"C") TO PROTECTED_BUFFER(1)
+++      5.700 P2: PROTECTED_BUFFER(1) RECEIVES INS(STRING:"B")
FROM INTERFACE(1)
+++      5.700 P2: PROTECTED_BUFFER(1) STARTING OPERATION INS(STRING:
"B")
+++      5.800 P2: PROTECTED_BUFFER(1) SENDS WR(STRING:"B",INT:
2),REPLY TO: ACCEPTED TO DATA_BASE(1)
+++      5.800 P2: PROTECTED_BUFFER(1) RECEIVES INS(STRING:"C")
FROM INTERFACE(1)
+++      6.000 P2: PROTECTED_BUFFER(1) IS IDLE
ACCEPTED FROM PROTECTED_BUFFER(1)
+++      6.700 P3: DATA_BASE(1) STARTING OPERATION WR(STRING:
"B",INT:2)
+++      6.800 P3: DATA_BASE(1) SENDS *REPLY() TO PROTECTED_BUFFER(1)
+++      7.000 P3: DATA_BASE(1) IS IDLE
+++      7.700 P2: PROTECTED_BUFFER(1) RECEIVES *REPLY() FROM DATA_BASE(1)
NO_CONSTRAINT
+++      7.900 P2: PROTECTED_BUFFER(1) STARTING OPERATION INS(STRING:
"C")
+++      8.000 P2: PROTECTED_BUFFER(1) SENDS WR(STRING:"C",INT:
3),REPLY TO: ACCEPTED TO DATA_BASE(1)
+++      8.200 P2: PROTECTED_BUFFER(1) IS IDLE
ACCEPTED FROM PROTECTED_BUFFER(1)
+++      8.700 P3: DATA_BASE(1) STARTING OPERATION WR(STRING:
"C",INT:3)
+++      8.800 P3: DATA_BASE(1) SENDS *REPLY() TO PROTECTED_BUFFER(1)
+++      9.000 P3: DATA_BASE(1) IS IDLE
+++      9.700 P2: PROTECTED_BUFFER(1) RECEIVES *REPLY() FROM DATA_BASE(1)
NO_CONSTRAINT
+++      9.900 P2: PROTECTED_BUFFER(1) IS IDLE
>>>      10.000 SYSTEM TERMINATED
>>>      10.000 INTERFACE(1) : ENTER CSSA COMMAND -
==> SEND INS("D") TO PB; SEND INS("E") TO PB; RUN
+++      10.000 P1: INTERFACE(1) SENDS INS(STRING:"D") TO PROTECTED_BUFFER(1)
+++      10.000 P1: INTERFACE(1) SENDS INS(STRING:"E") TO PROTECTED_BUFFER(1)
+++      10.100 P2: PROTECTED_BUFFER(1) RECEIVES INS(STRING:"D")
FROM INTERFACE(1)
+++      10.100 P2: PROTECTED_BUFFER(1) STARTING OPERATION INS(STRING:
"D")
+++      10.200 P2: PROTECTED_BUFFER(1) SENDS WR(STRING:"D",INT:
1),REPLY TO: ACCEPTED TO DATA_BASE(1)
+++      10.200 P2: PROTECTED_BUFFER(1) RECEIVES INS(STRING:"E")
FROM INTERFACE(1)
+++      10.400 P2: PROTECTED_BUFFER(1) IS IDLE
ACCEPTED FROM PROTECTED_BUFFER(1)
+++      11.100 P3: DATA_BASE(1) STARTING OPERATION WR(STRING:
"D",INT:1)
+++      11.200 P3: DATA_BASE(1) SENDS *REPLY() TO PROTECTED_BUFFER(1)
+++      11.400 P3: DATA_BASE(1) IS IDLE
+++      12.100 P2: PROTECTED_BUFFER(1) RECEIVES *REPLY() FROM DATA_BASE(1)
BUFFER_FULL
+++      12.300 P2: PROTECTED_BUFFER(1) IS IDLE
>>>      12.300 SYSTEM TERMINATED
>>>      12.300 INTERFACE(1) : ENTER CSSA COMMAND -
==> NOOBSERVE; SEND REM TO PB REPLY TO P; SEND REM TO PB REPLY TO P; RUN
>>>      21.400 SYSTEM TERMINATED
>>>      21.400 INTERFACE(1) : ENTER CSSA COMMAND -
==> MAILBOX;
MAILBOX OF INTERFACE(1) :
(1) *REPLY(STRING:"A")
(2) *REPLY(STRING:"B")
(3) *REPLY(STRING:"C")
(4) *REPLY(STRING:"D")

```

```

>>> 21.400 INTERFACE(1) : ENTER CSSA COMMAND -
==> SEND REM TO PB REPLY TO P; SEND REM TO PB REPLY TO P; RUN
>>> 24.500 SYSTEM TERMINATED
>>> 24.500 INTERFACE(1) : ENTER CSSA COMMAND -
==> MAILBOX
MAILBOX OF INTERFACE(1) :

```

&lt;==

&lt;==

```

(1) *REPLY(STRING:"A")
(2) *REPLY(STRING:"B")
(3) *REPLY(STRING:"C")
(4) *REPLY(STRING:"D")
(5) *REPLY(STRING:"E")

```

```

>>> 24.500 INTERFACE(1) : ENTER CSSA COMMAND -
==> DUMP DATA_BASE(1); DUMP PB

```

&lt;==

```

+++ 24.500 RUNTIME STACK OF DATA_BASE(1)

```

FACET RD_OR_WR ENV: SCRIPT DATA_BASE LINE: 11 ..... RCOUNT = 5
SCRIPT DATA_BASE ENV: LINE: 35 ..... SIZE = 3 ANONYM = (1..3) FILE(1) = "D" FILE(2) = "E" FILE(3) = "C"

```

+++ 24.500 RUNTIME STACK OF PROTECTED_BUFFER(1)

```

FACET BUFFER_EMPTY ENV: SCRIPT PROTECTED_BUFFER LINE: 60 .....
SCRIPT PROTECTED_BUFFER ENV: LINE: 100 ..... DATA_BASE = DATA_BASE DB = DATA_BASE(1) REM_COUNT = 2 INS_COUNT = 2 MAX = 3

```

>>> 24.500 INTERFACE(1) : ENTER CSSA COMMAND -
>>> 24.600 SYSTEM TERMINATED
MAILBOX OF INTERFACE(1) :

```

```

(1) *REPLY(STRING:"A")
(2) *REPLY(STRING:"B")
(3) *REPLY(STRING:"C")
(4) *REPLY(STRING:"D")
(5) *REPLY(STRING:"E")

```

```
>>> 24.600 INTERFACE(1) : ENTER CSSA COMMAND -
==> SEND INS("F") TO PB; RUN; MAILBOX
>>> 29.700 SYSTEM TERMINATED
MAILBOX OF INTERFACE(1) :
```

&lt;==

```
(1) *REPLY(STRING:"A")
(2) *REPLY(STRING:"B")
(3) *REPLY(STRING:"C")
(4) *REPLY(STRING:"D")
(5) *REPLY(STRING:"E")
(6) *REPLY(STRING:"F")
```

```
>>> 29.700 INTERFACE(1) : ENTER CSSA COMMAND -
==> TERMINATE
+++ 29.700 ALL EXISTING AGENTS:
```

&lt;==

AGENT	FACET	OPERATION	MAILBOX
P1: INTERFACE(1) *			*REPLY *REPLY *REPLY *REPLY *REPLY *REPLY
P2: PROTECTED_BUF	BUFFER_EM		
P3: DATA_BASE(1)	RD_OR_WR		

CSSA-SESSION-STATISTICS  
=====

```
SESSION STARTED AT 16:52:59.00
SESSION TERMINATED AT 17:06:42.00 ON 1981/12/07
REAL-TIME USED : 823.00 SEC.
CPU-TIME USED : 3.06 SEC.
SIMULATION TIME USED : 29.7000 SEC.
NUMBER OF AGENTS CREATED : 2
NUMBER OF MESSAGES SENT : 43
```

```

protected_buffer = monitor is create, append, remove;

rep = record[ slots:buffer, nonempty, nonfull: condition ]

create = proc() returns (cvt);
    return (rep#{slots:buffer#create(),
        nonempty,nonfull: condition#create()});
    end create;

append = proc(pb:cvt, x:message) ;
    if buffer$full(pb.slots) then condition$wait(pb.nonfull) end;
    buffer$append(pb.slots, x);
    condition$signal(pb.nonempty);
    end append;

remove = proc(pb:cvt) returns (message);
    if buffer$empty(pb.slots) then condition$wait(pb.nonempty) end;
    x:message := buffer$remove(pb.slots);
    condition$signal(pb.nonfull);
    return (x);
    end remove;

end bounded_buffer;

```

fig. 8 : Protected Buffer Monitor



	1	2	3	4	5	6	7	B L O C K N E S T I N G
1	type RDER_Prio is script							+1
2								
3	type DATA_BASE is script(int: SIZE)							+2
4	var array(1..SIZE) of string: FILE;							
5	facet RD_OR_WR is							+3
6	public: WR, RD, RD_ACKN;							
7	var int: RCOUNT := 0;							
8	operation WR(string: DATA; int: KEY) is							+4
9	reply;							
10	if KEY > 0 and KEY <= SIZE then							+5
11	FILE(KEY) := DATA; endif;							-5
12	endoperation							-4
13	operation RD(int: KEY) is							+3
14	RCOUNT := RCOUNT + 1;							+6
15	if KEY > 0 and KEY <= SIZE then							+7
16	reply (FILE(KEY)); endif;							-7
17	endoperation							-6
18	operation RD_ACKN(int: COUNT) assert COUNT = RCOUNT is							+3
19	reply;							+8
20	RCOUNT := 0;							
21	endoperation							-8
22	endfacet							+3
23	initial RD_OR_WR							-3
24	endscript;							+2
25								
26	var DATA_BASE: DB;							-2
27	var int: RDER_COUNT := 0;							+1
28	var bool: FIRST_OPER := true;							
29	facethead WRRS;							
30								
31	facet CREATION is							+9
32	public: CREATE;							
33	operation CREATE(int: SIZE) is							+10
34	DB := new DATA_BASE(SIZE);							
35	replace by WRRS;							-10
36	endoperation							+9
37	endfacet							-9
38								+1
39	facet RDERs is							+11
40	public: RD;							
41	operation RD(int: KEY) is							+12
42	send RD(KEY) to DB inherit;							
43	RDER_COUNT := RDER_COUNT + 1;							
44	endoperation							

fig. 9 : Readers Priority Script

	1	2	3	4	5	6	7	B L O C K N E S T I N G
57	endoperation							-12
58								*11
59	operation idle is replace by WRRS;							+13
60	endoperation							-13
61								*11
62	endfacet							-11
63								*1
64	facet WRRS is							+14
65	public: WR;							
66	var bool: FIRST_WR := true;							
67								
68	operation WR(string: DATA; int: KEY) assert FIRST_WR is							+15
69	oper: RD_ACKN;							
70	port: R_ACKN, W_ACKN;							
71								
72	if FIRST_OPER then FIRST_OPER := false; endif;							+16-16
73	FIRST_WR := false;							*15
74								
75	if RDER_COUNT > 0 then							+17
76	send RD_ACKN(RDER_COUNT) to DB reply to R_ACKN;							
77	wait R_ACKN;							
78	RDER_COUNT := 0;							
79	endif;							-17
80	send WR(DATA,KEY) to DB reply to W_ACKN;							*15
81	wait W_ACKN;							
82	endoperation							-15
83								*14
84	operation idle is							+18
85	if not FIRST_OPER then replace by RDER; endif;							+19-19
86	endoperation							-18
87								*14
88	endfacet							-14
89								*1
90	initial CREATION							
91								
92	endscript							-1

fig. 9 -continued-

PROGRAM GENERATED ON 1981/12/07 AT 15:49:41.00  
BY BMS-CSSA-COMPILER (VERS. 30 SEP 1981)

PROTOCOL OF CSSA SESSION ON 1981/12/07 AT 16:02:06.00  
=====

>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -  
==> ,DISPLAY

<==

IDENTIFIER	TYPE	VALUE
DATA_BASE RDER_Prio	SCRIPT SCRIPT	DATA_BASE RDER_Prio

```

>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -
==> VAR AGENT : PRT := NEW RDER_Prio
>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -
==> OPER : RD,WR,CREATE; PORT : P;
>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -
==> OBSERVE; SEND CREATE(1) TO PRT; SEND WR("A",1) TO PRT;RUN
+++ 0.000 P1: INTERFACE(1) SENDS CREATE(INT:1) TO RDER_Prio(1)
+++ 0.000 P1: INTERFACE(1) SENDS WR(STRING:"A",INT:1) TO RDER_Prio(1)
+++ 0.000 P2: RDER_Prio(1) IS IDLE
+++ 0.100 P2: RDER_Prio(1) RECEIVES CREATE(INT:1) FROM INTERFACE(1)
+++ 0.100 P2: RDER_Prio(1) STARTING OPERATION CREATE(INT:
1)
+++ 0.200 P2: RDER_Prio(1) CREATES DATA_BASE(1)(INT:1)
+++ 0.200 P2: RDER_Prio(1) RECEIVES WR(STRING:"A",INT:1)
FROM INTERFACE(1)
+++ 0.200 P3: DATA_BASE(1) IS IDLE
+++ 0.300 P2: RDER_Prio(1) PERFORMS FACETTING : CREATION --> WRRS
+++ 0.300 P2: RDER_Prio(1) STARTING OPERATION WR(STRING:
"A",INT:1)
+++ 0.800 P2: RDER_Prio(1) SENDS WR(STRING:"A",INT:1),REPLY TO: W_ACKN
TO DATA_BASE(1)
+++ 0.900 P2: RDER_Prio(1) IS IDLE
+++ 1.100 P3: DATA_BASE(1) RECEIVES WR(STRING:"A",INT:1),REPLY TO: W_ACKN
FROM RDER_Prio(1)
+++ 1.100 P3: DATA_BASE(1) STARTING OPERATION WR(STRING:
"A",INT:1)
+++ 1.200 P3: DATA_BASE(1) SENDS *REPLY() TO RDER_Prio(1)
+++ 1.400 P3: DATA_BASE(1) IS IDLE
+++ 2.100 P2: RDER_Prio(1) RECEIVES *REPLY() FROM DATA_BASE(1)
+++ 2.100 P2: RDER_Prio(1) STARTING OPERATION IDLE
+++ 2.300 P2: RDER_Prio(1) PERFORMS FACETTING : WRRS --> RDERs
+++ 2.300 P2: RDER_Prio(1) STARTING OPERATION IDLE
+++ 2.400 P2: RDER_Prio(1) PERFORMS FACETTING : RDERs --> WRRS
+++ 2.400 P2: RDER_Prio(1) STARTING OPERATION IDLE
+++ 2.600 P2: RDER_Prio(1) PERFORMS FACETTING : WRRS --> RDERs
+++ 2.600 P2: RDER_Prio(1) STARTING OPERATION IDLE
+++ 2.700 P2: RDER_Prio(1) PERFORMS FACETTING : RDERs --> WRRS
+++ 2.700 P2: RDER_Prio(1) STARTING OPERATION IDLE
+++ 2.900 P2: RDER_Prio(1) PERFORMS FACETTING : WRRS --> RDERs
+++ 2.900 P2: RDER_Prio(1) STARTING OPERATION IDLE
+++ 3.000 P2: RDER_Prio(1) PERFORMS FACETTING : RDERs --> WRRS
+++ 3.000 P2: RDER_Prio(1) STARTING OPERATION IDLE
+++ 3.200 P2: RDER_Prio(1) PERFORMS FACETTING : WRRS --> RDERs
+++ 3.200 P2: RDER_Prio(1) STARTING OPERATION IDLE
+++ 3.300 P2: RDER_Prio(1) PERFORMS FACETTING : RDERs --> WRRS
+++ 3.300 P2: RDER_Prio(1) STARTING OPERATION IDLE
+++ 3.500 P2: RDER_Prio(1) PERFORMS FACETTING : WRRS --> RDERs
+++ 3.500 P2: RDER_Prio(1) STARTING OPERATION IDLE
+++ 3.600 P2: RDER_Prio(1) PERFORMS FACETTING : RDERs, --> WRRS
+++ 3.600 P2: RDER_Prio(1) STARTING OPERATION IDLE
+++ 3.800 P2: RDER_Prio(1) PERFORMS FACETTING : WRRS --> RDERs
+++ 3.800 P2: RDER_Prio(1) STARTING OPERATION IDLE
>>> 3.800 REAL-TIME LIMIT EXCEEDED
>>> 3.800 INTERFACE(1) : ENTER CSSA COMMAND -

```

fig. 10 : Readers Priority Execution Protocol

```

==> SEND RD(1) TO PRT REPLY TO P; SEND WR("B",1) TO PRT;
+++ 3.800 P1: INTERFACE(1) SENDS RD(INT:1),REPLY TO: P TO RDER_PRIO(1)
+++ 3.800 P1: INTERFACE(1) SENDS WR(STRING:"B",INT:1) TO RDER_PRIO(1)
>>> 3.800 INTERFACE(1) : ENTER CSSA COMMAND -
==> SEND RD(1) TO PRT REPLY TO P; SEND RD(1) TO PRT REPLY TO P;
+++ 3.800 P1: INTERFACE(1) SENDS RD(INT:1),REPLY TO: P TO RDER_PRIO(1)
+++ 3.800 P1: INTERFACE(1) SENDS RD(INT:1),REPLY TO: P TO RDER_PRIO(1)
>>> 3.800 INTERFACE(1) : ENTER CSSA COMMAND -
==> RUN;
+++ 3.900 P2: RDER_PRIO(1) PERFORMS FACETTING : RDER --> WRRS
+++ 3.900 P2: RDER_PRIO(1) STARTING OPERATION IDLE
+++ 4.100 P2: RDER_PRIO(1) PERFORMS FACETTING : WRRS --> RDER
+++ 4.100 P2: RDER_PRIO(1) STARTING OPERATION IDLE
+++ 4.100 P2: RDER_PRIO(1) RECEIVES RD(INT:1),REPLY TO: P
      FROM INTERFACE(1)
+++ 4.100 P2: RDER_PRIO(1) RECEIVES RD(INT:1),REPLY TO: P
      FROM INTERFACE(1)
+++ 4.100 P2: RDER_PRIO(1) RECEIVES RD(INT:1),REPLY TO: P
      FROM INTERFACE(1)
+++ 4.200 P2: RDER_PRIO(1) PERFORMS FACETTING : RDER --> WRRS
+++ 4.200 P2: RDER_PRIO(1) STARTING OPERATION IDLE
+++ 4.400 P2: RDER_PRIO(1) PERFORMS FACETTING : WRRS --> RDER
+++ 4.400 P2: RDER_PRIO(1) STARTING OPERATION RD(INT:1)
+++ 4.500 P2: RDER_PRIO(1) SENDS RD(INT:1),REPLY TO: P TO DATA_BASE(1)
+++ 4.500 P2: RDER_PRIO(1) RECEIVES WR(STRING:"B",INT:1)
      FROM INTERFACE(1)
+++ 4.600 P2: RDER_PRIO(1) STARTING OPERATION RD(INT:1)
+++ 4.700 P2: RDER_PRIO(1) SENDS RD(INT:1),REPLY TO: P TO DATA_BASE(1)
+++ 4.800 P2: RDER_PRIO(1) STARTING OPERATION RD(INT:1)
+++ 4.900 P2: RDER_PRIO(1) SENDS RD(INT:1),REPLY TO: P TO DATA_BASE(1)
+++ 4.900 P3: DATA_BASE(1) RECEIVES RD(INT:1),REPLY TO: P
      FROM RDER_PRIO(1)
+++ 4.900 P3: DATA_BASE(1) STARTING OPERATION RD(INT:1)
+++ 5.000 P2: RDER_PRIO(1) STARTING OPERATION IDLE
+++ 5.100 P2: RDER_PRIO(1) PERFORMS FACETTING : RDER --> WRRS
+++ 5.100 P2: RDER_PRIO(1) STARTING OPERATION WR(STRING:
      "B",INT:1)
+++ 5.200 P3: DATA_BASE(1) SENDS *REPLY(STRING:"A") TO INTERFACE(1)
+++ 5.200 P3: DATA_BASE(1) RECEIVES RD(INT:1),REPLY TO: P
      FROM RDER_PRIO(1)
+++ 5.200 P3: DATA_BASE(1) STARTING OPERATION RD(INT:1)
+++ 5.500 P2: RDER_PRIO(1) SENDS RD_ACKN(INT:3),REPLY TO: R_ACKN
      TO DATA_BASE(1)
+++ 5.500 P3: DATA_BASE(1) SENDS *REPLY(STRING:"A") TO INTERFACE(1)
+++ 5.500 P3: DATA_BASE(1) IS IDLE
+++ 5.600 P2: RDER_PRIO(1) IS IDLE
+++ 5.800 P1: INTERFACE(1) RECEIVES *REPLY(STRING:"A") FROM DATA_BASE(1)
+++ 5.800 P1: INTERFACE(1) RECEIVES *REPLY(STRING:"A") FROM DATA_BASE(1)
+++ 5.900 P3: DATA_BASE(1) RECEIVES RD_ACKN(INT:3),REPLY TO: R_ACKN
      FROM RDER_PRIO(1)
+++ 5.900 P3: DATA_BASE(1) IS IDLE
+++ 6.000 P3: DATA_BASE(1) RECEIVES RD(INT:1),REPLY TO: P
      FROM RDER_PRIO(1)
+++ 6.000 P3: DATA_BASE(1) STARTING OPERATION RD(INT:1)
+++ 6.300 P3: DATA_BASE(1) SENDS *REPLY(STRING:"A") TO INTERFACE(1)
+++ 6.300 P3: DATA_BASE(1) STARTING OPERATION RD_ACKN(INT:
      3)
+++ 6.400 P3: DATA_BASE(1) SENDS *REPLY() TO RDER_PRIO(1)
+++ 6.500 P3: DATA_BASE(1) IS IDLE
+++ 6.800 P1: INTERFACE(1) RECEIVES *REPLY(STRING:"A") FROM DATA_BASE(1)
+++ 6.900 P2: RDER_PRIO(1) RECEIVES *REPLY() FROM DATA_BASE(1)
+++ 7.100 P2: RDER_PRIO(1) SENDS WR(STRING:"B",INT:1),REPLY TO: W_ACKN
      TO DATA_BASE(1)
+++ 7.200 P2: RDER_PRIO(1) IS IDLE
+++ 7.900 P3: DATA_BASE(1) RECEIVES WR(STRING:"B",INT:1),REPLY TO: W_ACKN
      FROM RDER_PRIO(1)
+++ 7.900 P3: DATA_BASE(1) STARTING OPERATION WR(STRING:
      "B",INT:1)
+++ 8.000 P3: DATA_BASE(1) SENDS *REPLY() TO RDER_PRIO(1)
+++ 8.200 P3: DATA_BASE(1) IS IDLE
+++ 8.900 P2: RDER_PRIO(1) RECEIVES *REPLY() FROM DATA_BASE(1)
+++ 8.900 P2: RDER_PRIO(1) STARTING OPERATION IDLE
+++ 9.100 P2: RDER_PRIO(1) PERFORMS FACETTING : WRRS --> RDER
+++ 9.100 P2: RDER_PRIO(1) STARTING OPERATION IDLE
+++ 9.200 P2: RDER_PRIO(1) PERFORMS FACETTING : RDER --> WRRS
+++ 9.200 P2: RDER_PRIO(1) STARTING OPERATION IDLE
>>> 9.300 REAL-TIME LIMIT EXCEEDED
>>> 9.300 INTERFACE(1) : ENTER CSSA COMMAND -

```

```

+++ 9.400 P2: RDER_Prio(1) PERFORMS FACETTING : WRRS --> RDERs
+++ 9.400 P2: RDER_Prio(1) STARTING OPERATION IDLE
+++ 9.500 P2: RDER_Prio(1) PERFORMS FACETTING : RDERs --> WRRS
+++ 9.500 P2: RDER_Prio(1) STARTING OPERATION IDLE
+++ 9.700 P2: RDER_Prio(1) PERFORMS FACETTING : WRRS --> RDERs
+++ 9.700 P2: RDER_Prio(1) STARTING OPERATION IDLE
+++ 9.800 P2: RDER_Prio(1) PERFORMS FACETTING : RDERs --> WRRS
+++ 9.800 P2: RDER_Prio(1) STARTING OPERATION IDLE
+++ 10.000 P2: RDER_Prio(1) PERFORMS FACETTING : WRRS --> RDERs
+++ 10.000 P2: RDER_Prio(1) STARTING OPERATION IDLE
+++ 10.100 P2: RDER_Prio(1) PERFORMS FACETTING : RDERs --> WRRS
+++ 10.100 P2: RDER_Prio(1) STARTING OPERATION IDLE
+++ 10.300 P2: RDER_Prio(1) PERFORMS FACETTING : WRRS --> RDERs
+++ 10.300 P2: RDER_Prio(1) STARTING OPERATION IDLE
+++ 10.400 P2: RDER_Prio(1) PERFORMS FACETTING : RDERs --> WRRS
+++ 10.400 P2: RDER_Prio(1) STARTING OPERATION IDLE
+++ 10.600 P2: RDER_Prio(1) PERFORMS FACETTING : WRRS --> RDERs
+++ 10.600 P2: RDER_Prio(1) STARTING OPERATION IDLE
+++ 10.700 P2: RDER_Prio(1) PERFORMS FACETTING : RDERs --> WRRS
+++ 10.700 P2: RDER_Prio(1) STARTING OPERATION IDLE
>>> 10.700 REAL-TIME LIMIT EXCEEDED
>>> 10.700 INTERFACE(1) : ENTER CSSA COMMAND -
==> TERMINATE
+++ 10.700 ALL EXISTING AGENTS:

```

&lt;==

AGENT	FACET	OPERATION	MAILBOX
P1: INTERFACE(1) *			*REPLY *REPLY *REPLY
P2: RDER_Prio(1) *	WRRS	IDLE	
P3: DATA_BASE(1)	RD_OR_WR		

#### CSSA-SESSION-STATISTICS

=====

```

SESSION STARTED AT 15:56:03.00
SESSION TERMINATED AT 16:17:13.00 ON 1981/12/07
REAL-TIME USED : 1270.00 SEC.
CPU-TIME USED : 2.07 SEC.
SIMULATION TIME USED : 10.7000 SEC.
NUMBER OF AGENTS CREATED : 2
NUMBER OF MESSAGES SENT : 18

```

```

readers_priority = monitor is create,
                    startread,
                    endread,
                    startwrite,
                    endwrite;

rep = record[readercount: int,
             busy:boolean,
             readers, writers:condition];

create = proc() returns (cvt);
        return(rep${readercount: 0,
                    busy:false,
                    readers,writers:condition$create{}});
        end create;

startread = proc(m:cvt);
            if m.busy then condition$wait(m.readers) end;
            m.readercount:= m.readercount + 1;
            condition$signal(m.readers);
            end startread;

endread = proc(m:cvt);
        m.readercount:= m.readercount - 1;
        if m.readercount:=0
            then condition$signal(m.writers)
            end;
        end endread;

startwrite = proc(m:cvt);
            if m.readercount > 0 | m.busy
                then condition$wait(m.writers)
                end;
            m.busy:=true;
            end startwrite;

endwrite = proc(m:cvt);
        m.busy:=false;
        if condition$queue(m.readers)
            then condition$signal(m.readers)
            else condition$signal(m.writers)
            end;
        end endwrite;

end readers_priority;

```

fig. 11a : Readers Priority Monitor

```

protected_data_base = cluster is create,read,write;

rep = record[m: readers_priority,d: data_base]

create = proc()returns(cvt);
        return (rep${m: readers_priority$create(),
                    d: data_base$create()});
        end create;

read = proc(pdb: cvt) returns(data);
        readers_priority$startread(pdb.m);
        x:data :=data_base$read(pdb.d);
        readers_priority$endread(pdb.m);
        return (x);
        end read;

write = proc(pdb: cvt, x:data);
        readers_priority$startwrite(pdb.m);
        data_base$write((pdb.d, x);
        readers_priority$endwrite(pdb.m);
        end write;

end protected_data_base;

```

fig. 11b : Readers Priority Protected Resource Module

```

first_come_first_serve = monitor is create, startread, endread, startwrite, endwrite;

rep = record[ busy: boolean,
               readercount: integer,
               users, writer: condition]

create = proc() returns (cvt);
        return(rep#{busy:false, readercount:0, users, writer: condition$create()});
        end create;

startread = proc(m: cvt)
            if m.busy | condition$queue(m.writer) | condition$queue(m.users)
                then condition$wait(m.users);
                end;
            m.readercount:=m.readercount + 1;
            condition$signal(m.users); %start all readers
        end startread;

endread = proc(m:cvt);
            m.readercount := m.readercount - 1;
            if m.readercount = 0
                then if condition$queue(m.writer)
                        then condition$signal(m.writer)
                        else condition$signal(m.users)
                    end;
                end;
            %anyone on the writers queue has been waiting longer than those on users queue
        end endread;

startwrite = proc(m:cvt);
            if m.readercount > 0 | m.busy
                then condition$wait(m.users);
                end;
            if m.readercount > 0
                then condition$wait(m.writer);
                end;
            m.busy := true;
        end startwrite;

endwrite = proc(m:cvt);
            m.busy:=false;
            condition$signal(m.users);
        end endwrite;

end first_come_first_serve;

```

fig. 12 : First Come First Serve Monitor



	1	2	3	4	5	6	7	BLOCKNESTING
1	type	ALARM_CLOCK	is					
2	script							+1
3	var	int:	NOW := 0;					
4								+2
5	facet	REGISTRY	is					+3
6	public:	TICK, WAKE_ME;						
7								+4
8	facet	RESUME	is					-4
9	private:	WAKE_UP;						+3
10								+5
11	operation	WAKE_UP(int: TIME)	assert TIME=NOW is					-5
12		reply;						-3
13	endoperation							+2
14								+6
15	operation	idle	is leave;					+7
16	endoperation							-7
17	endfacet							-6
18								+2
19								+8
20	operation	WAKE_ME(int: DELAY_TIME)	is					-8
21		oper: WAKE_UP;						+2
22								-2
23		if DELAY_TIME <= 0 then	reply;					+1
24		else send WAKE_UP(DELAY_TIME + NOW)	to self inherit; endif;					
25	endoperation							-1
26								
27	operation	TICK	is					
28		NOW := NOW + 1;						
29		setup RESUME;						
30	endoperation							
31								
32	endfacet							
33								
34	initial	REGISTRY						
35								
36	endscript							

fig. 13 : Alarmclock Script (Wrong Solution)

	1	2	3	4	5	6	7	B L O C K N E S T I N G
1	type ALARM_CLOCK is							+1
2	script							+2
3								
4	facet REGISTRY is							+3
5	public : TICK, WAKE_ME;							-3
6								+2
7	type CALL_ELEM is record int : WAKE_UP_TIME;							
8	--> int : NO_OF_SLEEPERS;							
9	endrecord;							
10	relation CALL_LIST of CALL_ELEM;							
11	var int : NOW := 0;							+4
12								
13	facet RESUME is							+5
14	private : WAKE_UP;							+6
15								+7
16	operation WAKE_UP(int : TIME) assert TIME = NOW is							
17	reply;							
18	find CALL_LIST(TIME) -> CE do							
19	if CE.NO_OF_SLEEPERS = 1 then							
20	delete CE in CALL_LIST;							
21	leave;							
22	else CE.NO_OF_SLEEPERS := CE.NO_OF_SLEEPERS - 1;							-7
23	endif;							-6
24	endfind;							-5
25	endoperation							-4
26	endfacet							+2
27								+8
28	operation WAKE_ME(int : DELAY_TIME) is							
29	oper : WAKE_UP;							
30	var CALL_ELEM : CE;							
31								
32	if DELAY_TIME <= 0 then reply;							+9
33	else find CALL_LIST(DELAY_TIME + NOW) -> CE1 do							+10
34	CE1.NO_OF_SLEEPERS := CE1.NO_OF_SLEEPERS + 1;							
35	otherwise do							
36	CE.WAKE_UP_TIME := DELAY_TIME + NOW;							
37	CE.NO_OF_SLEEPERS := 1;							
38	insert CE into CALL_LIST;							
39	endif;							-10
40	send WAKE_UP(DELAY_TIME + NOW) to self inherit;							+9
41	endif;							-9
42	endoperation							-8
43								+2
44	operation TICK is							+11
45	NOW := NOW + 1;							
46	find CALL_LIST(NOW) -> CE do							+12
47	setup RESUME;							
48	endfind;							-12
49	endoperation							-11
50								+2
51	endfacet							-2
52								+1
53	initial REGISTRY							
54								
55	endscript							-1

fig. 14 : Alarmclock Script (Correct Solution)

PROGRAM GENERATED ON 1981/12/22 AT 11:59:34.00  
BY BMS-CSSA-COMPILER (VERS. 30 SEP 1981)

PROTOCOL OF CSSA SESSION ON 1982/01/18 AT 11:05:29.00

>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -  
==> ,DISPLAY

<==

IDENTIFIER	TYPE	VALUE
ALARM_CLOCK	SCRIPT	ALARM_CLOCK

```

>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -
==> VAR AGENT : AC := NEW ALARM_CLOCK;
>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -
==> OPER : TICK,WAKE_ME,PORT : P
>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -
==> OBSERVE;
>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -
==> SEND WAKE_ME(3) TO AC REPLY TO P; RUN
+++ 0.000 P1: INTERFACE(1) SENDS WAKE_ME(INT:3),REPLY TO: P
      TO ALARM_CLOCK(1)
+++ 0.000 P2: ALARM_CLOCK(1) IS IDLE
+++ 0.100 P2: ALARM_CLOCK(1) RECEIVES WAKE_ME(INT:3),REPLY TO: P
      FROM INTERFACE(1)
+++ 0.100 P2: ALARM_CLOCK(1) STARTING OPERATION WAKE_ME(INT:
      3)
+++ 0.700 P2: ALARM_CLOCK(1) SENDS WAKE_UP(INT:3),REPLY TO: P
      TO ALARM_CLOCK(1)
+++ 0.700 P2: ALARM_CLOCK(1) RECEIVES WAKE_UP(INT:3),REPLY TO: P
      FROM ALARM_CLOCK(1)
+++ 0.700 P2: ALARM_CLOCK(1) IS IDLE
>>> 0.700 SYSTEM TERMINATED
>>> 0.700 INTERFACE(1) : ENTER CSSA COMMAND -
==> SEND TICK TO AC;RUN
+++ 0.700 P1: INTERFACE(1) SENDS TICK() TO ALARM_CLOCK(1)
+++ 0.800 P2: ALARM_CLOCK(1) RECEIVES TICK() FROM INTERFACE(1)
+++ 0.800 P2: ALARM_CLOCK(1) STARTING OPERATION TICK()
+++ 1.000 P2: ALARM_CLOCK(1) IS IDLE
>>> 1.000 SYSTEM TERMINATED
>>> 1.000 INTERFACE(1) : ENTER CSSA COMMAND -
==> SEND WAKE_ME(2) TO AC REPLY TO P; RUN
+++ 1.100 P1: INTERFACE(1) SENDS WAKE_ME(INT:2),REPLY TO: P
      TO ALARM_CLOCK(1)
+++ 1.200 P2: ALARM_CLOCK(1) RECEIVES WAKE_ME(INT:2),REPLY TO: P
      FROM INTERFACE(1)
+++ 1.200 P2: ALARM_CLOCK(1) STARTING OPERATION WAKE_ME(INT:
      2)
+++ 1.600 P2: ALARM_CLOCK(1) SENDS WAKE_UP(INT:3),REPLY TO: P
      TO ALARM_CLOCK(1)
+++ 1.600 P2: ALARM_CLOCK(1) RECEIVES WAKE_UP(INT:3),REPLY TO: P
      FROM ALARM_CLOCK(1)
+++ 1.600 P2: ALARM_CLOCK(1) IS IDLE
>>> 1.600 SYSTEM TERMINATED
>>> 1.600 INTERFACE(1) : ENTER CSSA COMMAND -
==> SEND TICK TO AC
+++ 1.600 P1: INTERFACE(1) SENDS TICK() TO ALARM_CLOCK(1)
>>> 1.600 INTERFACE(1) : ENTER CSSA COMMAND -
==> TRACE AC; SEND TICK TO AC; RUN
+++ 1.600 P1: INTERFACE(1) SENDS TICK() TO ALARM_CLOCK(1)
+++ 1.700 P2: ALARM_CLOCK(1) RECEIVES TICK() FROM INTERFACE(1)

```

<==

<==

<==

<==

fig. 15 : Alarmclock Execution Protocol

```

1.700  TRACING ALARM_CLOCK(1) / REGISTRY
SEARCHING NEXT_MESSAGE
FOUND : WAKE_UP(INT:3),REPLY TO: P
SEARCHING NEXT_MESSAGE
FOUND : WAKE_UP(INT:3),REPLY TO: P
SEARCHING NEXT_MESSAGE
FOUND : TICK()
STARTING OPERATION TICK()
LINE 45: NOW := 2
LINE 46: WAKE_UP_TIME := 2
LINE 46: RECORD NOT FOUND

>>> 1.700 REAL-TIME LIMIT EXCEEDED
>>> 1.700 INTERFACE(1) : ENTER CSSA COMMAND -
==>
>>> 1.700 INTERFACE(1) : ENTER CSSA COMMAND -
==> RUN

1.900  TRACING ALARM_CLOCK(1) / REGISTRY / TICK

1.900  TRACING ALARM_CLOCK(1) / REGISTRY
SEARCHING NEXT_MESSAGE
FOUND : WAKE_UP(INT:3),REPLY TO: P
SEARCHING NEXT_MESSAGE
FOUND : WAKE_UP(INT:3),REPLY TO: P
SEARCHING NEXT_MESSAGE - NO MESSAGE FOUND

+++ 1.900 P2: ALARM_CLOCK(1) IS IDLE
+++ 1.900 P2: ALARM_CLOCK(1) RECEIVES TICK() FROM INTERFACE(1)

1.900  TRACING ALARM_CLOCK(1) / REGISTRY
SEARCHING NEXT_MESSAGE
FOUND : WAKE_UP(INT:3),REPLY TO: P
SEARCHING NEXT_MESSAGE
FOUND : WAKE_UP(INT:3),REPLY TO: P
SEARCHING NEXT_MESSAGE
FOUND : TICK()
STARTING OPERATION TICK()
LINE 45: NOW := 3
LINE 46: WAKE_UP_TIME := 3
LINE 46: RECORD FOUND

>>> 1.900 REAL-TIME LIMIT EXCEEDED
>>> 1.900 INTERFACE(1) : ENTER CSSA COMMAND -
==> RUN

2.200  TRACING ALARM_CLOCK(1) / REGISTRY / TICK
LINE 47: FACETTING : REGISTRY --> RESUME
SEARCHING NEXT_MESSAGE
FOUND : WAKE_UP(INT:3),REPLY TO: P
LINE 16: TIME := 3
STARTING OPERATION WAKE_UP(INT:3)

2.300  TRACING ALARM_CLOCK(1) / RESUME / WAKE_UP
LINE 17: SEND *REPLY() TO INTERFACE(1)

2.300  TRACING ALARM_CLOCK(1) / RESUME / WAKE_UP
LINE 18: WAKE_UP_TIME := 3
LINE 18: RECORD FOUND
LINE 22: NO_OF_SLEEPERS := 1

2.600  TRACING ALARM_CLOCK(1) / RESUME / WAKE_UP

2.600  TRACING ALARM_CLOCK(1) / RESUME
SEARCHING NEXT_MESSAGE
FOUND : WAKE_UP(INT:3),REPLY TO: P
LINE 16: TIME := 3
STARTING OPERATION WAKE_UP(INT:3)

```

```

>>> 2.600 REAL-TIME LIMIT EXCEEDED
>>> 2.600 INTERFACE(1) : ENTER CSSA COMMAND -
==> RUN
+++ 2.600 P1: INTERFACE(1) RECEIVES *REPLY() FROM ALARM_CLOCK(1)
-----
2.700 TRACING ALARM_CLOCK(1) / RESUME / WAKE_UP
LINE 17: SEND *REPLY() TO INTERFACE(1)
-----
2.700 TRACING ALARM_CLOCK(1) / RESUME / WAKE_UP
LINE 18: WAKE_UP_TIME := 3
LINE 18: RECORD FOUND
LINE 20: RECORD DELETED
-----
3.100 TRACING ALARM_CLOCK(1) / RESUME / WAKE_UP
LINE 21: FACETTING : RESUME --> REGISTRY
-----
3.100 TRACING ALARM_CLOCK(1) / REGISTRY
SEARCHING NEXT_MESSAGE - NO MESSAGE FOUND
-----
+++ 3.100 P2: ALARM_CLOCK(1) IS IDLE
+++ 3.600 P1: INTERFACE(1) RECEIVES *REPLY() FROM ALARM_CLOCK(1)
>>> 3.700 SYSTEM TERMINATED
>>> 3.700 INTERFACE(1) : ENTER CSSA COMMAND -
==> SEND WAKE_ME(0) TO AC REPLY TO P; RUN
+++ 3.700 P1: INTERFACE(1) SENDS WAKE_ME(INT:0), REPLY TO: P
TO ALARM_CLOCK(1)
+++ 3.800 P2: ALARM_CLOCK(1) RECEIVES WAKE_ME(INT:0), REPLY TO: P
FROM INTERFACE(1)
-----
3.800 TRACING ALARM_CLOCK(1) / REGISTRY
SEARCHING NEXT_MESSAGE
FOUND : WAKE_ME(INT:0), REPLY TO: P
LINE 28: DELAY_TIME := 0
STARTING OPERATION WAKE_ME(INT:0)
-----
4.000 TRACING ALARM_CLOCK(1) / REGISTRY / WAKE_ME
LINE 32: SEND *REPLY() TO INTERFACE(1)
-----
4.000 TRACING ALARM_CLOCK(1) / REGISTRY / WAKE_ME
-----
4.000 TRACING ALARM_CLOCK(1) / REGISTRY / WAKE_ME
-----
4.000 TRACING ALARM_CLOCK(1) / REGISTRY
SEARCHING NEXT_MESSAGE - NO MESSAGE FOUND
-----
+++ 4.000 P2: ALARM_CLOCK(1) IS IDLE
>>> 4.000 REAL-TIME LIMIT EXCEEDED
>>> 4.000 INTERFACE(1) : ENTER CSSA COMMAND -
==> NOTRACE AC
>>> 4.000 INTERFACE(1) : ENTER CSSA COMMAND -
==> TERMINATE
+++ 4.000 ALL EXISTING AGENTS:

```

AGENT	FACET	OPERATION	MAILBOX
P1: INTERFACE(1) *			*REPLY *REPLY
P2: ALARM_CLOCK(1)	REGISTRY		

CSSA-SESSION-STATISTICS  
=====

```

SESSION STARTED AT 11:05:24.00
SESSION TERMINATED AT 11:15:47.00 ON 1982/01/18
REAL-TIME USED : 623.00 SEC.
CPU-TIME USED : 2.02 SEC.
SIMULATION TIME USED : 4.0000 SEC.
NUMBER OF AGENTS CREATED : 1
NUMBER OF MESSAGES SENT : 11

```

```

alarmclock = monitor is create, wakeme, tick;

pq=priority_queue;

rep= record[wakeup: pq, now: int];

create = proc() returns(cvt);
    return (rep${wakeup: pq$create(), now: 0});
end create;

wakeme = proc(ac: cvt, time: int)
    alarmsetting: int := time+ac.now;
    while ac.now < alarmsetting do
        pq$wait(ac.wakeup, alarmsetting)
    end;
    %the while statement is necessary because the first process on the
    %queue is awakened every tick.
    pq$signal(ac.wakeup);
    %in case the next process has same wakeup time.
end wakeme;

tick = procedure(ac:cvt);
    ac.now := ac.now + 1;
    pq$signal(ac.wakeup);
end tick;

end alarmclock;

```

fig. 16 : Alarmclock Monitor

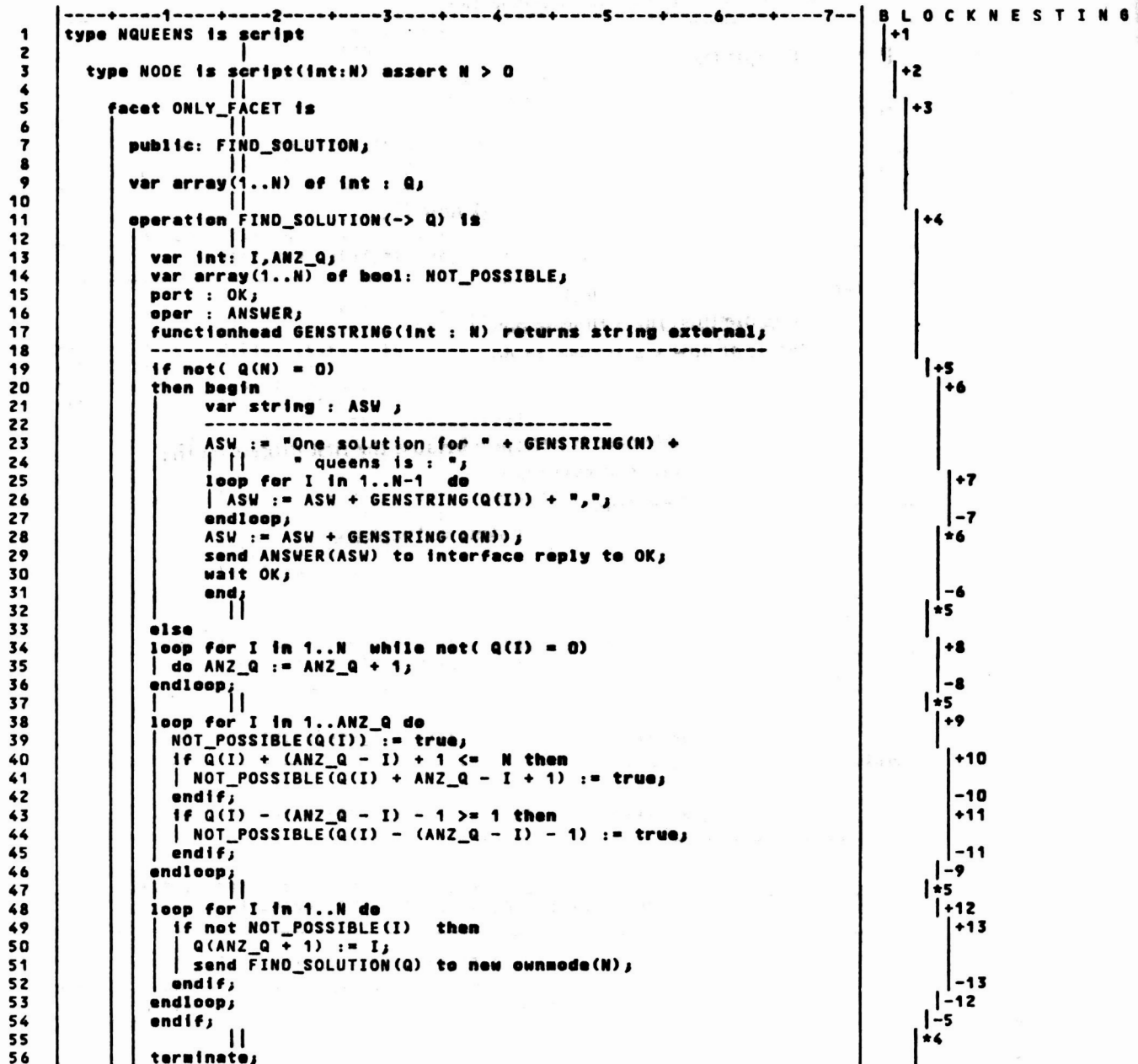


fig. 17 : N-Queens Script



fig. 17 -continued-



PROGRAM GENERATED ON 1982/01/18 AT 11:26:35.00  
BY BMS-CSSA-COMPILER (VERS. 30 SEP 1981)

PROTOCOL OF CSSA SESSION ON 1982/01/18 AT 11:43:19.00

>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -  
==> ,DISPLAY

<==

IDENTIFIER	TYPE	VALUE
NODE NQUEENS	SCRIPT SCRIPT	NODE NQUEENS

```

>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -
==> VAR AGENT : NQ := NEW NQUEENS; OPER : FIND_SOLUTION
>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -
==> OBSERVE, SEND FIND_SOLUTION(4) TO NQ; RUN
+++ 0.000 P1: INTERFACE(1) SENDS FIND_SOLUTION(INT:4) TO NQUEENS(1)
+++ 0.000 P2: NQUEENS(1) IS IDLE
+++ 0.100 P2: NQUEENS(1) RECEIVES FIND_SOLUTION(INT:4) FROM INTERFACE(1)
+++ 0.100 P2: NQUEENS(1) STARTING OPERATION FIND_SOLUTION(INT:
      4)
+++ 0.500 P2: NQUEENS(1) CREATES NODE(1) (INT:4)
+++ 0.500 P2: NQUEENS(1) SENDS FIND_SOLUTION(INT:1,INT:
      0,INT:0,INT:0) TO NODE(1)
+++ 0.500 P3: NODE(1) IS IDLE
+++ 0.700 P2: NQUEENS(1) CREATES NODE(2) (INT:4)
+++ 0.700 P2: NQUEENS(1) SENDS FIND_SOLUTION(INT:2,INT:
      0,INT:0,INT:0) TO NODE(2)
+++ 0.700 P4: NODE(2) IS IDLE
+++ 0.800 P4: NODE(2) RECEIVES FIND_SOLUTION(INT:2,INT:
      0,INT:0,INT:0) FROM NQUEENS(1)
+++ 0.800 P4: NODE(2) STARTING OPERATION FIND_SOLUTION(INT:
      2,INT:0,INT:0,INT:0)
+++ 0.900 P2: NQUEENS(1) CREATES NODE(3) (INT:4)
+++ 0.900 P2: NQUEENS(1) SENDS FIND_SOLUTION(INT:3,INT:
      0,INT:0,INT:0) TO NODE(3)
+++ 0.900 P5: NODE(3) IS IDLE
+++ 1.000 P5: NODE(3) RECEIVES FIND_SOLUTION(INT:3,INT:
      0,INT:0,INT:0) FROM NQUEENS(1)
+++ 1.000 P5: NODE(3) STARTING OPERATION FIND_SOLUTION(INT:
      3,INT:0,INT:0,INT:0)
+++ 1.100 P2: NQUEENS(1) CREATES NODE(4) (INT:4)
+++ 1.100 P2: NQUEENS(1) SENDS FIND_SOLUTION(INT:4,INT:
      0,INT:0,INT:0) TO NODE(4)
+++ 1.100 P6: NODE(4) IS IDLE
+++ 1.100 P2: NQUEENS(1) IS IDLE
+++ 1.100 P3: NODE(1) RECEIVES FIND_SOLUTION(INT:1,INT:
      0,INT:0,INT:0) FROM NQUEENS(1)
+++ 1.100 P3: NODE(1) STARTING OPERATION FIND_SOLUTION(INT:
      1,INT:0,INT:0,INT:0)
+++ 2.300 P5: NODE(3) CREATES NODE(5) (INT:4)
+++ 2.300 P5: NODE(3) SENDS FIND_SOLUTION(INT:3,INT:1,INT:
      0,INT:0) TO NODE(5)
+++ 2.300 P7: NODE(5) IS IDLE
+++ 2.400 P4: NODE(2) CREATES NODE(6) (INT:4)
+++ 2.400 P4: NODE(2) SENDS FIND_SOLUTION(INT:2,INT:4,INT:
      0,INT:0) TO NODE(6)

```

<==

<==

fig. 18 : N-Queens Execution Protocol

```

+++      2.400 P2: NODE(6) IS IDLE
+++      2.500 P2: NODE(6) RECEIVES FIND_SOLUTION(INT:2,INT:
              4,INT:0,INT:0) FROM NODE(2)
+++      2.500 P4: NODE(2) TERMINATED
+++      2.500 P2: NODE(6) STARTING OPERATION FIND_SOLUTION(INT:
              2,INT:4,INT:0,INT:0)
+++      2.500 P3: NODE(1) CREATES NODE(7) (INT:4)
+++      2.500 P3: NODE(7) IS IDLE
+++      2.500 P3: NODE(1) SENDS FIND_SOLUTION(INT:1,INT:3,INT:
              0,INT:0) TO NODE(7)
+++      2.500 P3: NODE(7) RECEIVES FIND_SOLUTION(INT:1,INT:
              3,INT:0,INT:0) FROM NODE(1)
+++      2.500 P3: NODE(7) STARTING OPERATION FIND_SOLUTION(INT:
              1,INT:3,INT:0,INT:0)
>>>      2.500 REAL-TIME LIMIT EXCEEDED
>>>      2.500 INTERFACE(1) : ENTER CSSA COMMAND -
==> ,STATUS
+++      2.500 ALL EXISTING AGENTS:

```

&lt;==

AGENT	FACET	OPERATION	MAILBOX
P1: INTERFACE(1) *			
P2: NQUEENS(1)	ONLY_FACE		
P3: NODE(1) *	ONLY_FACE	FIND_SOLU	N
P5: NODE(3) *	ONLY_FACE	FIND_SOLU	N
P6: NODE(4)	ONLY_FACE		
P7: NODE(5)	ONLY_FACE		
P2: NODE(6) *	ONLY_FACE	FIND_SOLU	N
P3: NODE(7) *	ONLY_FACE	FIND_SOLU	N

```

>>>      2.500 INTERFACE(1) : ENTER CSSA COMMAND -
==> ,SYSSTATUS
+++      2.500 SYSTEM STATUS:

```

&lt;==

PROC.	UTIL.	AGENTS
P1	100%	INTERFACE(1)*
P2	0%	NODE(6)* NQUEENS(1)
P3	56%	NODE(7)* NODE(1)*
P4	68%	
P5	60%	NODE(3)*
P6	1%	NODE(4)
P7	1%	NODE(5)

BUS	UTIL.	#MSGs	AVG-QLEN	MESSAGES
B1	5%	4	0.00	
B2	4%	2	0.00	
B3	5%	2	0.00	
B4	5%	2	0.00	
B5	0%	1	0.00	FIND_SOLUTION
B6	1%	0	0.00	

```

>>>      2.500 INTERFACE(1) : ENTER CSSA COMMAND -
==> ,RUN
+++      2.600 P6: NODE(4) RECEIVES FIND_SOLUTION(INT:4,INT:
              0,INT:0,INT:0) FROM NQUEENS(1)
+++      2.600 P6: NODE(4) STARTING OPERATION FIND_SOLUTION(INT:
              4,INT:0,INT:0,INT:0)
+++      2.700 P5: NODE(3) TERMINATED
+++      3.600 P7: NODE(5) RECEIVES FIND_SOLUTION(INT:3,INT:
              1,INT:0,INT:0) FROM NODE(3)
+++      3.600 P7: NODE(5) STARTING OPERATION FIND_SOLUTION(INT:
              3,INT:1,INT:0,INT:0)
+++      3.800 P6: NODE(4) CREATES NODE(8) (INT:4)
+++      3.800 P6: NODE(4) SENDS FIND_SOLUTION(INT:4,INT:1,INT:
              0,INT:0) TO NODE(8)
+++      3.800 P4: NODE(8) IS IDLE

```

&lt;==

```

+++ 3.800 P3: NODE(1) CREATES NODE(9) (INT:4)
+++ 3.800 P5: NODE(9) IS IDLE
+++ 4.100 P6: NODE(4) CREATES NODE(10) (INT:4)
+++ 4.100 P6: NODE(10) IS IDLE
+++ 4.100 P6: NODE(4) SENDS FIND_SOLUTION(INT:4,INT:2,INT:
      0,INT:0) TO NODE(10)
+++ 4.100 P6: NODE(10) RECEIVES FIND_SOLUTION(INT:4,INT:
      2,INT:0,INT:0) FROM NODE(4)
+++ 4.100 P6: NODE(10) STARTING OPERATION FIND_SOLUTION(INT:
      4,INT:2,INT:0,INT:0)
>>> 4.100 REAL-TIME LIMIT EXCEEDED
>>> 4.100 INTERFACE(1) : ENTER CSSA COMMAND -
==> ,RUN
+++ 4.200 P2: NODE(6) CREATES NODE(11) (INT:4)
+++ 4.200 P2: NODE(6) SENDS FIND_SOLUTION(INT:2,INT:4,INT:
      1,INT:0) TO NODE(11)
+++ 4.600 P2: NODE(6) TERMINATED
+++ 4.600 P7: NODE(11) IS IDLE
+++ 4.800 P3: NODE(7) TERMINATED
+++ 4.800 P3: NODE(1) SENDS FIND_SOLUTION(INT:1,INT:4,INT:
      0,INT:0) TO NODE(9)
+++ 4.900 P3: NODE(1) TERMINATED
+++ 5.200 P5: NODE(9) RECEIVES FIND_SOLUTION(INT:1,INT:
      4,INT:0,INT:0) FROM NODE(1)
+++ 5.200 P5: NODE(9) STARTING OPERATION FIND_SOLUTION(INT:
      1,INT:4,INT:0,INT:0)
+++ 5.300 P4: NODE(8) RECEIVES FIND_SOLUTION(INT:4,INT:
      1,INT:0,INT:0) FROM NODE(4)
+++ 5.300 P4: NODE(8) STARTING OPERATION FIND_SOLUTION(INT:
      4,INT:1,INT:0,INT:0)
+++ 5.400 P6: NODE(4) TERMINATED
+++ 5.600 P7: NODE(5) CREATES NODE(12) (INT:4)
+++ 5.600 P7: NODE(11) RECEIVES FIND_SOLUTION(INT:2,INT:
      4,INT:1,INT:0) FROM NODE(6)
+++ 5.600 P7: NODE(11) STARTING OPERATION FIND_SOLUTION(INT:
      2,INT:4,INT:1,INT:0)
+++ 5.600 P2: NODE(12) IS IDLE
>>> 5.600 REAL-TIME LIMIT EXCEEDED
>>> 5.600 INTERFACE(1) : ENTER CSSA COMMAND -
==> ,RUN
+++ 6.400 P6: NODE(10) TERMINATED
+++ 6.600 P7: NODE(5) SENDS FIND_SOLUTION(INT:3,INT:1,INT:
      4,INT:0) TO NODE(12)
+++ 7.000 P5: NODE(9) CREATES NODE(13) (INT:4)
+++ 7.000 P5: NODE(9) SENDS FIND_SOLUTION(INT:1,INT:4,INT:
      2,INT:0) TO NODE(13)
+++ 7.000 P3: NODE(13) IS IDLE
+++ 7.100 P2: NODE(12) RECEIVES FIND_SOLUTION(INT:3,INT:
      1,INT:4,INT:0) FROM NODE(5)
+++ 7.100 P2: NODE(12) STARTING OPERATION FIND_SOLUTION(INT:
      3,INT:1,INT:4,INT:0)
+++ 7.200 P4: NODE(8) CREATES NODE(14) (INT:4)
+++ 7.200 P4: NODE(14) IS IDLE
+++ 7.200 P4: NODE(8) SENDS FIND_SOLUTION(INT:4,INT:1,INT:
      3,INT:0) TO NODE(14)
+++ 7.200 P4: NODE(14) RECEIVES FIND_SOLUTION(INT:4,INT:
      1,INT:3,INT:0) FROM NODE(8)
+++ 7.200 P4: NODE(14) STARTING OPERATION FIND_SOLUTION(INT:
      4,INT:1,INT:3,INT:0)
>>> 7.200 REAL-TIME LIMIT EXCEEDED
>>> 7.200 INTERFACE(1) : ENTER CSSA COMMAND -
==> ,RUN
+++ 7.300 P5: NODE(9) TERMINATED
+++ 7.700 P7: NODE(5) TERMINATED
+++ 8.000 P7: NODE(11) CREATES NODE(15) (INT:4)
+++ 8.000 P7: NODE(11) SENDS FIND_SOLUTION(INT:2,INT:4,INT:
      1,INT:3) TO NODE(15)
+++ 8.000 P5: NODE(15) IS IDLE
+++ 8.100 P3: NODE(13) RECEIVES FIND_SOLUTION(INT:1,INT:
      4,INT:2,INT:0) FROM NODE(9)
+++ 8.100 P3: NODE(13) STARTING OPERATION FIND_SOLUTION(INT:
      1,INT:4,INT:2,INT:0)
+++ 8.200 P7: NODE(11) TERMINATED
+++ 8.400 P4: NODE(8) TERMINATED
+++ 9.300 P2: NODE(12) CREATES NODE(16) (INT:4)
+++ 9.300 P2: NODE(12) SENDS FIND_SOLUTION(INT:3,INT:1,INT:
      4,INT:2) TO NODE(16)

```

&lt;==

&lt;==

&lt;==

```

+++ 9.300 P6: NODE(16) IS IDLE
+++ 9.400 P5: NODE(15) RECEIVES FIND_SOLUTION(INT:2,INT:
      4,INT:1,INT:3) FROM NODE(11)
+++ 9.400 P5: NODE(15) STARTING OPERATION FIND_SOLUTION(INT:
      2,INT:4,INT:1,INT:3)
+++ 9.600 P2: NODE(12) TERMINATED
+++ 9.900 P4: NODE(14) TERMINATED
One solution for 4 queens is : 2,4,1,3"),REPLY TO: OK TO INTERFACE(1)
+++ 10.400 P5: NODE(15) IS IDLE
+++ 10.600 P3: NODE(13) TERMINATED
+++ 10.700 P6: NODE(16) RECEIVES FIND_SOLUTION(INT:3,INT:
      1,INT:4,INT:2) FROM NODE(12)
+++ 10.700 P6: NODE(16) STARTING OPERATION FIND_SOLUTION(INT:
      3,INT:1,INT:4,INT:2)
One solution for 4 queens is : 2,4,1,3"),REPLY TO: OK FROM NODE(15)
One solution for 4 queens is : 3,1,4,2"),REPLY TO: OK TO INTERFACE(1)
+++ 11.700 P6: NODE(16) IS IDLE
One solution for 4 queens is : 3,1,4,2"),REPLY TO: OK FROM NODE(16)
>>> 12.100 SYSTEM TERMINATED
>>> 12.100 INTERFACE(1) : ENTER CSSA COMMAND -
==> MAILBOX
MAILBOX OF INTERFACE(1) :

(1) ANSWER(STRING:"One solution for 4 queens is : 2,4,1,3"),REPLY TO: OK
(2) ANSWER(STRING:"One solution for 4 queens is : 3,1,4,2"),REPLY TO: OK

>>> 12.100 INTERFACE(1) : ENTER CSSA COMMAND -
==> REPLY 1
+++ 12.100 P1: INTERFACE(1) SENDS *REPLY() TO NODE(15)
>>> 12.100 INTERFACE(1) : ENTER CSSA COMMAND -
==> MAILBOX
MAILBOX OF INTERFACE(1) :

(1) ANSWER(STRING:"One solution for 4 queens is : 2,4,1,3"),REPLY TO: OK
(2) ANSWER(STRING:"One solution for 4 queens is : 3,1,4,2"),REPLY TO: OK

>>> 12.100 INTERFACE(1) : ENTER CSSA COMMAND -
==> REPLY 2
+++ 12.100 P1: INTERFACE(1) SENDS *REPLY() TO NODE(16)
>>> 12.100 INTERFACE(1) : ENTER CSSA COMMAND -
==> ,STATUS
+++ 12.100 ALL EXISTING AGENTS:

```

AGENT	FACET	OPERATION	MAILBOX
P1: INTERFACE(1) *			ANSWER ANSWER
P2: NQUEENS(1)	ONLY_FACE		
P5: NODE(15)	ONLY_FACE	FIND_SOLU	N
P6: NODE(16)	ONLY_FACE	FIND_SOLU	N

```

>>> 12.100 INTERFACE(1) : ENTER CSSA COMMAND -
==> ,RUN
+++ 12.300 P5: NODE(15) RECEIVES *REPLY() FROM INTERFACE(1)
+++ 12.300 P6: NODE(16) RECEIVES *REPLY() FROM INTERFACE(1)
+++ 12.400 P5: NODE(15) TERMINATED
+++ 12.400 P6: NODE(16) TERMINATED
>>> 12.400 SYSTEM TERMINATED
>>> 12.400 INTERFACE(1) : ENTER CSSA COMMAND -
==> ,STATUS
+++ 12.400 ALL EXISTING AGENTS:

```

AGENT	FACET	OPERATION	MAILBOX
P1: INTERFACE(1) *			ANSWER ANSWER
P2: NQUEENS(1)	ONLY_FACE		

```

>>> 12.400 INTERFACE(1) : ENTER CSSA COMMAND -
==> SEND FIND_SOLUTION(2) TO NQ;RUN
+++ 12.400 P1: INTERFACE(1) SENDS FIND_SOLUTION(INT:2) TO NQUEENS(1)
+++ 12.500 P2: NQUEENS(1) RECEIVES FIND_SOLUTION(INT:2) FROM INTERFACE(1)
+++ 12.500 P2: NQUEENS(1) STARTING OPERATION FIND_SOLUTION(INT:
      2)
+++ 12.900 P2: NQUEENS(1) CREATES NODE(17)(INT:2)
+++ 12.900 P2: NQUEENS(1) SENDS FIND_SOLUTION(INT:1,INT:
      0) TO NODE(17)
+++ 12.900 P7: NODE(17) IS IDLE
+++ 13.100 P2: NQUEENS(1) CREATES NODE(18)(INT:2)
+++ 13.100 P2: NODE(18) IS IDLE
+++ 13.100 P2: NQUEENS(1) SENDS FIND_SOLUTION(INT:2,INT:
      0) TO NODE(18)
+++ 13.100 P2: NODE(18) RECEIVES FIND_SOLUTION(INT:2,INT:
      0) FROM NQUEENS(1)
+++ 13.100 P2: NODE(18) STARTING OPERATION FIND_SOLUTION(INT:
      2,INT:0)
+++ 13.600 P7: NODE(17) RECEIVES FIND_SOLUTION(INT:1,INT:
      0) FROM NQUEENS(1)
+++ 13.600 P7: NODE(17) STARTING OPERATION FIND_SOLUTION(INT:
      1,INT:0)
+++ 14.100 P2: NQUEENS(1) IS IDLE
+++ 14.300 P2: NODE(18) TERMINATED
+++ 14.800 P7: NODE(17) TERMINATED
>>> 14.800 SYSTEM TERMINATED
>>> 14.800 INTERFACE(1) : ENTER CSSA COMMAND -
==> ,STATUS
+++ 14.800 ALL EXISTING AGENTS:

```

AGENT	FACET	OPERATION	MAILBOX
P1: INTERFACE(1) *			ANSWER ANSWER
P2: NQUEENS(1)	ONLY_FACE		

```

>>> 14.800 INTERFACE(1) : ENTER CSSA COMMAND -
==> MAILBOX
MAILBOX OF INTERFACE(1) :

```

```

(1) ANSWER(STRING:"One solution for 4 queens is : 2,4,1,3"),REPLY TO: OK
(2) ANSWER(STRING:"One solution for 4 queens is : 3,1,4,2"),REPLY TO: OK

```

```

>>> 14.800 INTERFACE(1) : ENTER CSSA COMMAND -
==> TERMINATE
+++ 14.800 ALL EXISTING AGENTS:

```

AGENT	FACET	OPERATION	MAILBOX
P1: INTERFACE(1) *			ANSWER ANSWER
P2: NQUEENS(1)	ONLY_FACE		

#### CSSA-SESSION-STATISTICS

```

=====
SESSION STARTED AT 11:42:59.00
SESSION TERMINATED AT 11:49:31.00 ON 1982/01/18
REAL-TIME USED : 392.00 SEC.
CPU-TIME USED : 6.34 SEC.
SIMULATION TIME USED : 14.8000 SEC.
NUMBER OF AGENTS CREATED : 19
NUMBER OF MESSAGES SENT : 24

```

	1	2	3	4	5	6	7	BLOCKNESTING
1	type START_NETWORK is							
2	script(int : UP1, GEN_ZUF1, UP2, GEN_ZUF2;							+1
3	bool : SHOW1, SHOW2) assert UP1 > 0 and UP2 > 0							
4								
5	type TRANSMITTER is							
6	script(agent : CH; int : WINDOW; agent : H)							+2
7								
8	procedure ACTIVATE_TRANSMIT(int:N) is							+3
9	var int: I;							
10	oper: SEND_PACKET;							
11								
12	loop for I in 1..N do							+4
13	send SEND_PACKET to self;							
14	endloop;							-4
15	endprocedure;							-3
16								+2
17	var string: MESS:="",							
18	var int: LENGTH_OF_MESS, CURRENT_MESS_ID := 0;							
19								
20	facethead TRANSMISSION1;							
21								
22	facet START_TRANSMISSION is							+5
23	public: SEND_MESS;							
24								
25	operation SEND_MESS(string: M) is							+6
26	functionhead LENGTH (string: P1) returns int external;							
27	oper: NEXT_MESS;							
28								
29	CURRENT_MESS_ID := CURRENT_MESS_ID + 1;							
30	MESS := M;							
31	LENGTH_OF_MESS := LENGTH(MESS);							
32	if LENGTH_OF_MESS < WINDOW then							+7
33	call ACTIVATE_TRANSMIT(LENGTH_OF_MESS);							
34	else call ACTIVATE_TRANSMIT(WINDOW);							
35	endif;							-7
36	send NEXT_MESS to H;							+6
37	replace by TRANSMISSION1;							
38	endoperation							-6
39								+5
40	endfacet							-5
41								+2
42	facet TRANSMISSION1 is							+8
43	public: ACKNOWLEDGE;							
44	private: SEND_PACKET;							
45								
46	functionhead MOD (int: P1,P2) returns int external;							
47	functionhead SUBSTR (string:P1,int:P2,P3) returns string external;							
48	functionhead GENSTRING(int:P1) returns string external;							
49	type PACKET is record int:CHAR_NO;--> string:CHAR,int: MESS_ID;							+9
50	bool:ENDE; endrecord;							-9
51	var PACKET: P;							+8
52	relation WINDOW_PACKETS of PACKET;							
53	var int: S_NO,LWE := 1;							
54	var bool: IDLE_ENABLED := false;							
55	oper: COLLECT;							
56								

fig. 19 : Window Mechanism Script

```

57 facet TRANSMISSION2 is
58   include : ACKNOWLEDGE,SEND_PACKET,
59
60   operation idle is
61     find WINDOW_PACKETS(LWE) --> P do
62       | send COLLECT(P) to CH;
63     endfind;
64     IDLE_ENABLED := false;
65     leave;
66   endoperation
67
68 endfacet
69
70
71 operation SEND_PACKET is
72
73   P.CHAR := SUBSTR(MESS,S_NO,S_NO);
74   P.CHAR_NO := S_NO;
75   P.MESS_ID := CURRENT_MESS_ID;
76   if S_NO = LENGTH_OF_MESS then
77     | P.ENDE := true;
78   else P.ENDE := false;
79   endif;
80   insert P into WINDOW_PACKETS;
81   send COLLECT(P) to CH;
82   S_NO := S_NO + 1;
83   if S_NO = 2 then
84     | IDLE_ENABLED := true;
85     setup TRANSMISSION2;
86   endif;
87
88 endoperation
89
90
91 operation ACKNOWLEDGE(int:N,MESS_ID) is
92   if MESS_ID = CURRENT_MESS_ID then
93     if N = LENGTH_OF_MESS + 1 then
94       | replace by START_TRANSMISSION;
95     else
96       if LWE < N and N <= LWE + WINDOW then
97         if N + WINDOW - 1 <= LENGTH_OF_MESS then
98           | call ACTIVATE_TRANSMIT(N - LWE);
99         else call ACTIVATE_TRANSMIT(LENGTH_OF_MESS -
100           (LWE + WINDOW - 1));
101       endif;
102
103       LWE := N;
104       if not IDLE_ENABLED then
105         | IDLE_ENABLED := true;
106         setup TRANSMISSION2;
107       endif;
108     else
109       if LWE = N then
110         | print("N= " + GENSTRING(N) + " ignored");
111       else print("N = " + GENSTRING(N) + " out of range");
112       endif;
113     endif;
114   endif;
115
116 endoperation

```

fig. 19 -continued-

117				*8
118		endfacet		-8
119				*2
120		initial START_TRANSMISSIO		
121				
122		endscript;		-2
123				*1
124				
125				
126		type COLLECTOR is		
127		script(agent : H; int : WINDOW)		+23
128				
129		var int:CURRENT_MESS_ID:= 1;		
130		var int:LWE_OLD;		
131		var agent: CH;		
132				
133		facethead COLLECTING;		
134				
135		facet STARTING is		+24
136		public: STARTC;		
137				
138		operation STARTC(->CH) is		+25
139		replace by COLLECTING;		
140		endoperation		-25
141		endfacet		-24
142				+23
143		facet COLLECTING is		+26
144		public: COLLECT;		
145				
146		var bool: FIRST_TIME:=true;		
147		var int: I;		
148		var string: MESS:= "";		
149		type RECEIVED is record int: S_NO;-->		+27
150		bool: REC; string: CHAR; endrecord;		-27
151		var RECEIVED : RC;		+26
152		relation R of RECEIVED;		
153		var int: LWE := 1;		
154		var int: LAST;		
155				
156		operation COLLECT(int:CHAR_NO;string:CHAR;int:MESS_ID;bool:ENDE) is		+28
157		port: ACCEPTED;		
158		oper: RECEIVE_MESS,ACKNOWLEDGE;		
159		-----		
160		if FIRST_TIME then		+29
161		loop for I in 1..WINDOW do		+30
162		RC.S_NO := I;		
163		RC.REC := false;		
164		RC.CHAR := "";		
165		insert RC into R;		
166		endloop;		-30
167		FIRST_TIME := false;		+29
168		endif;		-29
169				+28
170		if MESS_ID = CURRENT_MESS_ID then		+31
171		find R(CHAR_NO - LWE + 1) --> RC do		+32
172				
173		RC.CHAR := CHAR;		
174		if not (CHAR_NO = LWE) then		+33
175		RC.REC := true;		
176		insert RC into R;		

fig. 19 -continued-



177	if ENDE then LAST := CHAR_NO; endif;	+34-34
178	send ACKNOWLEDGE(LWE,CURRENT_MESS_ID) to CH;	*33
179	else (* CHAR_NO = LWE *)	+35
180	find R(1) --> RC1 do	+36
181	if ENDE then	
182	send ACKNOWLEDGE(LWE + 1,CURRENT_MESS_ID) to CH;	
183	MESS := MESS + RC1.CHAR;	
184	send RECEIVE_MESS(MESS) to H reply to ACCEPTED;	
185	CURRENT_MESS_ID := CURRENT_MESS_ID + 1;	
186	wait ACCEPTED;	
187	LWE_OLD := LWE + 1;	
188	replace by COLLECTING;	
189	(* EFFECT: new initialization *)	
190	(* of local variables *)	
191	else begin	+37
192	var int: J;	
193	MESS := MESS + RC1.CHAR;	
194	loop for J1 in 2..WINDOW until EXITLOOP do	+38
195	find R(J1) --> RC do	+39
196	if RC.REC then	+40
197	MESS := MESS + RC.CHAR;	
198	else signal EXITLOOP;	
199	endif;	
200	endfind;	-40
201	if J1 = WINDOW then J := WINDOW + 1;endif;	-39
202	exit EXITLOOP is	+41-41
203	J := J1;	*38
204	endloop;	-38
205	LWE := LWE + J - 1;	*37
206	send ACKNOWLEDGE(LWE,CURRENT_MESS_ID) to CH;	
207	if LWE = LAST + 1 then	+42
208	send RECEIVE_MESS(MESS) to H	
209	reply to ACCEPTED;	
210	CURRENT_MESS_ID := CURRENT_MESS_ID + 1;	
211	wait ACCEPTED;	
212	LWE_OLD := LWE;	
213	replace by COLLECTING;	
214	else begin	+43
215	I := 2;	
216	loop while I + J - 1 <= WINDOW do	+44
217	find R(I) --> RC2 do	+45
218	find R(I + J - 1) --> RC3 do	+46
219	delete RC2 in R;	
220	RC2.REC := RC3.REC;	
221	RC2.CHAR := RC3.CHAR;	
222	I := I + 1;	
223	insert RC2 into R;	
224	endfind;	-46
225	endfind;	-45
226	endloop;	-44
227	loop while I <= WINDOW do	+47
228	find R(I) --> RC2 do	+48
229	delete RC2 in R;	
230	RC2.REC := false;	
231	insert RC2 into R;	
232	endfind;	-48
233	I := I + 1;	*47
234	endloop;	-47

fig. 19 -continued-

```

237 |         end;                                     | -43
238 |     endif;                                       | -42
239 | and;                                           | -37
240 |     endif;                                       | -36
241 |     endif;                                       | -35
242 | endif;                                           | -33
243 | otherwise do send ACKNOWLEDGE(LWE,CURRENT_MESS_ID) to CH; | *32
244 |     endif;                                       | -32
245 | else send ACKNOWLEDGE(LWE_OLD,CURRENT_MESS_ID - 1) to CH; | *31
246 |     endif;                                       | -31
247 | endoperation                                    | -28
248 |                                                 | *26
249 | endfacet                                        | -26
250 |                                                 | *23
251 | initial STARTING;                               |
252 |                                                 |
253 | endsript;                                       | -23
254 |                                                 | *1
255 |                                                 |
256 | type HOST is                                   |
257 | script                                         | +49
258 |                                                 |
259 | var agent : TRM;                              |
260 |                                                 |
261 | facethead WORKING;                             |
262 |                                                 |
263 | facet INITH is                                | +50
264 | |                                               |
265 | | public : STARTH;                            |
266 | |                                               |
267 | | operation STARTH(->TRM) is                  | +51
268 | | | replace by WORKING;                      |
269 | | endoperation                                | -51
270 | |                                             | *50
271 | endfacet                                      | -50
272 |                                                 | *49
273 |                                                 |
274 | facet WORKING is                              | +52
275 | |                                               |
276 | | public : TRANSMIT,RECEIVE_MESS,NEXT_MESS;   |
277 | | private : TWIST_AND_SEND;                   |
278 | |                                               |
279 | | var int : I,J;                             |
280 | | var bool : READY_TO_TRM,READY_TO_TWIST := true; |
281 | | var string : M;                           |
282 | | type REC_ELEM is record int : IX; --> string : MESS; endrecord; | +53-53
283 | | relation REC_MESS of REC_ELEM;              | *52
284 | | var REC_ELEM : RE;                         |
285 | |                                               |
286 | | functionhead LENGTH( string : P1 ) returns int external; |
287 | |                                               |
288 | | operation TRANSMIT(->M) assert READY_TO_TRM is | +54
289 | | | if M = "" then                          | +55
290 | | | | print(" <empty message> cannot be transmitted"); |
291 | | | else                                     |
292 | | | | I := 0;                               |
293 | | | | READY_TO_TRM := false;                |
294 | | | | send TWIST_AND_SEND to self;           |
295 | | | endif;                                  | -55
296 | | endoperation;                             | -54

```

fig. 19 -continued-

297			+52
298	operation NEXT_MESS is		+56
299	READY_TO_TWIST := true;		
300	endoperation;		-56
301			+52
302	operation TWIST_AND_SEND assert READY_TO_TWIST is		+57
303	oper : SEND_MESS;		
304	functionhead SUBSTR(string:P1,int:P2,int:P3)		
305	returns string external;		
306	-----		
307	READY_TO_TWIST := false;		
308	I := I + 1;		
309	send SEND_MESS(M) to TRM;		
310			
311	if I < LENGTH(M) then		+58
312	M := SUBSTR(M,2,LENGTH(M)) + SUBSTR(M,1,1);		
313	send TWIST_AND_SEND to self;		
314	else READY_TO_TRM := true;		
315	endif;		-58
316	endoperation;		-57
317			+52
318	operation RECEIVE_MESS(string : MESS) is		+59
319	reply;		
320	J := J + 1;		
321	RE.IX := J;		
322	RE.MESS := MESS;		
323	insert RE into REC_MESS;		
324			
325	if J >= LENGTH(MESS) then		+60
326	begin port : OK;		+61
327	oper : ACKN;		
328	functionhead GENSTRING(int : P1) returns string external;		
329	var string : OUT;		
330	-----		
331	send ACKN("Host received fully twisted message")		
332	to interface reply to OK;		
333	loop for J in 1..LENGTH(MESS) do		+62
334	find REC_MESS(J) -> MS do		+63
335	OUT := GENSTRING(J) + ". message : " + MS.MESS;		
336	send ACKN(OUT) to interface;		
337	endif;		-63
338	endloop;		-62
339	J := 0;		+61
340	wait OK;		
341	end;		-61
342	endif;		-60
343	endoperation;		-59
344			+52
345	endfacet		-52
346			+49
347	initial INITH;		
348			
349	endscript;		-49
350			+1
351			
352			
353	type FAULTY_CHANNEL is		
354	script(agent : C)		+64
355			
356	var agent : T;		

fig. 19 -continued-

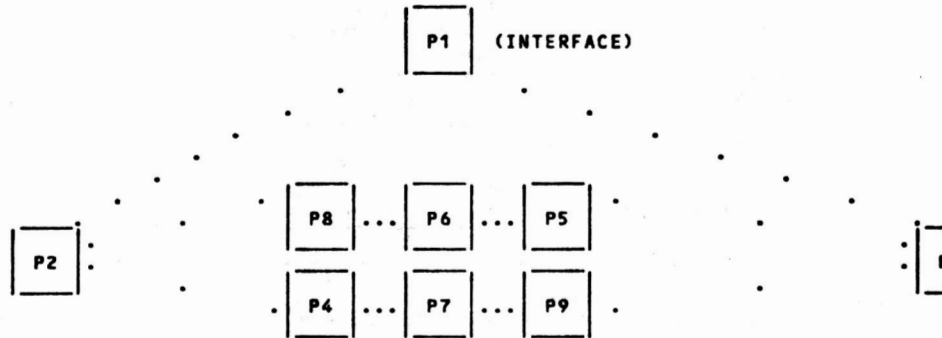
<pre> 357 var int : UP, GEN_ZUF, ZUF, I; 358 var bool : SHOW; 359 360 functionhead RANDOM(int:LOW,UP) returns int external; 361 362 facethead TRANSMITTING; 363 364 facet INITIALIZATION is 365   public : STARTCH; 366 367   operation STARTCH(-&gt;T,-&gt;UP,-&gt;GEN_ZUF,-&gt;SHOW) is 368     if UP &gt; 10 then UP := 10; endif; 369     loop for I in 1..GEN_ZUF do 370       ZUF := RANDOM(1,UP); 371     endloop; 372     replace by TRANSMITTING; 373   endoperation 374 375 endfacet 376 377 facet TRANSMITTING is 378   public : COLLECT, ACKNOWLEDGE, SET_CHNL_PARMS; 379 380   functionhead GENSTRING(int:P1) returns string external; 381   var int : COUNT := 1; 382   var string : MESS; 383 384   operation COLLECT(int:P1; string:P2; int: P3; bool:P4) is 385 386     if SHOW then 387       MESS := "(" + GENSTRING(P1) + " , " + P2 + " , " 388             + GENSTRING(P3) + " , " ; 389       if P4 then MESS := MESS + "T" ; 390       else MESS := MESS + "F" ; 391     endif; 392   endif; 393 394   if COUNT = 1 then ZUF := RANDOM(1,UP); endif; 395   if COUNT &lt; ZUF then 396     send COLLECT(P1,P2,P3,P4) to C; 397     COUNT := COUNT + 1; 398     if SHOW then print("sends COLLECT" + MESS); endif; 399   else COUNT := 1; 400     if SHOW then 401       print("forgets COLLECT" + MESS); 402     endif; 403   endif; 404 405 endoperation 406 407 operation ACKNOWLEDGE(int:P1,P2) is 408 409   if SHOW then 410     MESS := "(" + GENSTRING(P1) + " , " + GENSTRING(P2) + " )"; 411   endif; 412 413   if COUNT = 1 then ZUF := RANDOM(1,UP); endif; 414 415   if COUNT &lt; ZUF then 416     send ACKNOWLEDGE(P1,P2) to T; </pre>	<pre> *64 +65 +66 +67-67 +68 -68 *66 -66 *65 -65 *64 +69 +70 +71 +72 -72 -71 *70 +73-73 +74 +75-75 *74 +76 -76 -74 *70 -70 *69 +77 +78 -78 *77 +79-79 *77 +80 </pre>
--	--

fig. 19 -continued-

417	COUNT := COUNT + 1;	*80
418	if SHOW then print("sends ACKNOWLEDGE" + MESS); endif;	+81-81
419	else COUNT := 1;	*80
420	if SHOW then	+82
421	print("forgets ACKNOWLEDGE" + MESS);	
422	endif;	-82
423	endif;	-80
424		*77
425	endoperation	-77
426		*69
427	operation SET_CHNL_PARMS(-> UP, GEN_ZUF, SHOW) assert UP > 0 is	+83
428		
429	if UP > 10 then UP := 10; endif;	+84-84
430	loop for I in 1..GEN_ZUF do	+85
431	ZUF := RANDOM(1,UP);	
432	endloop;	-85
433	COUNT := 1;	*83
434		
435	endoperation	-83
436		*69
437	endfacet	-69
438		*64
439	initial INITIALIZATION;	
440		
441		
442	endscript;	-64
443		*1
444		
445		
446	facet ONLY_FCT is	+86
447	public: START_SYSTEM;	
448	operation START_SYSTEM(int:WINDOW_SIZE) assert WINDOW_SIZE > 0 is	+87
449		
450	const HOST : H1 := new HOST;	
451	const HOST : H2 := new HOST;	
452		
453	const COLLECTOR : C2 := new COLLECTOR(H1, WINDOW_SIZE);	
454	const COLLECTOR : C1 := new COLLECTOR(H2, WINDOW_SIZE);	
455		
456	const FAULTY_CHANNEL : CH1 := new FAULTY_CHANNEL(C1);	
457	const FAULTY_CHANNEL : CH2 := new FAULTY_CHANNEL(C2);	
458		
459	const TRANSMITTER : T1 := new TRANSMITTER(CH1, WINDOW_SIZE, H1);	
460	const TRANSMITTER : T2 := new TRANSMITTER(CH2, WINDOW_SIZE, H2);	
461		
462	oper: STARTCH, STARTH, STARTC, ACQ_TO_HOST;	
463		
464	-----	
465		
466	send STARTCH(T1, UP1, GEN_ZUF1, SHOW1) to CH1;	
467	send STARTCH(T2, UP2, GEN_ZUF2, SHOW2) to CH2;	
468		
469	send STARTH(T1) to H1;	
470	send STARTH(T2) to H2;	
471		
472	send STARTC(CH1) to C1;	
473	send STARTC(CH2) to C2;	
474		
475	send ACQ_TO_HOST(H1, H2, CH1, CH2) to interface;	
476		
477		*87
478	terminate;	
479		
480	endoperation	-87
481		*86
482	endfacet	-86
483		*1
484	initial ONLY_FCT;	
485		
486	endscript	-1

BMS-CSSA-COMPILER - DATE OF RELEASE: 30 SEP 1981 NO ERROR DETECTED  
 END OF COMPILING ON 1982/01/21 AT 14:28:44.00 RETURN CODE = 0  
 COMPILE-TIME (CPU) = 28.34 SEC. EXECUTION-TIME = 177.00 SEC.  
 NUMBER OF SOURCE-LINES READ = 486 NUMBER OF TOKENS = 2171  
 NUMBER OF OBJECT-RECORDS GENERATED = 5254

-----  
C S S A - S I M U L A T I O N - S Y S T E M  
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PROGRAM GENERATED ON 1981/11/17 AT 12:27:21.00  
BY BMS-CSSA-COMPILER (VERS. 30 SEP 1981)

PROTOCOL OF CSSA SESSION ON 1981/12/04 AT 12:08:57.00  
-----

>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -  
==> ,DISPLAY

<==

IDENTIFIER	TYPE	VALUE
COLLECTOR	SCRIPT	COLLECTOR
FAULTY_CHANNEL	SCRIPT	FAULTY_CHANNEL
HOST	SCRIPT	HOST
START_NETWORK	SCRIPT	START_NETWORK
TRANSMITTER	SCRIPT	TRANSMITTER

>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -  
==> VAR AGENT : NTW := NEW START\_NETWORK(600,10,500,0,TRUE,TRUE)  
>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -  
==> VAR AGENT : H1,H2,CH1,CH2; OPER : START\_SYSTEM,TRANSMIT,SET\_CHNL\_PARMS  
>>> 0.000 INTERFACE(1) : ENTER CSSA COMMAND -  
==> SEND START\_SYSTEM(2) TO NTW; RUN; MAILBOX  
>>> 10.300 SYSTEM TERMINATED  
MAILBOX OF INTERFACE(1) :

<==

<==

<==

(1) ACQ\_TO\_HOST(AGENT:HOST(1),AGENT:HOST(2),AGENT:FAULTY\_CHANNEL(1)  
,AGENT:FAULTY\_CHANNEL(2))

>>> 10.300 INTERFACE(1) : ENTER CSSA COMMAND -  
==> RECEIVE 1 (H1,H2,CH1,CH2); STATUS; DISPLAY  
+++ 10.300 ALL EXISTING AGENTS:

<==

AGENT	FACET	OPERATION	MAILBOX
P1: INTERFACE(1) *			ACQ_TO_HOST
P3: HOST(1)	WORKING		
P4: HOST(2)	WORKING		
P5: COLLECTOR(1)	COLLECTIN		
P6: COLLECTOR(2)	COLLECTIN		
P7: FAULTY_CHANNE	TRANSMITT		
P8: FAULTY_CHANNE	TRANSMITT		
P9: TRANSMITTER(1)	START_TRA		
P2: TRANSMITTER(2)	START_TRA		

fig. 20 : Window Mechanism Execution Protocol

IDENTIFIER	TYPE	VALUE
CH1	AGENT	FAULTY_CHANNEL(1)
CH2	AGENT	FAULTY_CHANNEL(2)
COLLECTOR	SCRIPT	COLLECTOR
FAULTY_CHANNEL	SCRIPT	FAULTY_CHANNEL
HOST	SCRIPT	HOST
H1	AGENT	HOST(1)
H2	AGENT	HOST(2)
NTW	AGENT	START_NETWORK(1)
SET_CHNL_PARMS	LITERAL	
START_NETWORK	SCRIPT	START_NETWORK
START_SYSTEM	LITERAL	
TRANSMIT	LITERAL	
TRANSMITTER	SCRIPT	TRANSMITTER

```

>>> 10.300 INTERFACE(1) : ENTER CSSA COMMAND -
==> SEND TRANSMIT("ABCD") TO H1; SEND TRANSMIT("XYZ") TO H2 ;RUN
*** 16.700 FAULTY_CHANNEL(2) : forgets COLLECT(1 , X , 1 , F)
*** 16.800 FAULTY_CHANNEL(1) : sends COLLECT(1 , A , 1 , F)
*** 17.800 FAULTY_CHANNEL(1) : sends COLLECT(2 , B , 1 , F)
*** 17.900 FAULTY_CHANNEL(2) : sends COLLECT(2 , Y , 1 , F)
*** 18.800 FAULTY_CHANNEL(1) : sends COLLECT(1 , A , 1 , F)
*** 18.900 FAULTY_CHANNEL(2) : sends COLLECT(1 , X , 1 , F)
*** 21.900 FAULTY_CHANNEL(1) : sends ACKNOWLEDGE(2,1)
*** 22.300 FAULTY_CHANNEL(2) : sends ACKNOWLEDGE(1,1)
*** 23.500 TRANSMITTER(2) : N= 1 ignored
*** 24.000 FAULTY_CHANNEL(2) : sends ACKNOWLEDGE(3,1)
*** 24.500 FAULTY_CHANNEL(1) : sends ACKNOWLEDGE(3,1)
*** 25.900 FAULTY_CHANNEL(1) : sends ACKNOWLEDGE(3,1)
*** 27.100 FAULTY_CHANNEL(1) : forgets COLLECT(3 , C , 1 , F)
*** 28.300 FAULTY_CHANNEL(1) : sends COLLECT(2 , B , 1 , F)
*** 28.500 FAULTY_CHANNEL(2) : sends COLLECT(3 , Z , 1 , T)
*** 29.300 FAULTY_CHANNEL(1) : sends COLLECT(4 , D , 1 , T)
*** 29.500 FAULTY_CHANNEL(2) : sends COLLECT(3 , Z , 1 , T)
*** 30.300 FAULTY_CHANNEL(1) : sends COLLECT(3 , C , 1 , F)
*** 30.700 TRANSMITTER(1) : N= 3 ignored
*** 32.000 FAULTY_CHANNEL(2) : sends ACKNOWLEDGE(4,1)
*** 32.800 FAULTY_CHANNEL(1) : sends ACKNOWLEDGE(3,1)
*** 33.800 FAULTY_CHANNEL(1) : sends ACKNOWLEDGE(3,1)
*** 34.800 FAULTY_CHANNEL(1) : sends ACKNOWLEDGE(5,1)
*** 34.800 TRANSMITTER(1) : N= 3 ignored
*** 35.800 TRANSMITTER(1) : N= 3 ignored
*** 36.600 FAULTY_CHANNEL(2) : forgets COLLECT(1 , Y , 2 , F)
*** 37.900 FAULTY_CHANNEL(2) : sends COLLECT(2 , Z , 2 , F)
*** 39.800 FAULTY_CHANNEL(2) : sends COLLECT(1 , Y , 2 , F)
*** 40.300 FAULTY_CHANNEL(1) : sends COLLECT(1 , B , 2 , F)
*** 41.400 FAULTY_CHANNEL(1) : sends COLLECT(2 , C , 2 , F)
*** 42.300 FAULTY_CHANNEL(1) : forgets COLLECT(1 , B , 2 , F)
*** 44.900 FAULTY_CHANNEL(2) : sends ACKNOWLEDGE(4,1)
*** 45.900 FAULTY_CHANNEL(2) : sends ACKNOWLEDGE(1,2)
*** 47.100 TRANSMITTER(2) : N= 1 ignored
*** 47.500 FAULTY_CHANNEL(2) : sends ACKNOWLEDGE(3,2)
*** 48.300 FAULTY_CHANNEL(1) : sends ACKNOWLEDGE(2,2)
*** 50.800 FAULTY_CHANNEL(1) : sends ACKNOWLEDGE(3,2)
*** 52.000 FAULTY_CHANNEL(2) : sends COLLECT(3 , X , 2 , T)
*** 53.000 FAULTY_CHANNEL(2) : sends COLLECT(3 , X , 2 , T)
*** 53.500 FAULTY_CHANNEL(1) : sends COLLECT(2 , C , 2 , F)
*** 54.500 FAULTY_CHANNEL(1) : sends COLLECT(3 , D , 2 , F)
*** 55.500 FAULTY_CHANNEL(1) : sends COLLECT(4 , A , 2 , T)
*** 55.500 FAULTY_CHANNEL(2) : sends ACKNOWLEDGE(4,2)
*** 56.400 FAULTY_CHANNEL(1) : forgets COLLECT(3 , D , 2 , F)
*** 57.300 FAULTY_CHANNEL(1) : sends ACKNOWLEDGE(3,2)
*** 58.500 TRANSMITTER(1) : N= 3 ignored
*** 59.000 FAULTY_CHANNEL(1) : sends ACKNOWLEDGE(4,2)
*** 60.100 FAULTY_CHANNEL(2) : forgets COLLECT(1 , Z , 3 , F)
*** 61.000 FAULTY_CHANNEL(1) : sends ACKNOWLEDGE(5,2)
*** 61.300 FAULTY_CHANNEL(2) : forgets COLLECT(2 , X , 3 , F)
*** 62.400 FAULTY_CHANNEL(2) : sends COLLECT(1 , Z , 3 , F)
*** 63.700 FAULTY_CHANNEL(1) : sends COLLECT(4 , A , 2 , T)
*** 65.700 FAULTY_CHANNEL(1) : sends COLLECT(1 , C , 3 , F)
*** 66.800 FAULTY_CHANNEL(1) : sends COLLECT(2 , D , 3 , F)
*** 67.700 FAULTY_CHANNEL(1) : forgets COLLECT(1 , C , 3 , F)

```

&lt;==

fig. 20 -continued-

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*** 68.400 FAULTY_CHANNEL(2) : sends ACKNOWLEDGE(4,2)
*** 70.000 FAULTY_CHANNEL(2) : sends ACKNOWLEDGE(2,3)
*** 72.200 FAULTY_CHANNEL(1) : sends ACKNOWLEDGE(5,2)
*** 73.700 FAULTY_CHANNEL(1) : sends ACKNOWLEDGE(2,3)
*** 74.400 FAULTY_CHANNEL(2) : forgets COLLECT(3 , Y , 3 , T)
*** 75.600 FAULTY_CHANNEL(2) : sends COLLECT(2 , X , 3 , F)
*** 76.300 FAULTY_CHANNEL(1) : sends ACKNOWLEDGE(3,3)
*** 79.000 FAULTY_CHANNEL(1) : sends COLLECT(2 , D , 3 , F)
*** 79.600 FAULTY_CHANNEL(2) : forgets ACKNOWLEDGE(3,3)
*** 80.000 FAULTY_CHANNEL(1) : sends COLLECT(3 , A , 3 , F)
*** 80.900 FAULTY_CHANNEL(1) : forgets COLLECT(4 , B , 3 , T)
*** 82.100 FAULTY_CHANNEL(1) : sends COLLECT(3 , A , 3 , F)
*** 82.900 FAULTY_CHANNEL(1) : sends ACKNOWLEDGE(3,3)
*** 84.100 TRANSMITTER(1) : N= 3 ignored
*** 84.600 FAULTY_CHANNEL(1) : sends ACKNOWLEDGE(4,3)
*** 85.600 FAULTY_CHANNEL(1) : sends ACKNOWLEDGE(4,3)
*** 87.200 TRANSMITTER(1) : N= 4 ignored
*** 88.500 FAULTY_CHANNEL(1) : sends COLLECT(4 , B , 3 , T)
*** 91.800 FAULTY_CHANNEL(1) : sends ACKNOWLEDGE(5,3)
*** 96.500 FAULTY_CHANNEL(1) : sends COLLECT(1 , D , 4 , F)
*** 97.500 FAULTY_CHANNEL(1) : forgets COLLECT(2 , A , 4 , F)
*** 99.700 FAULTY_CHANNEL(1) : sends COLLECT(1 , D , 4 , F)
*** 103.200 FAULTY_CHANNEL(1) : sends ACKNOWLEDGE(2,4)
*** 104.600 FAULTY_CHANNEL(1) : sends ACKNOWLEDGE(2,4)
*** 106.400 TRANSMITTER(1) : N= 2 ignored
*** 107.700 FAULTY_CHANNEL(1) : sends COLLECT(3 , B , 4 , F)
*** 108.700 FAULTY_CHANNEL(1) : sends COLLECT(2 , A , 4 , F)
*** 111.200 FAULTY_CHANNEL(1) : sends ACKNOWLEDGE(2,4)
*** 112.400 TRANSMITTER(1) : N= 2 ignored
*** 112.900 FAULTY_CHANNEL(1) : sends ACKNOWLEDGE(4,4)
*** 117.400 FAULTY_CHANNEL(1) : sends COLLECT(4 , C , 4 , T)
*** 118.300 FAULTY_CHANNEL(1) : forgets COLLECT(4 , C , 4 , T)
*** 120.800 FAULTY_CHANNEL(1) : sends ACKNOWLEDGE(5,4)
>>> 127.400 SYSTEM TERMINATED
>>> 127.400 INTERFACE(1) : ENTER CSSA COMMAND -
==> MAILBOX
MAILBOX OF INTERFACE(1) :

```

```

(1) ACQ_TO_HOST(AGENT:HOST(1),AGENT:HOST(2),AGENT:FAULTY_CHANNEL(1)
      ,AGENT:FAULTY_CHANNEL(2))
(2) ACKN(STRING:"Host received fully twisted message"),REPLY TO: OK
(3) ACKN(STRING:"3. message : CDAB")
(4) ACKN(STRING:"4. message : DABC")
(5) ACKN(STRING:"2. message : BCDA")
(6) ACKN(STRING:"1. message : ABCD")

```

```

>>> 127.400 INTERFACE(1) : ENTER CSSA COMMAND -
==> REPLY 2, SEND TRANSMIT("A") TO H1;RUN
*** 133.600 FAULTY_CHANNEL(1) : forgets COLLECT(1 , A , 5 , T)
*** 134.700 FAULTY_CHANNEL(1) : forgets COLLECT(1 , A , 5 , T)
>>> 134.700 SYSTEM TERMINATED
>>> 134.700 INTERFACE(1) : ENTER CSSA COMMAND -
==> TERMINATE
+++ 134.700 ALL EXISTING AGENTS:

```

AGENT	FACET	OPERATION	MAILBOX
P1: INTERFACE(1) *			ACQ_TO_HOST ACKN ACKN ACKN ACKN ACKN
P3: HOST(1)	WORKING		
P4: HOST(2)	WORKING		
P5: COLLECTOR(1)	COLLECTIN		
P6: COLLECTOR(2)	COLLECTIN		
P7: FAULTY_CHANNE	TRANSMITT		
P8: FAULTY_CHANNE	TRANSMITT		
P9: TRANSMITTER(1	TRANSMISS		
P2: TRANSMITTER(2	TRANSMISS		

#### CSSA-SESSION-STATISTICS

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=====
SESSION STARTED AT 12:08:27.00
SESSION TERMINATED AT 12:18:17.00 ON 1981/12/04
REAL-TIME USED : 592.00 SEC.
CPU-TIME USED : 12.27 SEC.
SIMULATION TIME USED : 134.7000 SEC.
NUMBER OF AGENTS CREATED : 9
NUMBER OF MESSAGES SENT : 224

```







