



Deutsches
Forschungszentrum
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**Research
Report**

RR-93-29

Representing Belief in Multi-Agent Worlds via Terminological Logics

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September 1993

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DFKI-RR-93-29

This work has been supported by a grant from The Federal Ministry for Research and Technology (FKZ IWT-9201).

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Representing Belief in Multi-Agent Worlds via Terminological Logics

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Abstract

In multi-agent systems a group of autonomous intelligent systems, called agents, acts and cooperates in a world in order to achieve certain goals. Such systems are in general assumed to have no central control structure and hence each agent can only perform actions that are based on his local knowledge and on his local beliefs. In the literature knowledge of agents is mostly represented under the view that knowledge is true belief. On the other hand, if agents are acting in a (real) world their knowledge often is obtained by perception and communication, and hence typically is not true. Thus, the use of belief—where agents may have false beliefs—seems more appropriate than the use of knowledge in multi-agent systems.

Terminological logics provide a well-investigated and decidable fragment of first-order logics that is much more expressive than propositional logic and well suited to describe a world agents are acting in. However, knowledge or belief of agents can only be represented in a very limited way. In this paper we investigate how terminological logics can be extended in such a way that belief of agents can be represented in an adequate manner. We therefore exemplarily extend the concept language \mathcal{ALC} by a modal operator \Box , which is indexed by agents. Thereby, $\Box_i\varphi$ represents the fact “agent i believes φ ”. This belief operator will be interpreted in terms of possible worlds using the well-known modal logic KD45.

This extended language \mathcal{ALC}_B provides a uniform formalism to describe both, a world agents are acting in and the beliefs agents have about this world and about their own and other agents' beliefs. Thus, it can be seen as a two-dimensional extension of \mathcal{ALC} which allows both, reasoning about objective facts that hold in the world and reasoning on the level of possible worlds. We will give sound and complete algorithms to check consistency of the represented beliefs and to decide whether an \mathcal{ALC}_B -sentence is logically entailed by the beliefs of agents. Hence, when acting in a world agents can use beliefs which are explicitly represented as well as implicit beliefs that are entailed by their knowledge base.

Contents

1	Introduction	3
2	Syntax and Semantics of \mathcal{ALC}_B	5
2.1	The Concept Language \mathcal{ALC}	5
2.2	The Extended Language \mathcal{ALC}_B	7
3	Testing Satisfiability of \mathcal{ALC}_B-formulas	9
3.1	The \mathcal{ALC}_B Frame Algorithm	9
4	Optimization and Computing \mathcal{ALC}_B-Inferences	20
4.1	Computing \mathcal{ALC} test sets	21
4.2	Restricted \mathcal{ALC} -TBoxes	28
4.3	Computing \mathcal{ALC}_B -Inferences	30
5	Conclusion	31

1 Introduction

Research on the field of multi-agent systems deals with the question how a group of autonomous intelligent systems, called agents, can cooperate in order to achieve certain goals (see, e.g., [6, 15]). As an example, a forwarding agent a and a shipping agent b may cooperate in order to carry out overseas transportation orders.

Although the tasks that multi-agent systems are required to perform are normally stated in terms of the global behavior of the system, the actions that an agent performs can depend only on his local knowledge and on his local beliefs. Thus, there is a close relationship between knowledge, belief, and action in multi-agent systems (see, e.g., [27, 28, 16]). Suppose, in the above example agent a wants to offer a price for carrying out some transportation order o_1 . If he believes that there is no other forwarding agent who also can carry out o_1 , he will most likely offer another price as in the case where he believes that there is a competitor for this order. And if he even *knows* that there is no competitor for this order, he perhaps offers an exorbitant price. In a recent paper [21] we investigated how knowledge of agents can be represented on the basis of terminological logics, whereby we used the classical view of knowledge as true belief. That means, an agent knows φ if he believes that φ holds and φ actual does hold. On the other hand, as pointed out in, say [23], the knowledge represented in a knowledge base typically is not true. Thus, the use of beliefs—where agents may have false beliefs—seems more appropriate than knowledge for formalizing the reasoning and deduction of a knowledge base. In the current paper we concern with the question how agents can be equipped with beliefs about the world they are acting in, about beliefs of other agents, and also about their own beliefs. Thereby, it should be taken into consideration that different agents may have different beliefs about the same notions. For example, forwarding agent a may believe that *company XY* is a rich company and a good client, while forwarding agent B believes that *company XY* is rich but not a good client. Using the language \mathcal{ALC}_B , which is presented in the next section, this can be formalized by

$$\begin{aligned} \Box_a(\text{company XY: rich-company} \sqcap \text{good-client}) \quad \text{and} \\ \Box_b(\text{company XY: rich-company} \sqcap \neg \text{good-client}) \end{aligned}$$

respectively, where $\Box_i\varphi$ is to be read as “agent i believes φ ”.

Since the work of Hintikka [18], modal logics have widely been accepted to be an adequate formalism for representing knowledge and belief of agents. The intuitive idea here is that besides the real world agents can imagine a number of other worlds (situations) to be possible. By imposing various conditions on this possibility relation, we can capture a number of interesting axioms. For example, if we require that the world that the agent finds himself in is always one of the worlds he considers possible (which amounts to saying that the possibility relation is reflexive), then it follows that the agent does not know false facts. When using a possibility relation which captures axioms of knowledge (belief) an agent is said to know (believe) a fact φ iff φ is true

in all worlds he thinks to be possible. For example, an agent knows (believes) that there exists a monster of Loch Ness if there is such a monster in all worlds he considers possible. To express the beliefs of an agent a in this approach a binary operator $BELIEF(a, \varphi)$ is used, where φ is a formula over some logical language \mathcal{L} . If we want to devise a formalism for representing the beliefs of agents we have to take two decisions. Firstly, we have to decide what the general properties of belief are we want this formalism to capture. Secondly, we have to choose a suitable logical representation language \mathcal{L} which allows to describe the beliefs of agents.

There are many approaches to determine axioms characterizing belief (see, e.g., [22, 28, 24, 25, 26, 12, 17]). We will use the following axiomatization which has been most commonly used in the literature. The first of these properties states that an agent does not believe false facts. That means, an agent cannot believe both a fact and its negation, though he can believe facts which actually do not hold in the world. Secondly, if an agent believes a fact then he believes that he believes it (positive introspection), and if he does not believe in a fact then he believes that he does not believe in it (negative introspection). From this it follows, e.g., that agents believe that their beliefs are true (weak reflexivity). Finally, the probably most important property is that agents can reason on the basis of their beliefs. For example, suppose agent a believes that each truck which is owned by John can be used to transport gasoline and he believes that John owns the truck *truck-1*. In this case, agent a must be able to conclude that John's truck *truck-1* can (probably) be used to transport gasoline, and thus may negotiate with John for a transportation order.

As logical language to describe belief of agents we will use a terminological logic. Terminological logics provide a well-investigated and decidable fragment of first-order logics that is much more expressive than propositional logic. They are based on the work of Brachman and Schmolze [9] and have been developed as a structured formalism to describe the relevant concepts of a problem domain and the interactions between these concepts. Starting with atomic concepts (unary predicates) and roles (binary predicates), one therefore defines complex concepts with the help of operators provided by a concept language, and interactions between (complex) concepts are expressed by a set of so-called terminological axioms. On the other hand, by so-called assertional axioms, objects can be associated with concepts and relationships between objects can be defined via roles. For example, we can use these logics to represent facts like "each truck which is owned by John can be used to transport gasoline" or "John owns *truck-1* which is a truck".

In the literature, a lot of concept languages have been considered (see, e.g., [8, 29, 3]). But they all have in common that they are only suitable for representing objective facts about the world, and knowledge or beliefs of agents can only be represented in a very limited way. Thus, we need an extended concept language which allows the representation of belief according to the above given (informal) axiomatization. Since the work of Schild [31] it is known that the concept language \mathcal{ALC} provides a

terminological logic which is a notational variant of the propositional modal logic $K_{(m)}$. However, it is not investigated there how to extend this logic to a two-dimensional logic which allows reasoning on both the objective level and the level of possible worlds. In order to combine both levels one has to define syntax and semantics of an extended language. Baader and Ohlbach [5] present a multi-dimensional extension of \mathcal{ALC} , where multi-modal operators can be used at all levels of the concept terms and they can be used to modify both concepts and roles. However, the underlying logic is simply the basic modal logic K , and it is not yet clear how to extend their approach in such a way that modal logics different from K can be handled. Moreover, they could not succeed in proving completeness of their satisfiability algorithm.

In this paper we will present a different extended language where (sequences of) modal operators are only allowed in front of terminological and assertional axioms. This language allows one to interpret the modal operators w.r.t. modal logics different from K , e.g., $S4$ (see [21]) or $KD45$ (in the present paper). This language, called \mathcal{ALC}_B , can be seen as a two-dimensional representation language with terminological and assertional axioms as primitives where each primitive may describe a part of the world and each agent can believe a set of such primitives to hold in the world. The modal operators, which are indexed with agents, are interpreted in terms of possible worlds in such a way that they satisfy the above axiomatization of belief, what amounts in using the modal logic $KD45$. Thus, the resulting language provides a uniform formalism to describe both, a world agents are acting in as well as the beliefs agents have about this world and about their own and other agents' beliefs. We will give sound and complete algorithms for deciding satisfiability of \mathcal{ALC}_B -formulas and for testing whether an \mathcal{ALC}_B -formula is entailed by a given set of \mathcal{ALC}_B -formulas. Hence, when acting in a world agents can use beliefs which are explicitly represented as well as implicit beliefs that are entailed by their knowledge base.

2 Syntax and Semantics of \mathcal{ALC}_B

In this section we will formally introduce the language \mathcal{ALC}_B which extends the concept language \mathcal{ALC} by a modal operator \Box_i for each agent. Syntax and semantics of \mathcal{ALC} and \mathcal{ALC}_B are given in Subsections 2.1 and 2.2, respectively.

2.1 The Concept Language \mathcal{ALC}

Terminological logics provide two formalisms to describe a problem domain: a terminological formalism to represent taxonomical knowledge by defining concepts, which can be seen as sets of objects, and an assertional formalism which can be used to describe concrete objects. Therefore, one starts with a set of *atomic concepts* (unary predicates) and a set of *roles* (binary predicates). In the concept language \mathcal{ALC} *concepts* are then

built up from atomic concepts, the *top concept* \top , the *bottom concept* \perp , and roles inductively by:

1. Each atomic concept, \top , and \perp are concepts.

2. If C and D are concepts and R is a role, then

(a) $C \sqcap D$ (*concept conjunction*),

(b) $C \sqcup D$ (*concept disjunction*),

(c) $\neg C$ (*concept negation*),

(d) $\forall R.C$ (*value restriction*), and

(e) $\exists R.C$ (*exists restriction*)

are concepts.

An interpretation I is a function over some non-empty domain Δ^I which maps each atomic concept C to a subset C^I of Δ^I , each role R to a subset R^I of $\Delta^I \times \Delta^I$, \top to Δ^I , and \perp to \emptyset . Furthermore, \sqcap is interpreted as set intersection, \sqcup as set union, and \neg as set complement w.r.t. Δ^I . The value and the exists restrictions are interpreted by

$$[\forall R.C]^I = \{d \in \Delta^I \mid \forall d' : (d, d') \in R^I \rightarrow d' \in C^I\}$$

$$[\exists R.C]^I = \{d \in \Delta^I \mid \exists d' : (d, d') \in R^I \wedge d' \in C^I\}$$

For example, if *man* and *truck* are atomic concepts and *owns* is a role we can define the concept of men who own a truck by $\text{man} \sqcap \exists \text{owns.truck}$.

The taxonomical knowledge of a problem domain can be defined by an *ALC-TBox* (*terminology*), which consists of a finite set of terminological axioms. A *terminological axiom* is of the form

- $C = D$ (concept equivalence) or
- $C \neq D$ (negated concept equivalence)

where C, D are concepts. An interpretation I satisfies $C = D$ iff $C^I = D^I$ and it satisfies $C \neq D$ iff $C^I \neq D^I$. An interpretation I satisfies an *ALC-TBox* \mathcal{T} iff I satisfies each axiom in \mathcal{T} . For example, if *carrier*, *person*, and *truck* are concepts and *owns* is a role, we can define exactly the persons who own a truck to be a carrier by

$$\text{carrier} = \text{person} \sqcap \exists \text{owns.truck}.$$

The assertional formalism of *ALC* allows one to introduce concrete objects by stating that they are instances of concepts and roles: If a is an object and C a concept, then $a : C$ is a *concept instance*. If a and b are objects and R is a role, then aRb is

a *role instance*. Concept instances and role instances are called *assertional axioms*, and a finite set of assertional axioms is called an *ALC-ABox*. An interpretation I maps objects to elements of its domain Δ^I and satisfies $a : C$ iff $a^I \in C^I$, and aRb iff $(a^I, b^I) \in R^I$. We assume that different objects in an ABox are mapped to different elements in Δ^I (*unique name assumption*). An interpretation I satisfies an *ALC-ABox* \mathcal{A} iff I satisfies each axiom in \mathcal{A} . As an example, if *John* and *truck-1* are objects, we can express that John owns truck-1 which is a truck by the assertional axioms

$$\text{John owns truck-1} \quad \text{and} \quad \text{truck-1} : \text{truck}.$$

Thus, we can describe the relevant concepts of a problem domain by terminological axioms, i.e., by an *ALC-TBox*, and properties of objects as well as relations between them by assertional axioms, i.e., by an *ALC-ABox*. We say an interpretation I satisfies a set Ax_1, \dots, Ax_n of terminological and assertional axioms iff I satisfies each of these axioms. We then write $I \models Ax_1, \dots, Ax_n$.

For sake of simplicity we will sometimes use the expressions $C \sqsubseteq D$ and $C \not\sqsubseteq D$ where C and D are concepts. An interpretation I satisfies $C \sqsubseteq D$ iff $C^I \subseteq D^I$ and it satisfies $C \not\sqsubseteq D$ iff $C^I \not\subseteq D^I$. The next lemma states that these expressions are abbreviations for certain terminological axioms.

Lemma 2.1 *Let C and D be concepts, and let I be an interpretation. Then*

1. *I satisfies $C \sqsubseteq D$ iff I satisfies $\neg C \sqcup D = \top$.*
2. *I satisfies $C \not\sqsubseteq D$ iff I satisfies $\neg C \sqcup D \neq \top$.*

Proof: For 1., firstly suppose I satisfies $C \sqsubseteq D$. Then for each element d in Δ^I either $d \in [\neg C]^I$ or both $d \in C^I$ and $d \in D^I$ holds. That means, I satisfies $\neg C \sqcup (C \sqcap D) = \top$ what can be simplified to $\neg C \sqcup D = \top$. Conversely, suppose I satisfies $\neg C \sqcup D = \top$. Then for each element $d \in \Delta^I$ either $d \notin C^I$ or $d \in D^I$ holds. Thus, from $d \in C^I$ follows $d \in D^I$, i.e., $C^I \subseteq D^I$. The proof of 2. is analogous. \square

For example, if *truck* and *vehicle* are concepts we can define each truck to be a vehicle by $\text{truck} \sqsubseteq \text{vehicle}$, what is an abbreviation for $\neg \text{truck} \sqcup \text{vehicle} = \top$.

2.2 The Extended Language \mathcal{ALC}_B

Now we will introduce the language \mathcal{ALC}_B which extends \mathcal{ALC} by a new operator \Box_i ; for each agent i .¹ We allow these operators in front of terminological and assertional

¹In the following, we will abbreviate agents by numbers, and we suppose only a finite number of agents to be given.

axioms. Thereby, the operator \Box_i , read as “agent i believes”, allows us to express the beliefs agent i has about the world, about beliefs of other agents, and about his own beliefs. We extend the definition of terminological and assertional axioms as follows.

- If TA is a terminological axiom, then $\Box_i TA$ and $\neg\Box_i TA$ are terminological axioms as well.
- If CI is a concept instance, then $\Box_i CI$ and $\neg\Box_i CI$ are concept instances as well.
- If RI is a role instance, then $\Box_i RI$ is a role instance as well.

Note, that we do not allow formulas of the form $\neg\Box_i(aRb)$. The reason for this restriction is that such axioms would be equivalent to stating that there exists a world in which the role instance aRb does *not* hold. And negation of roles is not allowed in \mathcal{ALC} .

These extended assertional and terminological axioms are called \mathcal{ALC}_B -formulas and can, e.g., be used to state that agent i believes that each truck is a vehicle by

$$\Box_i (\text{truck} \sqsubseteq \text{vehicle}).$$

Analogously, the \mathcal{ALC}_B -formulas $\Box_i\neg\Box_j (\text{vehicle-1} : \text{truck})$ and $\Box_i\neg\Box_i (\text{vehicle-1} : \text{truck})$ are to be read as “agent i believes that agent j doesn’t believe that vehicle-1 is a truck” and “agent i believes that he doesn’t believe truck-1 to be a truck”, respectively. Allowing \Box_i immediately in front of concepts (possibly $\Box_i C$ may be interpreted as “the set of individuals agent i believes to be a C ”) causes essential algorithmic problems and is out of the scope of this paper.

We will interpret the operators \Box_i in terms of *possible worlds*, i.e., besides the real world there exist a number of worlds agents consider to be possible. If agent i considers world w' possible at world w , we say w' is *accessible from w* by agent i . The *accessibility relation* of agent i is given by all pairs (w, w') such that w' is accessible from w by agent i . Since different worlds are possible in our approach, the interpretation of concepts and roles in \mathcal{ALC}_B -formulas depends on the world we are currently speaking of. That means, in different worlds concepts may contain different objects and roles may contain different pairs of objects. This will be expressed by taking an additional parameter, the *world parameter*, into consideration when interpreting concepts and roles. Formally, we use the notion of a K -interpretation K_I which consists of a non-empty domain Δ^{K_I} and maps objects to elements in Δ^{K_I} while satisfying the unique name assumption, atomic concepts to subsets of $\Delta^{K_I} \times \mathcal{W}$, \top to $\Delta^{K_I} \times \mathcal{W}$, \perp to $\emptyset \times \mathcal{W}$, and roles to subsets of $\Delta^{K_I} \times \Delta^{K_I} \times \mathcal{W}$. Furthermore, \sqcap is interpreted as set intersection, \sqcup as set union, and \neg as set complement w.r.t. $\Delta^{K_I} \times \mathcal{W}$, and the value and exists restrictions are interpreted by

$$\begin{aligned} [\forall R.C]^{K_I} &= \{(d, w) \mid (d', w) \in C^{K_I} \text{ for each } d' \text{ with } (d, d', w) \in R^{K_I}\} \\ [\exists R.C]^{K_I} &= \{(d, w) \mid (d', w) \in C^{K_I} \text{ for some } d' \text{ with } (d, d', w) \in R^{K_I}\}. \end{aligned}$$

Definition 2.2 A Kripke structure K is a triple $(\mathcal{W}, \Gamma, K_I)$. Thereby, \mathcal{W} is a non-empty set of worlds, Γ is a finite set of accessibility relations, one accessibility relation γ_i for each agent i , and K_I is a K -interpretation.

The satisfiability of an \mathcal{ALC}_B -formula F in a Kripke structure $K = (\mathcal{W}, \Gamma, K_I)$ and a world $w \in \mathcal{W}$, written as $K, w \models F$, is recursively defined by:

$$\begin{aligned} K, w \models C = D & \text{ iff } \{d \mid (d, w) \in C^{K_I}\} = \{d \mid (d, w) \in D^{K_I}\} \\ K, w \models C \neq D & \text{ iff } \{d \mid (d, w) \in C^{K_I}\} \neq \{d \mid (d, w) \in D^{K_I}\} \\ K, w \models a : C & \text{ iff } (a, w) \in C^{K_I} \\ K, w \models aRb & \text{ iff } (a, b, w) \in R^{K_I} \\ K, w \models \Box_i G & \text{ iff } K, w' \models G \text{ for each world } w' \text{ with } (w, w') \in \gamma_i \\ K, w \models \neg \Box_i G & \text{ iff there is a world } w' \text{ with } (w, w') \in \gamma_i \text{ and } K, w' \not\models G \end{aligned}$$

where G is an \mathcal{ALC}_B -formula, C, D are concepts, a, b are objects, and R is a role.

A set F_1, \dots, F_n of \mathcal{ALC}_B -formulas is *satisfiable* iff there exists a Kripke structure $K = (\mathcal{W}, \Gamma, K_I)$ and a world $w_0 \in \mathcal{W}$ such that $K, w_0 \models F_i$ for $i = 1, \dots, n$. We then write $K \models F_1, \dots, F_n$.

In the following we will use the notion *modality* to denote (negated) indexed \Box operators, and \mathcal{ALC}_B -formulas without any modalities are called *ALC-formulas*. For example, the \mathcal{ALC}_B -formula $\Box_i \neg \Box_j (\text{vehicle-1} : \text{truck})$ contains the modalities \Box_i and $\neg \Box_j$, and the \mathcal{ALC}_B -formula $\text{vehicle-1} : \text{truck}$ is an *ALC-formula*.

3 Testing Satisfiability of \mathcal{ALC}_B -formulas

Using \mathcal{ALC}_B -formulas, a “real world” and belief of agents can be defined as follows. The real world is given by a finite set of *ALC-formulas*, and the belief of agent i is given by a finite set of \mathcal{ALC}_B -formulas with the leading modality \Box_i . Of course, we do not only want to represent a world and beliefs of agents, but we are interested in algorithms to test (i) consistency of the represented facts, i.e., whether a given set of \mathcal{ALC}_B -formulas is satisfiable, and (ii) whether an \mathcal{ALC}_B -formula is a logical consequence of a given set of \mathcal{ALC}_B -formulas. In this section we will give an algorithm for testing satisfiability of a set of \mathcal{ALC}_B -formulas. Building upon this we will show how to decide whether or not an \mathcal{ALC}_B -formula is a logical consequence from a given set of \mathcal{ALC}_B -formulas in Section 4.

3.1 The \mathcal{ALC}_B Frame Algorithm

We will now present an algorithm for testing satisfiability of a finite set F_1, \dots, F_n of \mathcal{ALC}_B -formulas. By definition, a set F_1, \dots, F_n of \mathcal{ALC}_B -formulas is satisfiable iff there

exists a Kripke structure K such that $K \models F_1, \dots, F_n$. Of course, we are not interested in arbitrary Kripke structures to satisfy F_1, \dots, F_n , but only in Kripke structures which interpret the belief operators \Box in F_1, \dots, F_n in such a way that they satisfy the properties described in Section 1. We therefore introduce the notion of KD45 Kripke structures.

Definition 3.1 A set F_1, \dots, F_n of \mathcal{ALC}_B -formulas is KD45-satisfiable iff there exists a Kripke structure $K = (\mathcal{W}, \Gamma, K_I)$ which satisfies F_1, \dots, F_n and which has the properties

- (P1) if $K, w \models \Box_i F$ then $K, w \models \neg \Box_i \neg F$
- (P2) if $K, w \models \Box_i F$ then $K, w \models \Box_i \Box_i F$
- (P3) if $K, w \models \neg \Box_i F$ then $K, w \models \Box_i \neg \Box_i F$

for each \mathcal{ALC}_B -formula F , for each agent i , and for each world $w \in \mathcal{W}$.² A Kripke structure which satisfies (P1), (P2), and (P3) is called KD45 Kripke structure.

Property (P1) corresponds to “an agent cannot believe in both a fact and its negation”, (P2) to “if an agent believes something, then he believes that he believes it”, and (P3) to “if an agent does not believe in a fact then he believes that he does not believe in this fact”. The property “agents must be able to reason on the basis of their beliefs”, is guaranteed by choosing Kripke structures for the representation of belief (cf., e.g., [17]).

It is a well-known fact that $K = (\mathcal{W}, \Gamma, K_I)$ is a KD45 Kripke structure if the accessibility relation γ_i of each agent i is serial, Euclidean, and transitive (see, e.g., [25]). A relation $\gamma \subseteq \mathcal{W} \times \mathcal{W}$ is

- *serial* iff for each u in \mathcal{W} there is a v in \mathcal{W} such that $(u, v) \in \gamma$,
- *Euclidean* iff for all u, v, w in \mathcal{W} holds: if $(u, v) \in \gamma$ and $(u, w) \in \gamma$ then $(v, w) \in \gamma$,
- *transitive* iff for all u, v, w in \mathcal{W} holds: if $(u, v) \in \gamma$ and $(v, w) \in \gamma$ then $(u, w) \in \gamma$.

We will use the standard notation $\Diamond_i F$ as an abbreviation for $\neg \Box_i \neg F$ such that the properties (P1) and (P3) can be rewritten as

- (P1') if $K, w \models \Box_i F$ then $K, w \models \Diamond_i F$
- (P3') if $K, w \models \Diamond_i F$ then $K, w \models \Box_i \Diamond_i F$.

In the following we will use the one or the other version of these two properties, whichever is more appropriate.

²Since these properties hold for *arbitrary* worlds this amounts in saying that all these properties are mutually believed, i.e., each agent's belief has these properties, each agent believes that each agent's belief has these properties and so on.

To keep notation simple we transform \mathcal{ALC}_B -formulas into negation normal form. An \mathcal{ALC}_B -formula (concept) is in *negation normal form* iff in the formula (concept) negation signs occur immediately in front of atomic concepts only. Concepts can be transformed into an equivalent negation normal form by the rules

$$\begin{array}{lll} \neg\neg C & \rightarrow & C \\ \neg\top & \rightarrow & \perp \\ \neg\perp & \rightarrow & \top \\ \neg(C \sqcap D) & \rightarrow & \neg C \sqcup \neg D \\ \neg(C \sqcup D) & \rightarrow & \neg C \sqcap \neg D \\ \neg(\forall R.C) & \rightarrow & \exists R.\neg C \\ \neg(\exists R.C) & \rightarrow & \forall R.\neg C \end{array}$$

where C is a concept and R is a role (see, e.g., [20]). Building upon this, \mathcal{ALC}_B -formulas can be transformed into negation normal form by applying the rules

$$\begin{array}{lll} \neg\neg F & \rightarrow & F \\ \neg\Box_i F & \rightarrow & \Diamond_i \neg F \\ \neg\Diamond_i F & \rightarrow & \Box_i \neg F \\ \neg(C = D) & \rightarrow & a_n : (C \sqcap \neg D) \sqcup (\neg C \sqcap D) \\ \neg(C \neq D) & \rightarrow & C = D \\ \neg(a : C) & \rightarrow & a : \neg C \end{array}$$

to the outermost negation sign. Thereby, F is an \mathcal{ALC}_B -formula, C, D are concepts, a is an object, and a_n is a new object. For example, the negation normal form of the \mathcal{ALC}_B -formula

$$\neg\Box_i(A = \neg(\forall R.C)) \quad \text{is} \quad \Diamond_i(a_n : (A \sqcap \forall R.C) \sqcup (\neg A \sqcap \exists R.\neg C))$$

where a_n is a new object. The next lemma states that an \mathcal{ALC}_B -formula is KD45-satisfiable iff its negation normal form is KD45-satisfiable.

Lemma 3.2 *Let F be an \mathcal{ALC}_B -formula, F^* be the negation normal form of F , and K be a KD45 Kripke structure. Then $K \models F$ iff there is a KD45 Kripke structure K' such that $K' \models F^*$.*

Proof: If we apply one of the rules $\neg\neg F \rightarrow F$, $\neg\Box_i F \rightarrow \Diamond_i \neg F$, $\neg\Diamond_i F \rightarrow \Box_i \neg F$, or $\neg(a : C) \rightarrow a : \neg C$ to an \mathcal{ALC}_B -formula F , then K obviously satisfies the formula on the left hand side of the rule iff K satisfies the right hand side of the rule.

Now suppose $K = (\mathcal{W}, \Gamma, K_I)$ and $K \models \neg(C = D)$, i.e., $K, w_0 \models \neg(C = D)$ for some world w_0 in \mathcal{W} . In this case there is an element $u \in \Delta^{K_I}$ such that (u, w_0) is either in C^{K_I} and in $[\neg D]^{K_I}$ or in $[\neg C]^{K_I}$ and in D^{K_I} . Let now $K' = (\mathcal{W}', \Gamma', K'_I)$ be a Kripke structure which is identical to K but $a_n^{K_I} := u$. Then, K' is a KD45 Kripke structure³ and $K', w_0 \models a_n : (C \sqcap \neg D) \sqcup (\neg C \sqcap D)$. Conversely, if K' is a KD45 Kripke structure such that $K', w_0 \models a_n : (C \sqcap \neg D) \sqcup (\neg C \sqcap D)$, then obviously $K', w_0 \models \neg(C = D)$. \square

If F is an \mathcal{ALC}_B -formula in negation normal form it has a (possibly empty) leading sequence $o^* = o_{i_1} \dots o_{i_m}$ of non-negated modalities where each o_{i_j} is either \Box or \Diamond and

³Note, that a_n is a *new* element and the unique name assumption only has to hold for objects occurring in an ABox.

each index i_j is an agent. We now replace each subsequence of modalities indexed with the same agent in \circ^* by the last modality in this subsequence. The obtained \mathcal{ALC}_B -formula is called the *KD45 normal form* F' of F . For example, the KD45 normal form of $\Box_1 \Diamond_1 \Diamond_2 \Box_2 \Box_1 (a : C)$ is given by $\Diamond_1 \Box_2 \Box_1 (a : C)$. As an immediate consequence of Proposition 4.27 in [10], for each KD45 Kripke structure $K = (\mathcal{W}, \Gamma, K_I)$ and for each world $w \in \mathcal{W}$ holds that $K, w \models F$ iff $K, w \models F'$.

Assumption: In the following we suppose each \mathcal{ALC}_B -formula to be in KD45 normal form (and thus especially in negation normal form).

To formulate a calculus for testing KD45 satisfiability of \mathcal{ALC}_B formulas we introduce the notions of labeled \mathcal{ALC}_B -formulas and of a world constraint system. A *labeled \mathcal{ALC}_B -formula* consists of an \mathcal{ALC}_B -formula F together with a label w , written as $F \parallel w$. Thereby, w is a constant representing a world in which F holds. A *world constraint* is either a labeled \mathcal{ALC}_B -formula or a term $w \bowtie_i w'$. The constants w and w' represent worlds and \bowtie_i represents the accessibility relation of agent i . A *world constraint system* is a finite, non-empty set of world constraints.

A Kripke structure $K = (\mathcal{W}, \Gamma, K_I)$ satisfies a world constraint system W iff for each label w in W there is a world $w^K \in \mathcal{W}$ such that (i) $K, w^K \models F$ for each world constraint $F \parallel w$ in W and (ii) $(w^K, v^K) \in \gamma_i$ for each world constraint $w \bowtie_i v$ in W . A world constraint system W is (KD45-)satisfiable iff there exists a (KD45) Kripke structure which satisfies W .

For testing KD45-satisfiability of a set F_1, \dots, F_n of \mathcal{ALC}_B -formulas we firstly translate them into a world constraint system. The world constraint system W is *induced by* F_1, \dots, F_n iff $W = \{F_1 \parallel w_0, \dots, F_n \parallel w_0\}$, where w_0 is a new constant (which represents the real world). Obviously, F_1, \dots, F_n are KD45-satisfiable iff W is KD45-satisfiable.

KD45-satisfiability of a world constraint system W is tested by the *frame algorithm* which has a world constraint system as input that is induced by a finite number of \mathcal{ALC}_B -formulas and which successively adds new world constraints to W by applying the four propagation rules in Figure 1. Thus, the result of the frame algorithm with input W is a (modified) world constraint system W' . We will call each world constraint system that can be obtained from W by a finite number of propagation rule applications a *derived system* in the following.

The intuitive idea behind these propagation rules is as follows: Firstly, for $W \rightarrow_{\Diamond} W'$, if there is a world constraint $\Diamond_i F \parallel w$ in W we add a world v such that (i) v is accessible from w by agent i and (ii) $F \parallel v$ holds. Furthermore, whenever $\Diamond_i F_j \parallel w$ is in W we add $\Diamond_i F_j \parallel v$ because of property (P'_3), and whenever $\Box_i G_k \parallel w$ is in W we add both $\Box_i G_k \parallel v$ and $G_k \parallel v$ because of property (P_2) and the definition of \Box_i . This rule is similar to the unsigned prefixed KD45 tableaux rules in [14]. Secondly, for $W \rightarrow_{\Box} W'$, if $\Box_i G_1 \parallel w, \dots, \Box_i G_m \parallel w$ are in W but there is no world u accessible from w by agent i , we have to introduce a new world v —accessible from w by agent i —where $G_1 \parallel v, \dots, G_m \parallel v$ and $\Box_i G_1 \parallel v, \dots, \Box_i G_m \parallel v$ holds. This is due to the properties

$W \rightarrow_{\Diamond} \{w \bowtie_i v, F \parallel v, \Diamond_i F_1 \parallel v, \dots, \Diamond_i F_n \parallel v, G_1 \parallel v, \Box_i G_1 \parallel v, \dots, G_m \parallel v, \Box_i G_m \parallel v\} \cup W$
 if $\Diamond_i F \parallel w, \Diamond_i F_1 \parallel w, \dots, \Diamond_i F_n \parallel w$ are the world constraints with leading modality \Diamond_i in W , $\Box_i G_1 \parallel w, \dots, \Box_i G_m \parallel w$ are the world constraints with leading modality \Box_i in W , there is no label u in W such that the world constraints $F \parallel u, \Diamond_i F_1 \parallel u, \dots, \Diamond_i F_n \parallel u, G_1 \parallel u, \dots, G_m \parallel u, \Box_i G_1 \parallel u, \dots, \Box_i G_m \parallel u$ are exactly the labeled \mathcal{ALC}_B -formulas with label u in W , and v is a new label.

$W \rightarrow_{\Diamond_0} \{w \bowtie_i u\} \cup W$
 if $\Diamond_i F \parallel w, \Diamond_i F_1 \parallel w, \dots, \Diamond_i F_n \parallel w$ are the world constraints with leading modality \Diamond_i in W , $\Box_i G_1 \parallel w, \dots, \Box_i G_m \parallel w$ are the world constraints with leading modality \Box_i in W , there is a label u in W such that the world constraints $F \parallel u, \Diamond_i F_1 \parallel u, \dots, \Diamond_i F_n \parallel u, G_1 \parallel u, \dots, G_m \parallel u, \Box_i G_1 \parallel u, \dots, \Box_i G_m \parallel u$ are exactly the labeled \mathcal{ALC}_B -formulas with label u in W , and $w \bowtie_i u$ is not in W .

$W \rightarrow_{\Box} \{w \bowtie_i v, G_1 \parallel v, \Box_i G_1 \parallel v, \dots, G_m \parallel v, \Box_i G_m \parallel v\} \cup W$
 if no world constraint of the form $\Diamond_i F \parallel w$ is in W , $\Box_i G_1 \parallel w, \dots, \Box_i G_m \parallel w$ are the world constraints with leading modality \Box_i in W , there is no label u in W such that the world constraints $G_1 \parallel u, \dots, G_m \parallel u, \Box_i G_1 \parallel u, \dots, \Box_i G_m \parallel u$ are exactly the labeled \mathcal{ALC}_B -formulas with label u in W , and v is a new label.

$W \rightarrow_{\Box_0} \{w \bowtie_i u\} \cup W$
 if no world constraint of the form $\Diamond_i F \parallel w$ is in W , $\Box_i G_1 \parallel w, \dots, \Box_i G_m \parallel w$ are the world constraints with leading modality \Box_i in W , there is a label u in W such that $G_1 \parallel u, \dots, G_m \parallel u, \Box_i G_1 \parallel u, \dots, \Box_i G_m \parallel u$ are exactly the labeled \mathcal{ALC}_B -formulas with label u in W , and $w \bowtie_i u$ is not in W .

Figure 1: Propagation rules of the \mathcal{ALC}_B frame algorithm.

(P'_1) and (P_2) of KD45 Kripke structures. Finally, the rules \rightarrow_{\Diamond_0} and \rightarrow_{\Box_0} are used to guarantee termination of applying propagation rules to a derived system.

Now we will show that the \mathcal{ALC}_B frame algorithm has the following two important properties. Firstly, it terminates for every world constraint system W as input which is induced by a finite set F_1, \dots, F_n of \mathcal{ALC}_B -formulas. Secondly, if W' is the result of the frame algorithm with input W , then F_1, \dots, F_n are KD45-satisfiable iff for each label w in W' the set of \mathcal{ALC} -formulas with label w in W' is satisfiable.

Termination of the frame algorithm is stated by the next proposition. Its proof will employ techniques which have been developed for proving termination of term rewriting systems (see [11]).

Proposition 3.3 *If W is a world constraint system which is induced by a finite set of \mathcal{ALC}_B -formulas, there is no infinite chain of applications of propagation rules to W .*

In order to prove this proposition we will map derived systems to multisets which are equipped with a well-founded strict partial ordering \gg . Multisets are like sets, but multiple occurrences of elements are allowed, e.g., $\{2, 3, 3, 4\}$ is a multiset which is different from the set $\{2, 3, 4\}$.

A given ordering $>$ on elements in a set S can be extended to an ordering \gg on finite multisets over S as follows. If M and M' are finite multisets over S then $M \gg M'$ iff M' is obtained from M by replacing one or more elements in M by a finite number of elements in S , each of which is smaller than one of the replaced elements (w.r.t. $>$). For example, $\{2, 3, 3, 4\}$ is larger than the multisets $\{2, 3, 1, 2, 3\}$ and $\{4\}$. Dershowitz and Manna [11] showed that the ordering \gg on finite multisets over S is well-founded if the original ordering on S is so.

We will use a mapping Ψ which maps derived systems to multisets whose elements are pairs of non-negative integers. These pairs are ordered lexicographically from left to right, i.e., (c_1, c_2) is larger than (c'_1, c'_2) iff (i) $c_1 > c'_1$ or (ii) $c_1 = c'_1$ and $c_2 > c'_2$. Since the orderings on both components are well-founded, the lexicographical ordering on these pairs is also well-founded. Finite multisets of these pairs are now compared w.r.t. the multiset ordering which is induced by this lexicographical ordering. This is the well-founded ordering \gg mentioned above.

In order to simplify notation of the mapping Ψ we introduce the following notions. If $F \parallel w$ and $G \parallel w$ are labeled \mathcal{ALC}_B -axioms with the same label w , we say that $G \parallel w$ is a *modal subformula* of $F \parallel w$ iff there is a (possibly empty) sequence \circ^* of modalities such that \circ^*G and F are identical. For example, $\Diamond_j \Box_i \Diamond_j F \parallel w$, $\Box_i \Diamond_j F \parallel w$, $\Diamond_j F \parallel w$, and $F \parallel w$ are all modal subformulas of $\Diamond_j \Box_i \Diamond_j F \parallel w$. If $F \parallel w$ is a labeled \mathcal{ALC}_B -formula we denote the set $\{G \mid G \parallel w \text{ is a modal subformula of } F \parallel w\}$ by $MSub(F \parallel w)$. For a derived system W , $MSub^W(w)$ denotes the set $\bigcup_{F \parallel w \in W} MSub(F \parallel w)$.

Now we can define the mapping Ψ as follows.

Definition 3.4 *Let W_0 be a world constraint system which is induced by a finite number of \mathcal{ALC}_B -formulas, let W be a world constraint system which is derived from W_0 by applications of the propagation rules of the \mathcal{ALC}_B frame algorithm, and let w be a label in W . We define*

1. \mathcal{S} is the power set of $\{MSub^{W_0}(w_0) \mid w_0 \text{ is the (only) label in } W_0\}$, and $N_{\mathcal{S}}$ is the number of elements in \mathcal{S} .

2. N_L^W is the number of labels in W which are different from w_0 .
3. $\psi_1^W(w)$ is the number of modalities in the set $\{F \parallel w \mid F \parallel w \in W\}$.
4. $\psi_2^W(w)$ is the cardinality of $\{w \bowtie_j v \mid w \bowtie_j v \in W \text{ for some agent } j \text{ and label } v\}$.
5. The mapping $\psi^W(w)$ is now defined by the pair

$$(N_S - N_L^W, \psi_1^W(w) - \psi_2^W(w))$$

of integers, and $\Psi(W)$ is the multiset $\{\psi^W(w) \mid w \text{ is a label in } W\}$.

Given a derived system W and a label w in W we firstly show that both components of $\psi^W(w)$ are non-negative integers. It is easy to verify that each of the four propagation rules only introduces labeled \mathcal{ALCB} -formulas $H_1 \parallel v, \dots, H_n \parallel v$ such that each H_i is an \mathcal{ALCB} -formula in $MSub^{W_0}(w_0)$. In other words, the set $\{H_1, \dots, H_n\}$ is an element in \mathcal{S} . Because of the definition of the propagation rules a new label v —together with some labeled \mathcal{ALCB} -formulas $H_1 \parallel v, \dots, H_n \parallel v$ —is only introduced if there is no label u in W such that the world constraints $H_1 \parallel u, \dots, H_n \parallel u$ are exactly the labeled \mathcal{ALCB} -formulas with label u in W . That means, for each element in \mathcal{S} at most one new label can be introduced, and hence $N_S - N_L^W \geq 0$ for each derived system W . On the other hand, since each propagation rule can be applied to a modality in a labeled \mathcal{ALCB} -axiom at most once (because of the disjoint preconditions of the four propagation rules) and adds exactly one world constraint of the form $w \bowtie_j v$, we can conclude that $\psi_1^W(w) - \psi_2^W(w) \geq 0$ for each label w in W .

We will now show that each chain $W_0 \rightarrow_1 W_1 \rightarrow_2 \dots \rightarrow_n W_n$ of propagation rule applications to derived systems corresponds to the decreasing chain $\Psi(W_0) \gg \Psi(W_1) \gg \dots \gg \Psi(W_n)$. Thus, proposition 3.3 is an immediate consequence of the next lemma.

Lemma 3.5 *If W' is obtained from the derived system W by an application of a propagation rule, then $\Psi(W) \gg \Psi(W')$.*

Proof: We have to show this lemma for each of the propagation rules.

(1) Assume that W' is obtained from W by applying the \rightarrow_\diamond or the \rightarrow_\square rule to one or more world constraints labeled with label w . Such a rule application adds a world constraint $w \bowtie_j v$ to W , where j is an agent and v is a new label. Hence, the number of labels in W' is greater than the number of labels in W and for each label u in W' the first component of $\psi^{W'}(u)$ is less than the first component of $\psi^W(w)$, i.e., $\Psi(W) \gg \Psi(W')$.

(2) Now assume that W' is obtained from W by an application of the \rightarrow_{\diamond_0} or the \rightarrow_{\square_0} rule. Such a propagation rule application does not introduce a new label

to W . Thus, the first components of $\Psi^W(u)$ and $\Psi^{W'}(u)$ are identical for each label u . Let us now consider how the values $\psi_1^W(w)$ and $\psi_2^W(w)$ are changed. Obviously, $\psi_1^W(w) = \psi_1^{W'}(w)$ and $\psi_2^W(w) = \psi_2^{W'}(w) + 1$. From this it follows immediately that $\Psi^W(w) \gg \Psi^{W'}(w)$. \square

Thus, the application of the frame algorithm to a world constraint system W induced by a finite set of \mathcal{ALC}_B -formulas F_1, \dots, F_n terminates and results a world constraint system, say W' . In order to test KD45-satisfiability of W' , for each label w in W' we determine the set of all those \mathcal{ALC}_B -formulas which are labeled by w and which do not contain any modality. That means, such a set contains only \mathcal{ALC} -formulas and is therefore called the \mathcal{ALC} test set of label w . More formally, if W' is a world constraint system, the \mathcal{ALC} test set $A(w)$ of label w in W' is given by the set

$$\{F \mid F \parallel w \in W' \text{ and } F \text{ does not contain any modality}\}.$$

We will show in the following that a finite set F_1, \dots, F_n of \mathcal{ALC}_B -formulas is KD45-satisfiable iff the \mathcal{ALC} test set $A(w)$ of each label in W' is satisfiable. Thereby, W' is the result of the frame algorithm with input $\{F_1 \parallel w_0, \dots, F_n \parallel w_0\}$. One direction is given by the next lemma.

Lemma 3.6 *Let W be a world constraint system which is induced by the finite set $\{F_1, \dots, F_n\}$ of \mathcal{ALC}_B -formulas, and let W' be the result of the frame algorithm with input W . If $K = (\mathcal{W}, \Gamma, K_I)$ is a KD45 Kripke structure which satisfies W , then for each label w in W' there is a world $w^K \in \mathcal{W}$ such that $K, w^K \models F$ for each labeled \mathcal{ALC}_B -formula $F \parallel w$ in W' .*

Proof: Since W' is the result of the frame algorithm with input W there is a chain $W = W_0 \rightarrow_1 W_1 \rightarrow_2 \dots \rightarrow_k W_k = W'$ with $\rightarrow_i \in \{\rightarrow_\diamond, \rightarrow_{\diamond_0}, \rightarrow_\square, \rightarrow_{\square_0}\}$ for $i \in \{1, \dots, k\}$. We will show that K satisfies each labeled \mathcal{ALC}_B -formula in W' by induction over the number of rule applications in this chain. By assumption, $K = (\mathcal{W}, \Gamma, K_I)$ satisfies $W_0 = \{F_1 \parallel w_0, \dots, F_n \parallel w_0\}$, i.e., there is a world w_0^K in \mathcal{W} such that $K, w_0^K \models F_1, \dots, K, w_0^K \models F_n$.

We thus can assume that, after j rule applications, for each label w in W_j there is a world w^K in \mathcal{W} such that $K, w^K \models F$ for each labeled \mathcal{ALC}_B -formula $F \parallel w$ in W_j . If $W_j \rightarrow_j W_{j+1}$ there are four possibilities. Firstly, suppose $W_j \rightarrow_\diamond W_{j+1}$ by applying the \rightarrow_\diamond rule to $\diamond_i F \parallel w$ in W_j . In this case, there are labeled \mathcal{ALC}_B -formulas $\diamond_i F \parallel w, \diamond_i F_1 \parallel w, \dots, \diamond_i F_n \parallel w, \square_i G_1 \parallel w, \dots, \square_i G_m \parallel w$ in W_j , and W_{j+1} is given by $W_j \cup \{w \bowtie_i v, F \parallel v, \diamond_i F_1 \parallel v, \dots, \diamond_i F_n \parallel v, \square_i G_1 \parallel v, \dots, \square_i G_m \parallel v, G_1 \parallel v, \dots, G_m \parallel v\}$

where v is a new label. By induction hypothesis there is a world w^K in \mathcal{W} such that (i) $K, w^K \models \diamond_i F$ and (ii) $K, w^K \models \square_i G_j$ for $j = 1, \dots, m$. Because of (i) there is a world v^K in \mathcal{W} (not necessarily different from w^K) such that $(w^K, v^K) \in \gamma_i$ and

$K, v^K \models F$. Furthermore, because of (ii) and property (P2) of KD45 Kripke structures, both $K, w^K \models \Box_i G_j$ and $K, w^K \models \Box_i \Box_i G_j$ holds for $j = 1, \dots, m$. And thus, since $(w^K, v^K) \in \gamma_i$, especially $K, v^K \models G_j$ and $K, v^K \models \Box_i G_j$ holds for $j = 1, \dots, m$. Finally, $K, w^K \models \Diamond_i F_j$ for $j = 1, \dots, n$ by induction hypotheses, i.e., $K, w^K \models \Box_i \Diamond_i F_j$ because of property (P3') of KD45 Kripke structures. Hence, also $K, v^K \models \Diamond_i F_j$ holds for $j = 1, \dots, n$. Summing up, for each labeled \mathcal{ALC}_B -formula $F \parallel v$ in W_j we have $K, v^K \models F$.

Secondly, if $W_j \rightarrow_{\Box} W_{j+1}$ there are world constraints $\Box_i G_1 \parallel w, \dots, \Box_i G_m \parallel w$ in W_j and $W_{j+1} = W_j \cup \{w \boxtimes_i v, G_1 \parallel v, \Box_i G_1 \parallel v, \dots, G_m \parallel v, \Box_i G_m \parallel v\}$ where v is a new label. By induction hypothesis, $K, w^K \models \Box_i G_j$ for $j = 1, \dots, m$ and for some world $w^K \in \mathcal{W}$. Because of $K, w^K \models \Box_i G_1$ and property (P1') of KD45 Kripke structures especially $K, w^K \models \Diamond_i G_1$ holds. Hence, there is a world v^K (not necessarily different from w^K) such that $(w^K, v^K) \in \gamma_i$ and $K, v^K \models G_1$. Furthermore, because of property (P2) of KD45 Kripke structures, both $K, w^K \models \Box_i G_j$ and $K, w^K \models \Box_i \Box_i G_j$ holds for $j = 1, \dots, m$. And thus, since $(w^K, v^K) \in \gamma_i$, especially $K, v^K \models G_j$ and $K, v^K \models \Box_i G_j$ holds for $j = 1, \dots, m$. Summing up, we have $K, v^K \models \Box_i G_j$ and $K, v^K \models G_j$ for $j = 1, \dots, m$.

Finally, if $W_j \rightarrow_{\Diamond} W_{j+1}$ or $W_j \rightarrow_{\Box} W_{j+1}$ there is nothing to show since these rules do not change the set of labeled \mathcal{ALC}_B -formulas in W_j at all. \square

The next lemma states that a world constraint system W' , which is a result of the frame algorithm, is KD45-satisfiable if the \mathcal{ALC} test set of each label in W' is satisfiable.

Lemma 3.7 *Let W be a world constraint system which is induced by a finite set of \mathcal{ALC}_B -formulas, and let W' be the result of the frame algorithm with input W . If the \mathcal{ALC} test set $A(w)$ of each label w in W' is satisfiable, then W is KD45-satisfiable.*

Proof: Let K be the Kripke structure $(\mathcal{W}, \Gamma, K_I)$ where

- \mathcal{W} is given by the set of all labels in W' .
- Γ consists of an accessibility relation γ_i for each agent i . Thereby, γ_i is given by the Euclidean and transitive closure of the set $\{(w, v) \mid w \boxtimes_i v \text{ in } W'\} \cup \{(w, w) \mid w \text{ is a label in } W' \text{ and for no label } v \text{ in } W' \text{ there is } w \boxtimes_i v \text{ in } W'\}$. It is easy to verify that each accessibility relation γ_i is serial.
- K_I is given such that $K, w \models F$ for each labeled \mathcal{ALC}_B -formula $F \parallel w$ in W' where F does not contain any modality. Such a K -interpretation K_I exists, since we assumed the \mathcal{ALC} test set of each label in W' to be satisfiable. Given interpretations I_1, \dots, I_n which satisfy the \mathcal{ALC} test sets of each label in W' respectively, the construction of K_I is straightforward.

Obviously, each γ_i is Euclidean, transitive, and serial and hence K is a KD45 Kripke structure.

We will now prove that K satisfies each world constraint c in W' . If c is of the form $w \boxtimes_i v$ there is nothing to show because of the definition of K . The fact $K \models F \parallel w$ for each labeled \mathcal{ALC}_B -formula $F \parallel w$ in W' can be shown by induction over the number of modalities in F . If F does not contain any modality, then $K, w \models F$ because of the construction of K . Thus we can assume that $K, w \models F$ for each labeled \mathcal{ALC}_B -formula $F \parallel w$ in W' such that F contains n modalities.

If F contains $n + 1$ modalities, there are two possibilities: the leading modality is either \Box_i or \Diamond_i . Firstly, suppose W' contains a world constraint $\Box_i F \parallel w$, where F has n modalities. We then have to show that $K, w \models \Box_i F$, i.e., that $K, v \models F$ for each v such that $(w, v) \in \gamma_i$. Since $\Box_i F \parallel w$ is in W' , during the frame algorithm a propagation rule has been applied to the world constraints with label w , such that $w \boxtimes_i u$ is in W' for some label u . It is easy to verify that $w \boxtimes_i v \in W'$ for some label v (not necessarily different from w) if $\Box_i F \parallel w$ or $\Diamond_i F \parallel w$ is in W' for some \mathcal{ALC}_B -formula F . Thus, because each γ_i is transitive and Euclidean, there are two possibilities for $(w, v) \in \gamma_i$:

1. there is a path $w = w_{i_1} \boxtimes_i w_{i_2}, w_{i_2} \boxtimes_i w_{i_3}, \dots, w_{i_{k-1}} \boxtimes_i w_{i_k} = v$ in W' , or
2. there are two paths starting with some label z , namely

$$z = w_{i_1} \boxtimes_i w_{i_2}, w_{i_2} \boxtimes_i w_{i_3}, \dots, w_{i_{k-1}} \boxtimes_i w_{i_k} = v \text{ and}$$

$$z = \tilde{w}_{j_1} \boxtimes_i \tilde{w}_{j_2}, \tilde{w}_{j_2} \boxtimes_i \tilde{w}_{j_3}, \dots, \tilde{w}_{j_{l-1}} \boxtimes_i \tilde{w}_{j_l} = w \text{ in } W'.$$

For case 1., assume that W' contains a world constraint $\Box_i F \parallel w$. Because of the definition of the propagation rules it holds that, whenever $\Box_i F \parallel w$ is in W' and $w \boxtimes_i v$ (or $w \boxtimes_i w$) is added to some derived system W_j , then both $F \parallel v$ and $\Box_i F \parallel v$ are in W_j . Analogously, whenever a world constraint $v \boxtimes_i u$ is added to a system $W_{j'}$ (with $j' \geq j + 1$), the derived system $W_{j'}$ contains $\Box_i F \parallel u$ and $F \parallel u$, and so on. Hence, since none of the propagation rules deletes a world constraint, we can conclude that $F \parallel v$ is in W' if there is a path $w = w_{i_1} \boxtimes_i w_{i_2}, w_{i_2} \boxtimes_i w_{i_3}, \dots, w_{i_{k-1}} \boxtimes_i w_{i_k} = v$ in W' . And, by induction hypothesis, we know that $K, v \models F$ because F contains only n modalities.

For the second case, assume the two above given paths starting with label z are in W' . We firstly show that

$$(\star) \quad \text{if } \Box_i F \parallel w \text{ is in } W', \text{ then } \Box_i F \parallel z \text{ is also in } W'.$$

It is sufficient to show that $\Box_i F \parallel u$ is in W_j whenever $\Box_i F \parallel u'$ is in W_{j+1} , where $W_j \rightarrow_\Diamond W_{j+1}$, $W_j \rightarrow_{\Diamond\Diamond} W_{j+1}$, $W_j \rightarrow_\Box W_{j+1}$, or $W_j \rightarrow_{\Box\Diamond} W_{j+1}$, and $u \boxtimes_i u'$ is added to W_j by this rule application. This holds because of the definitions of the four propagation rules and since \mathcal{ALC}_B -formulas are in KD45 normal form, such that (\star) follows immediately.

1. Let W be the world constraint system which is induced by F_1, \dots, F_n .
2. Let W' be the result of the frame algorithm with input W .
3. For each label w in W' do: If the \mathcal{ALC} test set of w is not satisfiable, then STOP and return "KD45-unsatisfiable".
4. Return "KD45-satisfiable".

Figure 2: The KD45-satisfiability algorithm.

That means, if $\Box_i F \parallel w$ is in W' we know that $\Box_i F \parallel z$ is in W' . As argued above, in this case W' contains the labeled \mathcal{ALC}_B -formula $F \parallel v$ because of the path $z \bowtie_i \dots \bowtie_i v$. Again, we know $K, v \models F$ because of the induction hypothesis.

Secondly suppose W' contains $\Diamond_i F \parallel w$. We then have to show that $K, v \models F$ for some world v such that $(w, v) \in \gamma_i$. This is obvious, since either (i) the \rightarrow_\Diamond rule has been applied to $\Diamond_i F \parallel w$ and added both world constraints $w \bowtie_i v$ and $F \parallel v$ for some label v , or (ii) the $\rightarrow_{\Diamond_\circ}$ rule has been applied to $\Diamond_i F \parallel w$ and added $w \bowtie_i u$ to W' such that $F \parallel u \in W'$. Summing up, K satisfies each world constraint in W' and hence in $W \subseteq W'$. Thus, W is KD45-satisfiable. \square

The next theorem summarizes the previous results.

Theorem 3.8 *Let F_1, \dots, F_n be a finite set of \mathcal{ALC}_B -formulas, and let W be the world constraint system which is induced by F_1, \dots, F_n . If W' is the result of the frame algorithm with input W , then the set F_1, \dots, F_n is KD45-satisfiable iff the \mathcal{ALC} test set $A(w)$ of each label w in W' is satisfiable.*

Proof: The \mathcal{ALC}_B -formulas F_1, \dots, F_n are KD45-satisfiable iff W is KD45-satisfiable. Firstly, suppose K is a KD45 Kripke structure which satisfies W . Then, because of Lemma 3.6, for each label w in W' there is a world $w^K \in \mathcal{W}$ such that $K, w^K \models F$ for each \mathcal{ALC}_B -formula $F \parallel w$ in W' . Thus, especially the \mathcal{ALC} test set of each label w in W' is satisfied by K . Conversely, suppose that the \mathcal{ALC} test set $A(w)$ of each label w in W' is satisfiable. Then W is KD45-satisfiable because of Lemma 3.7. \square

Summing up, we obtain the algorithm for testing KD45-satisfiability of \mathcal{ALC}_B -formulas F_1, \dots, F_n which is given in Figure 2. An algorithm for testing satisfiability of \mathcal{ALC} test sets has been given in [21].

Unfortunately, the number of labels in the constructed world constraint system is in the worst case exponential in the number of input \mathcal{ALC}_B -formulas: Let $\Diamond_1 F_1, \dots, \Diamond_1 F_n$ be \mathcal{ALC}_B -formulas where each F_i is an \mathcal{ALC} -formula. The induced world constraint

system W_0 is then given by

$$W_0 = \{\Diamond_1 F_1 \parallel w_0, \dots, \Diamond_1 F_n \parallel w_0\}.$$

Applying the \rightarrow_\Diamond rule successively to $\Diamond_1 F_1 \parallel w_0, \dots, \Diamond_1 F_n \parallel w_0$ results the world constraint systems

$$\begin{aligned} W_1 &= W_0 \cup \{w_0 \bowtie_1 w_1, F_1 \parallel w_1, \Diamond_1 F_2 \parallel w_1, \dots, \Diamond_1 F_n \parallel w_1\} \\ W_2 &= W_1 \cup \{w_0 \bowtie_1 w_2, \Diamond_1 F_1 \parallel w_2, F_2 \parallel w_2, \Diamond_1 F_3 \parallel w_2, \dots, \Diamond_1 F_n \parallel w_2\} \\ &\vdots \\ W_n &= W_{n-1} \cup \{w_0 \bowtie_1 w_n, \Diamond_1 F_1 \parallel w_n, \dots, \Diamond_1 F_{n-1} \parallel w_n, F_n \parallel w_n\} \end{aligned}$$

Each of the world constraint systems W_1, \dots, W_n contains $n-1$ labeled \mathcal{ALC}_B -formulas with a leading modality \Diamond_1 , i.e., to each of these derived systems the \rightarrow_\Diamond rule can be applied $n-1$ times.⁴ To each of the thereby obtained derived systems the \rightarrow_\Diamond rule can be applied $n-2$ times and so on. summing up, an exponential number of labels—and hence of labeled \mathcal{ALC}_B -formulas—is generated.

In order to test KD45-satisfiability of a set of \mathcal{ALC}_B -formulas we are mainly interested in the constructed \mathcal{ALC} test sets which have to be tested on satisfiability. The following example shows that the number of different \mathcal{ALC} test sets in a derived system W may essentially be smaller than the number of labels in W .

Example 3.9 *Let the \mathcal{ALC}_B -formulas $\Diamond_1 F_1, \Diamond_1 F_2, \Diamond_1 F_3$ be given where F_1, F_2, F_3 are \mathcal{ALC} -formulas. Applying the frame algorithm to*

$$\{\Diamond_1 F_1 \parallel w_0, \Diamond_1 F_2 \parallel w_0, \Diamond_1 F_3 \parallel w_0\}$$

results a world constraint system W' which has 13 different labels. However, only three different \mathcal{ALC} test sets (namely $\{F_1\}$, $\{F_2\}$, and $\{F_3\}$) are constructed.

An optimized version of the KD45-satisfiability algorithm—which does not generate a worst case exponential overhead of labeled \mathcal{ALC}_B -formulas—is presented in the following section.

4 Optimization and Computing \mathcal{ALC}_B -Inferences

In this section we consider optimizations of the KD45-satisfiability algorithm as well as the problem of computing logical inferences from given \mathcal{ALC}_B -formulas. In 4.1 we

⁴For sake of simplicity we disregard the \rightarrow_\Diamond rule here. It is easy to verify that applications of this rule do not influence the exponential behaviour of the frame algorithm with input $\Diamond_1 F_1, \dots, \Diamond_1 F_n$.

will present an algorithm which—based upon the results of the previous section—determines \mathcal{ALC} test sets without computing an exponential number of labeled \mathcal{ALC}_B -formulas. These sets then have to be tested on satisfiability, i.e., we have to decide whether or not a set of terminological and assertional axioms is satisfiable. Terminological axioms are defined to be of the form $C = D$ or $C \neq D$ where C, D are (complex) concepts. However, in most of the existing terminological representation systems terminological axioms are not allowed in this general form but they have to satisfy additional conditions. In 4.2 we investigate in which cases only such restricted terminological axioms have to be considered in order to test satisfiability of \mathcal{ALC} test sets. Finally, in 4.3 we will show how to decide whether or not a given \mathcal{ALC}_B -formula is logically entailed by a set of \mathcal{ALC}_B -formulas.

4.1 Computing \mathcal{ALC} test sets

The KD45-satisfiability algorithm presented in the previous section works in two phases: Firstly, a set of \mathcal{ALC} test sets is computed and then each of these sets is tested on satisfiability. The thereby used propagation rules of the frame algorithm have the advantage to “reflect” the properties of KD45 Kripke structures, and hence soundness and completeness of the KD45-satisfiability algorithm could be proved in a very natural way. On the other hand, the frame algorithm has the disadvantage to possibly construct a large number of labeled \mathcal{ALC}_B -formulas in order to determine a small number of \mathcal{ALC} test sets (cf. example 3.9).

In the following we will have a closer look at the properties of the frame algorithm. Building upon these properties we will then develop an algorithm which computes \mathcal{ALC} test sets immediately from the \mathcal{ALC}_B -formulas to be tested on KD45-satisfiability, i.e., we do no longer have to compute a (large) number of labels from which \mathcal{ALC} test sets then are extracted.

The main idea of this new approach is as follows: Suppose $o_1^*F_1, \dots, o_n^*F_n$ are \mathcal{ALC}_B -formulas which are to be tested on KD45-satisfiability, where each o_i^* is a (possibly empty) sequence of modalities and each F_i is an \mathcal{ALC} -formula. By looking at the structure of the sequences o_1^*, \dots, o_n^* we will then decide syntactically which elements in the power set of $\{F_1, \dots, F_n\}$ will be computed as an \mathcal{ALC} test set by the frame algorithm with input $o_1^*F_1 \parallel w_0, \dots, o_n^*F_n \parallel w_0$. The following example shows that this task is not obvious.

Example 4.1 Let A be a primitive concept and let C, D be concepts.

1. Let \mathcal{S} be the set $\{\Diamond_1(A = C), \Diamond_1(A = D)\}$. Applying the frame algorithm to the induced world constraint system $W = \{\Diamond_1(A = C) \parallel w_0, \Diamond_1(A = D) \parallel w_0\}$

results the derived system W' , given by the world constraints

$$\begin{array}{lll}
& \Diamond_1(A = C) \parallel w_0 & \Diamond_1(A = D) \parallel w_0 \\
w_0 \bowtie_1 w_1 & A = C \parallel w_1 & \Diamond_1(A = D) \parallel w_1 \\
w_0 \bowtie_1 w_2 & \Diamond_1(A = C) \parallel w_2 & A = D \parallel w_2 \\
w_1 \bowtie_1 w_3 & & A = D \parallel w_3 \\
w_2 \bowtie_1 w_4 & A = C \parallel w_4 &
\end{array}$$

The \mathcal{ALC} test set $A(w_0)$ of w_0 is empty, while $A(w_1) = A(w_4) = \{A = C\}$ and $A(w_2) = A(w_3) = \{A = D\}$.

2. On the other hand, starting with the set $S = \{\Box_1(A = C), \Box_1(A = D)\}$ leads to only one non-empty \mathcal{ALC} test set, namely $\{A = C, A = D\}$.

The leading sequences of modalities in the input \mathcal{ALC}_B -formulas obviously influence the determined \mathcal{ALC} test sets. ■

Let us firstly introduce the following notions. If W is a derived system, a subset consisting of world constraints $w_1 \bowtie_{i_1} w_2, \dots, w_{n-1} \bowtie_{i_{n-1}} w_n$ is called a *path* if the w_j are labels in W and each i_j is an agent. The *norm* of this path is given by the sequence $i_1 \dots i_n$, which arises from $i_1 \dots i_{n-1}$ by replacing each subsequence $i_j \dots i_j$ in $i_1 \dots i_{n-1}$ by i_j . For example, the norm of the path $w_1 \bowtie_{i_1} w_2, w_2 \bowtie_{i_1} w_3, w_3 \bowtie_{i_2} w_4, w_4 \bowtie_{i_2} w_5$ is given by $i_1 i_2$.

In the following we will present correlations between labeled \mathcal{ALC}_B -formulas which are taken as input of the frame algorithm and the labeled \mathcal{ALC}_B -formulas which are generated during the frame algorithm with this input. In order to simplify notation we thereby assume the input is of the form $\circ_1^* F_1 \parallel w_0, \dots, \circ_n^* F_n \parallel w_0$ where each \circ_i^* is a (possibly empty) sequence of modalities and each two \mathcal{ALC} -formulas F_i and F_j are syntactically different (e.g., by marking them with numbers). Since the presented correlations only depend on the modalities in the input \mathcal{ALC}_B -formulas this can be done without loss of generality.

Firstly, if W' is the result of an application of the frame algorithm, the next lemma states that there is a correlation between paths in W' and the sets of \mathcal{ALC}_B -formulas in W' .

Lemma 4.2 *Let W be a world constraint system which is induced by a finite set of \mathcal{ALC}_B -formulas, let W' be the result of the frame algorithm with input W , and let P be a path $w_0 \bowtie_{i_1} w_1, \dots, w_{n-1} \bowtie_{i_n} w_n$ in W' with norm $(P) = i_1 \dots i_n$. Then the set of labeled \mathcal{ALC}_B -formulas in W' with label w_n is a subset of $\{F \parallel w_n, \circ_{i_n} F \parallel w_n \mid \circ_{i_1} \dots \circ_{i_n} F \parallel w_0 \in W'\}$ where $\circ \in \{\Box, \Diamond\}$.*

Proof: Let $w_0 \bowtie_i u$ be in W' . If $\Diamond_i F_1 \parallel w_0, \dots, \Diamond_i F_n \parallel w_0, \Box_i G_1 \parallel w_0, \dots, \Box_i G_m \parallel w_0$ are the labeled \mathcal{ALC}_B -formulas in W' with label w_0 and with leading modality \Box_i or \Diamond_i , the labeled \mathcal{ALC}_B -formulas in W' with label u are a subset of S , given by

$$\{\Diamond_i F_1 \parallel u, \dots, \Diamond_i F_n \parallel u, \Box_i G_1 \parallel u, \dots, \Box_i G_m \parallel u, F_1 \parallel u, \dots, F_n \parallel u, G_1 \parallel u, \dots, G_m \parallel u\}$$

or, alternatively, by $\{F \parallel u, \Box_i F \parallel u \mid \Box_i F \parallel w_0 \in W'\}$. This follows immediately from the definition of the four propagation rules. Since we assumed \mathcal{ALC}_B -formulas to be in KD45 normal form, none of the world constraints $F_1 \parallel u, \dots, F_n \parallel u, G_1 \parallel u, \dots, G_m \parallel u$ has a leading modality \Box_i or \Diamond_i . Hence, if for some label u' there is $u \bowtie_i u'$ in W' , the labeled \mathcal{ALC}_B -formulas with label u' in W' also are a subset of S , and so on. Summing up, if there is a path $w_0 \bowtie_i w_{i_1}, \dots, w_{i_{k-1}} \bowtie_i w_k$ we know that the set of labeled \mathcal{ALC}_B -formulas with label w_k in W' is a subset of $S = \{F \parallel w_k, \Box_i F \parallel w_k \mid \Box_i F \parallel w_0 \in W'\}$.

Let us now consider a path P with $\text{norm}(P) = ij$. If there is a path $w_0 \bowtie_i w_{i_1}, \dots, w_{i_{k-1}} \bowtie_i w_k, w_k \bowtie_j w_{j_1}$, then the set of labeled \mathcal{ALC}_B -formulas with label w_k in W' is a subset of $S = \{F \parallel w_k, \Box_i F \parallel w_k \mid \Box_i F \parallel w_0 \in W'\}$. Hence, the set of labeled \mathcal{ALC}_B -formulas with label w_{j_1} in W' is a subset of $\{F \parallel w_{j_1}, \Box_j F \parallel w_{j_1} \mid \Box_j F \parallel w_k \in S\}$. Analogously to the argumentation above we obtain the following: If there is a path $w_0 \bowtie_i w_{i_1}, \dots, w_{i_{k-1}} \bowtie_i w_k, w_k \bowtie_j w_{j_1}, \dots, w_{j_{l-1}} \bowtie_j w_l$ in W' it follows that the set of labeled \mathcal{ALC}_B -formulas with label w_l in W' is a subset of $\{F \parallel w_l, \Box_j F \parallel w_l \mid \Box_j F \parallel w_k \in S\}$.

Since S consists only of elements $\Box_i F \parallel w_k$ or $F \parallel w_k$ such that $\Box_i F \parallel w_0$ is in W' it follows that $\Box_j F \parallel w_k \in S$ at most if $\Box_i \Box_j F \parallel w_0 \in W'$. That means, the set of labeled \mathcal{ALC}_B -formulas with label w_l in W' is a subset of $\{F \parallel w_l, \Box_j F \parallel w_l \mid \Box_i \Box_j F \parallel w_0 \in W'\}$.

With this argumentation the lemma follows immediately by induction. \square

In an \mathcal{ALC} test set only \mathcal{ALC} -formulas, i.e. \mathcal{ALC}_B -formulas without modalities, do occur. If W is a world constraint system and W' is the result of applying the frame algorithm to W , the next lemma provides a correlation between labeled \mathcal{ALC}_B -formulas in W' without modalities and labeled \mathcal{ALC}_B -formulas in W . Thereby, a sequence of modalities is called *sequence* for short. If S is the sequence $\Box_1 \dots \Box_n$, with $\Box \in \{\Diamond, \Box\}$, then $S[j]$ denotes \Box_j and $\text{indexes}(S)$ denotes $1, \dots, n$. Furthermore, an \mathcal{ALC}_B -formula F occurs in a derived system iff there is a (possibly empty) sequence \Box^* and a label w in W such that $\Box^* F \parallel w$ is in W .

Lemma 4.3 *Let W be a world constraint system which is induced by a finite number of \mathcal{ALC}_B -formulas, let W' be the result of the frame algorithm with input W , and let F and G be \mathcal{ALC} -formulas occurring in W . If there is a label w_n in W' such that $F \parallel w_n$ and $G \parallel w_n$ are both in W' , then there are two sequences S and S' such that (i) $SF \parallel w_0$ and $S'G \parallel w_0$ are in W and (ii) $\text{indexes}(S) = \text{indexes}(S')$.*

Proof: Let $w_0 \bowtie_{i_1} w_1, \dots, w_{n-1} \bowtie_{i_n} w_n$ be the path from w_0 to w_n in W' where each $w_{j-1} \bowtie_{i_j} w_j$ has been added into W' by an application of the \rightarrow_\Diamond or the \rightarrow_\Box rule. It is

sufficient to investigate this path since the \rightarrow_{\Diamond} and the \rightarrow_{\Box} rule do not add any new labeled \mathcal{ALC}_B -formulas at all, and it is easy to verify that there is exactly one such path.

If $F \parallel w_n$ and $G \parallel w_n$ are both in W' , there are labeled \mathcal{ALC}_B -formulas $\circ_{i_1} \dots \circ_{i_n} F \parallel w_0$ and $\circ_{i_1} \dots \circ_{i_n} G \parallel w_0$ in W where each \circ is either \Diamond or \Box and each i_j is an agent. This follows immediately by Lemma 4.2 and the fact that $H \parallel w_0 \in W'$ iff $H \parallel w_0 \in W$ for each \mathcal{ALC}_B -formula H since no propagation rule introduces a labeled \mathcal{ALC}_B -formula with label w_0 into a derived system. Thus, there are sequences S and S' such that $\text{indexes}(S) = \text{indexes}(S')$ and both $SF \parallel w_0$ and $S'G \parallel w_0$ are in W . \square

If W' is the result of applying the frame algorithm to a world constraint system W , the following proposition provides syntactical conditions that W satisfies whenever there is an \mathcal{ALC} test set $A(w)$ for some label w in W which contains a non-empty set of \mathcal{ALC} -axioms.

Proposition 4.4 *Let W be a world constraint system which is induced by a finite number of \mathcal{ALC}_B -formulas, let W' be the result of the frame algorithm with input W , and let F and G be \mathcal{ALC} -formulas occurring in W . If there is a label w_n in W' such that $F \parallel w_n$ and $G \parallel w_n$ are both in W' , then there are sequences S and S' such that (i) $\text{indexes}(S) = \text{indexes}(S')$, (ii) $SF \parallel w_0$ and $S'G \parallel w_0$ are in W , and (iii) there is no position j in S (respectively in S') such that $S[j] = S'[j] = \Diamond_{i_j}$ for some agent i_j .*

Proof: Because of Lemma 4.3 we know that there are sequences S and S' with $\text{indexes}(S) = \text{indexes}(S')$ and both $SF \parallel w_0$ and $S'G \parallel w_0$ are in W . Suppose j is the first position such that $S[j] = S'[j] = \Diamond_{i_j}$.

Starting with label w_0 we firstly show that for each path $w_0 \bowtie_{i_1} w_1, \dots, w_{n-1} \bowtie_{i_n} w_n$ in W' there is a label w_j in w_0, \dots, w_n such that $\Diamond_{i_j} \circ_1^* F \parallel w_j$ and $\Diamond_{i_j} \circ_2^* G \parallel w_j$ are in W' , where \circ_1^* and \circ_2^* are (possibly empty) sequences of modalities and $\text{indexes}(\circ_1^*) = \text{indexes}(\circ_2^*)$. Therefore we use the fact that for each position $k < j$ at least one of the modalities in $\{S[k], S'[k]\}$ is not a labeled \Diamond -modality.

If $\Diamond_{i_j} \circ_1^* F \parallel w_0$ and $\Diamond_{i_j} \circ_2^* G \parallel w_0$ are both in W' there is nothing to show. Otherwise, we have to distinguish two cases:

- (i) $\Diamond_i \circ^* \Diamond_{i_j} \circ_1^* F \parallel w_0$ and $\Box_i \circ^* \Diamond_{i_j} \circ_2^* G \parallel w_0$ are in W' or
- (ii) $\Box_i \circ^* \Diamond_{i_j} \circ_1^* F \parallel w_0$ and $\Box_i \circ^* \Diamond_{i_j} \circ_2^* G \parallel w_0$ are in W' ,

where i is an agent and \circ^* is a (possibly empty) sequence of modalities.⁵ Without loss of generality we do not consider the case where $\Box_i \circ^* \Diamond_{i_j} \circ_1^* F \parallel w_0$ and $\Diamond_i \circ^* \Diamond_{i_j} \circ_2^* G \parallel w_0$

⁵In order to simplify notation we will use \circ^* in both the sequences of F and of G . Formally we had to distinguish these occurrences since everything we know about them is that they do not differ in their indexes. However, here it should always be clear by context which occurrence is meant.

are in W' , since it is symmetrical to case (i). Note, that applying propagation rules to labeled \mathcal{ALC}_B -formulas with label w_0 whose leading modalities are not indexed with i —obtaining, say, label v —cannot result the first element in a path leading to a label w such that $F \parallel w$ and $G \parallel w$ are both in W' . This is due to the fact that in this case neither F nor G occurs in labeled \mathcal{ALC}_B -formulas labeled with v .

For case (i), after an application of the \rightarrow_\diamond or the \rightarrow_{\diamond_0} rule to $\diamond_i o^* \diamond_{i_j} o_1^* F \parallel w_0$, for some label v the world constraints

$$o^* \diamond_{i_j} o_1^* F \parallel v, \quad o^* \diamond_{i_j} o_2^* G \parallel v, \quad \Box_i o^* \diamond_{i_j} o_2^* G \parallel v \quad (\alpha)$$

are exactly the labeled \mathcal{ALC}_B -formulas in the currently derived world constraint system which are labeled with v and contain F or G . Analogously, after an application of the \rightarrow_\diamond or the \rightarrow_{\diamond_0} rule to $\diamond_i H \parallel w_0$ —where H is different from $o^* \diamond_{i_j} o_1^* F$ —for some label v the world constraints

$$\diamond_i o^* \diamond_{i_j} o_1^* F \parallel v, \quad o^* \diamond_{i_j} o_2^* G \parallel v, \quad \Box_i o^* \diamond_{i_j} o_2^* G \parallel v \quad (\beta)$$

are exactly the labeled \mathcal{ALC}_B -formulas in the currently derived world constraint system which are labeled with v and contain F or G .

For case (ii), after an application of the \rightarrow_\Box or the \rightarrow_{\Box_0} rule to labeled \mathcal{ALC}_B -formulas with label w_0 whose leading modalities are indexed by i , for some label v the world constraints

$$\Box_i o^* \diamond_{i_j} o_1^* F \parallel v, \quad o^* \diamond_{i_j} o_1^* F \parallel v, \quad \Box_i o^* \diamond_{i_j} o_2^* G \parallel v, \quad o^* \diamond_{i_j} o_2^* G \parallel v \quad (\gamma)$$

are exactly the labeled \mathcal{ALC}_B -formulas in the currently derived world constraint system which are labeled with v and contain F or G . And, finally, after an application of the \rightarrow_\diamond or the \rightarrow_{\diamond_0} rule to some labeled \mathcal{ALC}_B -formula $\diamond_i H \parallel w_0$ we obtain the same labeled \mathcal{ALC}_B -formulas as in case (γ) above.

Since we assumed \mathcal{ALC}_B -formulas to be in KD45 normal form, the leading modality in o^* is not indexed with i . Hence, for case (α), propagation rule applications to $\Box_i o^* \diamond_{i_j} o_2^* G \parallel v$ cannot result the first element in a path leading to a label w such that $F \parallel w$ and $G \parallel w$ are both in W' . Analogously, for case (β), applying propagation rules to $o^* \diamond_{i_j} o_2^* G \parallel v$ cannot lead to a label w such that $F \parallel w$ and $G \parallel w$ are both in W' . For the same reason in case (γ) either rule applications to labeled \mathcal{ALC}_B -formulas with leading modality i and label v , or to labeled \mathcal{ALC}_B -formulas whose leading modality is the first modality in o^* and label v , may lead to a label w such that $F \parallel w$ and $G \parallel w$ are both in W' .

Summing up, if v is the first element in a path $w_0 \bowtie_{i_1}, \dots, \bowtie_{i_n} w_n$, then either

- (1) $o^* \diamond_{i_j} o_1^* F \parallel v$ and $o^* \diamond_{i_j} o_2^* G \parallel v$ are in W' or
- (2) $\diamond_i o^* \diamond_{i_j} o_1^* F \parallel v$ and $\Box_i o^* \diamond_{i_j} o_2^* G \parallel v$ are in W' or
- (3) $\Box_i o^* \diamond_{i_j} o_1^* F \parallel v$ and $\Box_i o^* \diamond_{i_j} o_2^* G \parallel v$ are in W'

Hence, to v the same argumentation is applicable as to label w_0 above.

Let now w_j be a label in w_0, \dots, w_n such that $\Diamond_{i_j} o_1^* F \parallel w_j$ and $\Diamond_{i_j} o_2^* G \parallel w_j$ are in W' . Without loss of generality during the frame algorithm the \rightarrow_{\Diamond} (respectively the \rightarrow_{\Diamond_o}) rule has been applied to $\Diamond_{i_j} o_1^* F \parallel w_j$ before it has been applied to $\Diamond_{i_j} o_2^* G \parallel w_j$. If o_1^* (and hence o_2^*) is the empty sequence, then $F \parallel v$ and $\Diamond_{i_j} G \parallel v$ are in W' , where $w_j \bowtie_{i_j} v$ has been introduced by this rule application. But $F \parallel w_n$ cannot be in W' because of the remaining application of the \rightarrow_{\Diamond} rule to $\Diamond_{i_j} o^* G \parallel w_j$. On the other hand, if o_1^* (and hence o_2^*) is not empty, its leading modality is different from i_j (since \mathcal{ALC}_B -formulas are in KD45 normal form) and thus there cannot be a label w_n in W' such that $F \parallel w_n$ and $G \parallel w_n$ are both in W' because of the remaining application of the \rightarrow_{\Diamond} or the \rightarrow_{\Diamond_o} rule to $\Diamond_{i_j} o^* G \parallel w_j$.

In both cases this contradicts the assumption that there is a label w_n in W' such that $F \parallel w_n$ and $G \parallel w_n$ are both in W' , i.e., there cannot exist a position j such that $S[j] = S'[j] = \Diamond_{i_j}$ for some agent i_j . \square

Finally, we present a proposition which states that also the opposite direction of proposition 4.4 holds.

Proposition 4.5 *Let W be a world constraint system which is induced by a finite set of \mathcal{ALC}_B -formulas, let W' be the result of the frame algorithm with input W , and let F and G be \mathcal{ALC} -formulas which occur in W . If $SF \parallel w_0$ and $S'G \parallel w_0$ are in W , where S, S' are sequences such that $\text{indexes}(S) = \text{indexes}(S')$ and there is no position j in S (respectively in S') such that $S[j] = S'[j] = \Diamond_{i_j}$ for some agent i_j , then there is a label w_n in W' such that $F \parallel w_n$ and $G \parallel w_n$ are both in W' .*

Proof: Let S_* and S'_* arise from S and S' by deleting the first modality, respectively. If $SF \parallel w_0$ and $S'G \parallel w_0$ are in W , there are three cases:

1. $S[1] = S'[1] = \Box_{i_1}$ and the \rightarrow_{\Box} or the \rightarrow_{\Box_o} rule is applied to \mathcal{ALC}_B -formulas of the form $\Box_{i_1} H \parallel w_0$. In this case $SF \parallel w_1, S'G \parallel w_1, S_*F \parallel w_1$, and $S'_*G \parallel w_1$ are all in W' if $w_0 \bowtie_{i_1} w_1$ is introduced by this rule application.
2. Without loss of generality $S[1] = \Diamond_{i_1}, S'[1] = \Box_{i_1}$, and the \rightarrow_{\Diamond} or the \rightarrow_{\Diamond_o} rule is applied to some \mathcal{ALC}_B -formula $\Diamond_{i_1} H \parallel w_0$, where H is different from F . In this case $SF \parallel w_1$ and $S'G \parallel w_1$ are in W' if $w_0 \bowtie_{i_1} w_1$ is introduced by this rule application.
3. Without loss of generality $S[1] = \Diamond_{i_1}, S'[1] = \Box_{i_1}$, and the \rightarrow_{\Diamond} or the \rightarrow_{\Diamond_o} rule is applied to $SF \parallel w_0$. In this case $S_*F \parallel w_1$ and $S'_*G \parallel w_1$ are both in W' if $w_0 \bowtie_{i_1} w_1$ is introduced by this rule application. It does not influence the argumentation that $S'G \parallel w_1$ is also in W' .

For label w_1 the argumentation is the same such that the proposition follows immediately by induction. \square

Propositions 4.4 and 4.5 together give us an answer to the question which \mathcal{ALC} test sets are generated by an application of the frame algorithm. We summarize this result in the following theorem.

Theorem 4.6 *Let W be a world constraint system which is induced by a finite set of \mathcal{ALC}_B -formulas and let W' be the result of the frame algorithm with input W . Then there is a label w in W' such that the \mathcal{ALC} test set $A(w)$ contains the \mathcal{ALC} -formulas F_1, \dots, F_n iff there are sequences S_1, \dots, S_n in W such that*

- $S_1 F_1 \parallel w_0, \dots, S_n F_n \parallel w_0$ are in W ,
- $\text{indexes}(S_1) = \dots = \text{indexes}(S_n)$, and
- there is no position j such that for two sequences S' and S'' in $\{S_1, \dots, S_n\}$ holds $S'[j] = S''[j] = \Diamond_{i_j}$ for some agent i_j .

Proof: If the \mathcal{ALC} test set $A(w)$ contains the \mathcal{ALC} -axioms F_1, \dots, F_n , none of the F_j does contain modalities. Thus, because of Proposition 4.4 there are sequences S_1, \dots, S_n such that $S_1 F_1 \parallel w_0, \dots, S_n F_n \parallel w_0$ are in W and $\text{indexes}(S_1) = \dots = \text{indexes}(S_n)$. Now suppose, that $S' F' \parallel w_0$ and $S'' F'' \parallel w_0$ with F' and F'' in $\{F_1, \dots, F_n\}$ are in W such that $S'[j] = S''[j] = \Diamond_{i_j}$ for some agent i_j . Again because of Proposition 4.4, we can conclude that either $F' \parallel w$ or $F'' \parallel w$ is not in W' . This contradicts that $A(w)$ contains F_1, \dots, F_n . The other direction follows immediately from Proposition 4.5. \square

An optimized version of the KD45-satisfiability algorithm is given in Figure 3. For the input of this algorithm remember that each \mathcal{ALC}_B -formula is of the form SF where S is a (possibly empty) sequence of modalities and F is an \mathcal{ALC} -formula. In step 2. of the algorithm for each sequence S_i in $\{S_1, \dots, S_n\}$ the set S_i of all \mathcal{ALC} -formulas in $\{F_1, \dots, F_n\}$ is computed such that the conditions of Theorem 4.6 are satisfied. Multiple generations of the same \mathcal{ALC} test sets are avoided by testing whether or not the \mathcal{ALC} -formula F_i is already in an \mathcal{ALC} test set S_k .

Summing up, soundness, completeness, and termination of this algorithm follow immediately from the results in Section 3 and in this subsection. It is easy to verify that step 2. of this optimized KD45-satisfiability algorithm works in polynomial time in the length of the input \mathcal{ALC}_B -formulas and generates a polynomial number of \mathcal{ALC} test sets. Thus, in contrast to the KD45-satisfiability algorithm in the previous section, it does not generate a worst case exponential overhead of \mathcal{ALC}_B -formulas.

Let us now reconsider example 3.9: Given the \mathcal{ALC}_B -formulas $\Diamond_1 F_1$, $\Diamond_1 F_2$, and $\Diamond_1 F_3$ the optimized KD45-satisfiability algorithm constructs the \mathcal{ALC} test sets $\{F_1\}$,

1. Let S_1F_1, \dots, S_nF_n be the \mathcal{ALC}_B -formulas to be tested on KD45-satisfiability, where each S_i is a (possibly empty) sequence of modalities and each F_i is an \mathcal{ALC} -formula.
2. for $i := 1, \dots, n$ do
 - if $F_i \in S_k$ for some $k \in \{1, \dots, i-1\}$
 - then $S_i := \emptyset$
 - else
 - $S_i := \{F_i\}$
 - for $j := i+1, \dots, n$ do
 - if $\text{indexes}(S_i) = \text{indexes}(S_j)$ and
 - there is no position k such that $S_i[k] = S_j[k] = \Diamond_l$ for some agent l
 - then $S_i := S_i \cup \{F_j\}$
3. For each non-empty set s in $\{S_1, \dots, S_n\}$ do: If s is not satisfiable, then STOP and return "KD45-unsatisfiable".
4. Return "KD45-satisfiable".

Figure 3: The optimized KD45-satisfiability algorithm.

$\{F_2\}$, and $\{F_3\}$ without applying the frame algorithm. Analogously, from $\Box_1(A = C)$ and $\Box_1(A = D)$ the only \mathcal{ALC} test set $\{A = C, A = D\}$ is generated immediately from the syntactical structure of the input \mathcal{ALC}_B -formulas (cf. example 4.1).

4.2 Restricted \mathcal{ALC} -TBoxes

In Section 2 we defined an \mathcal{ALC} -TBox as a set of terminological axioms of the form $C = D$ and $C \neq D$, where C and D are concepts. However, most of the existing terminological representation and inference systems (e.g., `BACK` [30], `CLASSIC` [7], `KRIS` [2]) only allow terminological axioms of the form $A = C$, where A is a primitive concept and C is a concept. Such a terminological axiom is called (*concept*) *definition* of A .⁶ Building upon this, an \mathcal{ALC} -TBox is then defined as a finite set of terminological axioms which satisfies the following restrictions:

- each atomic concept appears at most once as the left hand side of a terminological axiom, and

⁶Often so-called concept specializations of the form $A \sqsubseteq C$ are allowed which abbreviate the terminological axiom $A = C \sqcap A^*$ where A^* is a new primitive concept.

- in this set cycles do not occur.

Thereby, a set \mathcal{S} of terminological axioms contains a *cycle* iff there exists a terminological axiom $A = C$ in \mathcal{S} such that A occurs in the concept C' which arises from C by iterated substitutions of primitive concepts in C by the right hand sides of their definition in \mathcal{S} . For example, if A and B are primitive concepts the sets $\{A = A\}$ and $\{A = C \sqcap B, B = D \sqcup \exists R.A\}$ of terminological axioms contain cycles. In the following we will call \mathcal{ALC} -TBoxes satisfying the additional conditions described above *restricted \mathcal{ALC} -TBoxes* in order to distinguish them from the \mathcal{ALC} -TBoxes defined in Section 2.

It can be shown that each restricted \mathcal{ALC} -TBox \mathcal{T} can be transformed into an equivalent restricted \mathcal{ALC} -TBox \mathcal{T}' such that each right hand side of a concept definition in \mathcal{T}' does only contain concepts which do not occur as a left hand side in \mathcal{T}' (see, e.g., [20]). For example, if A_1, A_2, A_3 are primitive concepts, the restricted \mathcal{ALC} -TBox $\mathcal{T} = \{A_1 = A_2 \sqcap A_3, A_2 = C \sqcup D, A_3 = \exists R.C\}$ can be transformed into $\mathcal{T}' = \{A_1 = (C \sqcup D) \sqcap \exists R.C, A_2 = C \sqcup D, A_3 = \exists R.C\}$. Thus, each primitive concept A on the left hand side of a terminological axiom $A = C$ in \mathcal{T}' can be seen as an abbreviation for the concept C . With this it is easy to verify that testing consistency of an \mathcal{ALC} -ABox \mathcal{A} w.r.t. a restricted \mathcal{ALC} -TBox \mathcal{T} is equivalent to only testing consistency of an \mathcal{ALC} -ABox \mathcal{A}' . Thereby, \mathcal{A}' arises from \mathcal{A} by successively replacing all primitive concepts by the right hand sides of their definitions in \mathcal{T} . The size of \mathcal{A}' is worst case exponential in the size of \mathcal{A} and \mathcal{T} (see, e.g., [20]) and testing satisfiability of \mathcal{A}' is known to be PSPACE-complete (see [19]). Possible optimizations are discussed in [1].

An algorithm for testing consistency of \mathcal{ALC} test sets which may contain terminological axioms as defined in Section 2 has been given in [21]. As an easy consequence of a result by Fischer and Ladner [13] this test is EXP-TIME complete. Moreover, when using more expressive terminological logics than \mathcal{ALC} this test becomes undecidable (for the terminological logic \mathcal{ALCF} this has been shown in [4]), while this is not the case when using restricted \mathcal{ALC} -TBoxes only.

Let now \mathcal{S} be a set of \mathcal{ALC}_B -formulas. Because of the above given discussion on efficiency and decidability of testing satisfiability of an \mathcal{ALC} -TBox and an \mathcal{ALC} -ABox it is an interesting question whether or not the terminological axioms in each \mathcal{ALC} test set which is generated from \mathcal{S} define a restricted \mathcal{ALC} -TBox. The answer to this question can be given with the help of Theorem 4.6 which can be used to formulate sufficient syntactical conditions which—if satisfied—guarantee that only restricted \mathcal{ALC} -TBoxes have to be tested in order to test KD45-satisfiability of a set of \mathcal{ALC}_B -formulas. For example, these conditions could be given by

1. agents only have positive beliefs, i.e., negation signs do not occur in front of \Box -operators, and
2. for each sequence S of modalities holds that the set of \mathcal{ALC}_B -formulas occurring in the scope of S define a restricted \mathcal{ALC} -TBox.

For practical applications, however, such conditions seem not reasonable and, even worse, when computing logical consequences (see next subsection) such syntactical conditions in general cannot be maintained. Hence, for testing satisfiability of \mathcal{ALC} test sets we in general have to take terminological axioms into account which do not define a restricted \mathcal{ALC} -TBox.

4.3 Computing \mathcal{ALC}_B -Inferences

We will now show how to use the KD45-satisfiability algorithm in order to test whether or not a given \mathcal{ALC}_B -formula is a logical consequence from a set F_1, \dots, F_n of \mathcal{ALC}_B -formulas. Again, we are only interested in KD45 Kripke structures and thus define: F is a *KD45 consequence* of F_1, \dots, F_n iff for each KD45 Kripke structure $K = (\mathcal{W}, \Gamma, K_I)$ and for each world w in \mathcal{W} holds: if $K, w \models F_1, \dots, F_n$, then $K, w \models F$. Firstly, let F be an \mathcal{ALC}_B -formula of the form $\circ^*(C = D)$, $\circ^*(C \neq D)$, or $\circ^*(a : C)$, where \circ^* is an abbreviation for a (possibly empty) sequence of modalities. Then, F is an KD45 consequence of F_1, \dots, F_n iff the set $F_1, \dots, F_n, [\neg F]^*$ of \mathcal{ALC}_B -formulas is not KD45-satisfiable, where $[\neg F]^*$ denotes the negation normal form of $\neg F$. Note, that $\neg F$ is an \mathcal{ALC}_B -formula if F is of the above described form.

If, on the other hand, F is of the form $\Box^*(aRb)$, where \Box^* is an abbreviation for a (possibly empty) sequence of non-negated indexed \Box operators, we cannot use this test method since negation signs are not allowed in \mathcal{ALC}_B -formulas which contain a role instance. To handle this case, we extend the notion of \mathcal{ALC}_K -formulas as follows: if R is a role, a, b are objects, and i_1, \dots, i_m are agents, then $\Diamond_{i_1} \dots \Diamond_{i_m}(aRb)$ is an \mathcal{ALC}_B -formula.

Alternatively, these \mathcal{ALC}_B -formulas could be defined by $\circ_{i_1} \dots \circ_{i_m}(aR'b)$ where (i) each \circ_{i_j} is either \Box_{i_j} or $\neg\Box_{i_j}$, (ii) R' is either R or $\neg R$, and (iii) the number of negation signs in $\circ_{i_1} \dots \circ_{i_m}(aR'b)$ is even. Using this definition it is easy to see that the negation normal form of the new \mathcal{ALC}_B -formulas does not contain negation of roles. Therefore, on a technical level we could allow such formulas as \mathcal{ALC}_K -axioms in Section 2. But a restriction like “the number of negation signs is even” seems not to be adequate when defining a language to describe beliefs of agents. However, for testing whether or not an \mathcal{ALC}_B -formula is entailed by a set F_1, \dots, F_n of \mathcal{ALC}_B -axioms, this definition turns out to be reasonable.

Note, that KD45-satisfiability of a set of \mathcal{ALC}_B -formulas can be handled by the KD45-satisfiability algorithm in Section 3 even if we use the above introduced extended definition of \mathcal{ALC}_B -formulas: Firstly, the algorithm only treats the modalities of \mathcal{ALC}_B -formulas, i.e., it works independently from the syntactical structure of formulas without modalities. Secondly, satisfiability of the resulting \mathcal{ALC} test sets still can be tested, since they do not contain negation of roles. And, finally, it does not matter whether aRb is in an \mathcal{ALC} test set because of an input \mathcal{ALC}_B -formula $\Box_{i_1} \dots \Box_{i_m}(aRb)$, or because

of an input \mathcal{ALC}_B -formula $\Diamond_{i_1} \dots \Diamond_{i_m}(aRb)$. Summing up, when using the extended definition of \mathcal{ALC}_B -formulas we need not to change the KD45-satisfiability algorithm at all. The following proposition provides a test whether or not an \mathcal{ALC}_B -formula $\Box_{i_1} \dots \Box_{i_n}(aRb)$ is entailed by a set of \mathcal{ALC}_B -formulas.

Proposition 4.7 *Let F_1, \dots, F_n be a finite set of \mathcal{ALC}_B -formulas such that F_1, \dots, F_n are KD45-satisfiable, let a and b be objects, let R be a role, and let i_1, \dots, i_m be agents. Then $\Box_{i_1} \dots \Box_{i_m}(aRb)$ is a KD45 consequence of F_1, \dots, F_n iff $\Box_{i_1} \dots \Box_{i_m}(aRb)$ is one of the \mathcal{ALC}_B -formulas in F_1, \dots, F_n .*

Proof: The test whether or not $\Box_{i_1} \dots \Box_{i_m}(aRb)$ is a KD45 consequence of F_1, \dots, F_n is equivalent to testing whether or not $F_1, \dots, F_n, \Diamond_{i_1} \dots \Diamond_{i_m}(a \neg Rb)$ is KD45-satisfiable. However, since $\Diamond_{i_1} \dots \Diamond_{i_m}(a \neg Rb)$ is not an \mathcal{ALC}_B -formula, this case cannot be handled by the KD45-satisfiability algorithm of Section 3. Alternatively, let us have a look at the application of the frame algorithm to the world constraint system W which is induced by $\{F_1, \dots, F_n, \Diamond_{i_1} \dots \Diamond_{i_m}(aR'b)\}$, where R' is a role which does not occur in F_1, \dots, F_n . Since F_1, \dots, F_n are KD45-satisfiable, each \mathcal{ALC} test set is satisfiable which is constructed by applying the frame algorithm to the world constraint system which is induced by $\{F_1, \dots, F_n\}$. Hence, it is easy to verify that the \mathcal{ALC}_B -formulas $F_1, \dots, F_n, \Diamond_{i_1} \dots \Diamond_{i_m}(aR'b)$ are KD45-satisfiable.

Let us now consider R' as an abbreviation for $\neg R$. Obviously, this does not influence the result of applying the frame algorithm to W , but it may influence satisfiability of the thereby computed \mathcal{ALC} test sets—say $A(w_1), \dots, A(w_k)$ —in which $aR'b$ occurs. It is easy to verify that, given an \mathcal{ALC} -TBox \mathcal{T} and an \mathcal{ALC} -ABox \mathcal{A} , the \mathcal{ALC} -formula aRb is logically entailed by \mathcal{T} and \mathcal{A} iff $aRb \in \mathcal{A}$. Hence, an \mathcal{ALC} test set $A(w_i) \in \{A(w_1), \dots, A(w_k)\}$ is unsatisfiable iff aRb and $aR'b$ are both in $A(w_i)$. Because of Theorem 4.6 the \mathcal{ALC} -formulas aRb and $aR'b$ can only be in the same \mathcal{ALC} test set if there is a labeled \mathcal{ALC}_B -formula $\Box_{i_1} \dots \Box_{i_m}(aRb)$ in W . From this the proposition follows immediately. \square

Summing up, we have now given algorithms for deciding KD45-satisfiability of a given set of \mathcal{ALC}_B -formulas, and, building upon this, for deciding whether or not a given \mathcal{ALC}_B -formula F is a KD45 consequence of a given set F_1, \dots, F_n of \mathcal{ALC}_B -formulas.

5 Conclusion

We have presented a two-dimensional extension of the concept language \mathcal{ALC} which allows both reasoning on the objective level and reasoning on the level of epistemic alternatives. In the obtained language \mathcal{ALC}_B , a world agents are acting in can be described by a set of terminological and assertional axioms. Furthermore, the beliefs agents have about this world, about the beliefs of other agents, and about their own

beliefs can be described by terminological and assertional axioms with a leading indexed \Box operator or a leading sequence of indexed \Box operators. We presented sound and complete algorithms to check consistency of the represented beliefs and to decide whether an \mathcal{ALC}_B -formula is logically entailed by a given set of \mathcal{ALC}_B -formulas. Thus, it is possible to equip agents with a decidable component to represent beliefs that is much more expressive than representing beliefs via propositional logic. Since the actions a single agent can perform are essentially based on his local beliefs this component can be seen as one of the basic parts in the architecture of agents.

The main restriction of the presented language \mathcal{ALC}_B lies in the fact that modalities are only allowed in front of terminological and assertional axioms. As an extension one might think of modalities in front of concepts as well. Such a language would allow to represent facts like “the things agent i believes to be expensive are exactly the things agent j believes to be cheap” by

$$\Box_i(\text{expensive}) = \Box_j(\text{cheap}).$$

Such an extended language, however, causes algorithmic problems that are beyond the scope of this paper and is currently investigated.

Acknowledgements. I am grateful to Hans-Jürgen Bärkert, Bernhard Hollunder, and Andreas Nonnengart for many discussions on the topic of this paper and for reading earlier drafts.

5 Conclusion

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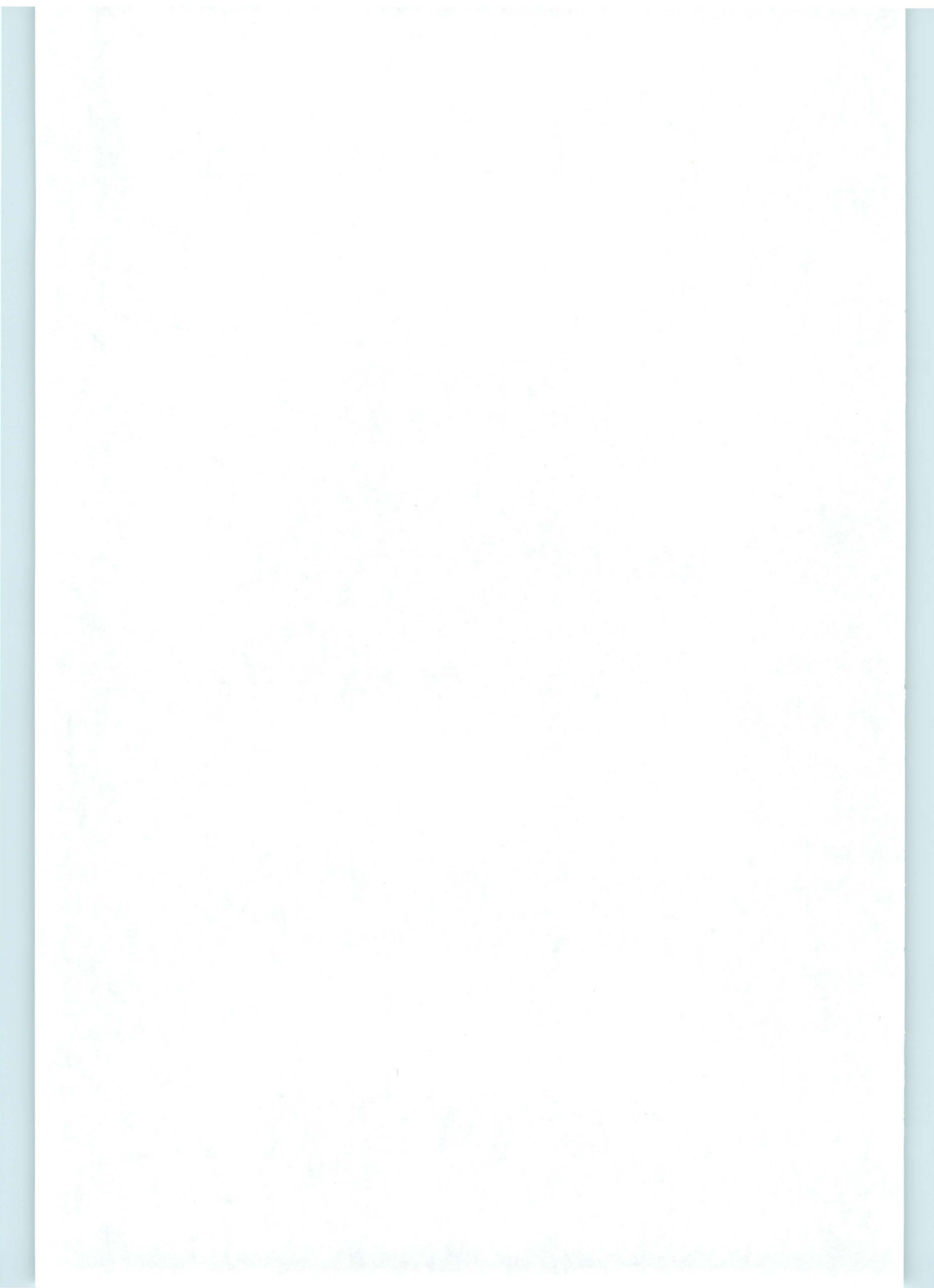
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